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Decomposing tournaments into paths

Allan Lo, Viresh Patel, Jozef Skokan and John Talbot

ABSTRACT

We consider a generalisation of Kelly’s conjecture which is due to Alspach, Mason, and Pullman from 1976. Kelly’s conjecture states that every regular tournament has an edge decomposition into Hamilton cycles, and this was proved by Kühn and Osthus for large tournaments. The conjecture of Alspach, Mason, and Pullman asks for the minimum number of paths needed in a path decomposition of a general tournament T . There is a natural lower bound for this number in terms of the degree sequence of T and it is conjectured that this bound is correct for tournaments of even order. Almost all cases of the conjecture are open and we prove many of them.

1. Introduction

There has been a great deal of recent activity in the study of decompositions of graphs and hypergraphs. The prototypical question in this area asks whether, for some given class \mathcal{C} of graphs, hypergraphs or directed graphs, the edge set of each $H \in \mathcal{C}$ can be decomposed into parts satisfying some given property. The development of the robust expanders technique by Kühn and Osthus [9] was a major breakthrough leading to the resolution of several conjectures concerning decompositions of (directed) graphs into spanning structures such as matchings and Hamilton cycles (see, for example, [4, 10]).

The problem we address in this paper is that of decomposing tournaments into directed paths. A *tournament* is an orientation of the complete graph, that is, one obtains a tournament by assigning a direction to each edge of the (undirected) complete graph. Let us begin, however, in the more general setting of directed graphs.

Let D be a directed graph with vertex set $V(D)$ and edge set $E(D)$. When referring to paths and cycles in directed graphs, we always mean directed paths and directed cycles. A *path decomposition* of D is a collection of paths P_1, \dots, P_k of D whose edge sets $E(P_1), \dots, E(P_k)$ partition $E(D)$. Given any directed graph D , it is natural to ask what the minimum number of paths is in a path decomposition of D . This is called the *path number* of D and is denoted $\text{pn}(D)$. A natural lower bound on $\text{pn}(D)$ is obtained by examining the degree sequence of D . For each vertex $v \in V(D)$, write $d_D^+(v)$ (respectively, $d_D^-(v)$) for the number of edges exiting (respectively, entering) v . Define the *excess* at v to be $\text{ex}(v) := d_D^+(v) - d_D^-(v)$ and similarly define the *positive* and *negative excess* at v to be, respectively, $\text{ex}^+(v) := \max\{\text{ex}(v), 0\}$

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and $\text{ex}^-(v) := \max\{-\text{ex}(v), 0\}$. It is easy to see that the excesses of all vertices sum to zero.

We note that in any path decomposition of D , at least $\text{ex}^+(v)$ paths must start at v and at least $\text{ex}^-(v)$ paths must end at v . Therefore, we have

$$\text{pn}(D) \geq \text{ex}(D) := \sum_{v \in V(D)} \text{ex}^+(v) = \sum_{v \in V(D)} \text{ex}^-(v) = \frac{1}{2} \sum_{v \in V(D)} |\text{ex}(v)|,$$

where $\text{ex}(D)$ is called the excess of D . Any digraph for which equality holds above is called *consistent*. Clearly, not every digraph is consistent; in particular, any nonempty digraph D of excess 0 cannot be consistent. However, Alspach, Mason, and Pullman [1] conjectured that every even tournament is consistent.

CONJECTURE 1.1 (Alspach, Mason, and Pullman [1]). Every tournament T with an even number of vertices satisfies $\text{pn}(T) = \text{ex}(T)$.

It is almost immediate to see that this conjecture is a considerable generalisation of Kelly’s conjecture stated below. We give the easy argument after Theorem 1.3.

CONJECTURE 1.2 (Kelly; see, for example, [3]). The edge set of every regular tournament can be decomposed into Hamilton cycles.

Almost 50 years after it was stated, Kühn and Osthus [9] finally proved Kelly’s conjecture for large tournaments using their powerful robust expanders technique, which was subsequently used to prove several other conjectures on edge decompositions of (directed) graphs [4, 10].

THEOREM 1.3 (Kühn and Osthus [9]). *Every sufficiently large regular tournament has a Hamilton decomposition.*

To see that Conjecture 1.1 implies Conjecture 1.2, take any regular $(n + 1)$ -vertex tournament T (so n must be even) and any $v \in V(T)$, and note that $\text{ex}(T - v) = n/2$. If Conjecture 1.1 holds, then $T - v$ can be decomposed into $n/2$ paths, so they must be Hamilton paths. Adding v back to $T - v$, it is easy to see that each path can be completed to a Hamilton cycle, giving a Hamilton decomposition of T . The converse is also easy to see. Thus, the special case of Conjecture 1.1 in which $\text{ex}(T) = n/2$ is equivalent to Kelly’s Conjecture. In general, however, $\text{ex}(T)$ can take a large range of values as the proposition below shows.

PROPOSITION 1.4. *If T is an n -vertex tournament with n even, then $n/2 \leq \text{ex}(T) \leq n^2/4$. Furthermore, each value in the range occurs.*

As we saw, the lower bound occurs for any almost-regular tournament and it is easy to verify that the upper bound occurs for the transitive tournament (in fact it occurs for any tournament with a vertex partition into two equal size parts A and B where all edges are directed from A to B). Alspach and Pullman [2] showed that for any tournament T , $\text{pn}(T) \leq n^2/4$ thus verifying Conjecture 1.1 for the special case $\text{ex}(T) = n^2/4$ (and this was generalised to digraphs [13]). Thus, the conjecture has been solved for the two extreme values of excess, namely, $n/2$ and $n^2/4$: for every other value of $\text{ex}(T)$ between $n/2$ and $n^2/4$ the conjecture remains open. Our main contribution is to solve many more cases of the conjecture.

THEOREM 1.5. *There exists $C > 0$ and $n_0 \in \mathbb{N}$ such that if T is an n -vertex tournament with $n \geq n_0$ even and $\text{ex}(T) > Cn$, then $\text{pn}(T) = \text{ex}(T)$.*

We make no attempt to optimise or even compute the value of C , but we note that it is not a tower-type constant that one associates with the use of the Szemerédi Regularity Lemma (a powerful tool used in extremal combinatorics and elsewhere). We prove this theorem in two steps. We will first prove the following weakening of the Theorem 1.5.

THEOREM 1.6. *There exists $\varepsilon > 0$ (we can take $\varepsilon = 1/18$) and $n_0 \in \mathbb{N}$ such that if T is a tournament on $n > n_0$ vertices with n even and $\text{ex}(T) \geq n^{2-\varepsilon}$, then $\text{pn}(T) = \text{ex}(T)$.*

The proof of this result is short and self-contained, relying on a novel application of the absorption technique due to Rödl, Ruciński, and Szemerédi [14] (with special forms appearing in earlier work, for example, [7]).

In the next step, we consider tournaments of excess smaller than $n^{2-\varepsilon}$ but bigger than Cn . Such tournaments are almost regular and are therefore amenable to the techniques used by Kühn and Osthus [9]. For tournaments of small excess, we will ultimately reduce the problem of showing that $\text{pn}(T) = \text{ex}(T)$ to the problem of showing that a regular oriented graph D of very high degree has an edge decomposition into Hamilton cycles; such a Hamilton decomposition of D is known to exist by the main result from [9].

1.1. Outline

In the next section, we give the basic notation we will use as well as some preliminary results needed in Section 3. In Section 3, we give the short proof of Theorem 1.6, which requires only Hall's Theorem and Menger's Theorem. In Section 4, we give further preliminaries needed for the remaining sections; in particular, we state the results related to robust expansion that we will need. At the end of Section 4, we give an overview of the arguments in Section 5 and Section 6 that allow us to extend Theorem 1.6 to Theorem 1.5. Section 5 contains a preliminary result, Lemma 5.1, that helps us to deal with certain problematic vertices that we encounter in Section 6. In Section 6, we prove Theorem 1.5 in a three-step reduction via Theorems 6.1, 6.7, and 6.12.

2. Notation and preliminaries

2.1. Notation

In this paper, a digraph refers to a directed graph without loops where we allow up to two edges between any pair x, y of distinct vertices, at most one in each direction. Occasionally we work with directed multigraphs which again have no loops, but where we permit more than two directed edges between any pair of distinct vertices. An oriented graph is a directed graph where we permit only one edge between any pair of distinct vertices. Given a digraph D , we write $V(D)$ for its vertex set and $E(D)$ for its edge set. We write xy for an edge directed from x to y .

We write $H \subseteq D$ to mean that H is a subdigraph of D , that is, $V(H) \subseteq V(D)$ and $E(H) \subseteq E(D)$. Given $X \subseteq V(D)$, we write $D - X$ for the digraph obtained from D by deleting all vertices in X , and $D[X]$ for the subdigraph of D induced by X . Given $F \subseteq E(D)$, we write $D - F$ for the digraph obtained from D by deleting all edges in F . If H is a subdigraph of D , we write $D - H$ for $D - E(H)$. For two subdigraphs H_1 and H_2 of D , we write $H_1 \cup H_2$ for the subdigraph with vertex set $V(H_1) \cup V(H_2)$ and edge set $E(H_1) \cup E(H_2)$. For a set of edges $F \subseteq E(D)$, we sometimes write $V(F)$ to denote the set of vertices incident to some edge in F .

If x is a vertex of a digraph D , then $N_D^+(x)$ denotes the *out-neighbourhood* of x , that is, the set of all those vertices y for which $xy \in E(D)$. Similarly, $N_D^-(x)$ denotes the *in-neighbourhood* of x , that is, the set of all those vertices y for which $yx \in E(D)$. For $S \subseteq V(D)$, we write $N_D^+(x, S)$ for all those vertices $y \in S$ such that $xy \in E(D)$ and correspondingly for $N_D^-(x, S)$. We write $d_D^+(x) := |N_D^+(x)|$ for the *outdegree* of x and $d_D^-(x) := |N_D^-(x)|$ for its *indegree*. Similarly, we write $d_D^\pm(x, S) := |N_D^\pm(x, S)|$. We denote the *minimum outdegree* of D by $\delta^+(D) := \min\{d_D^+(x) : x \in V(D)\}$ and the *minimum indegree* $\delta^-(D) := \min\{d_D^-(x) : x \in V(D)\}$. The *minimum semi-degree* of D is $\delta^0(D) := \min\{\delta^+(D), \delta^-(D)\}$ and the *minimum degree* is $\delta(D) := \min\{d^+(x) + d^-(x) : x \in V(D)\}$. We use $\Delta^\pm(D)$, $\Delta^0(D)$ and $\Delta(D)$ for the corresponding maximum degrees.

Whenever $X, Y \subseteq V(D)$ are disjoint, we write $E_D(X)$ for the set of edges of D having both endvertices in X , and $E_D(X, Y)$ for the set of edges of D that start in X and end in Y .

Unless stated otherwise, when we refer to paths and cycles in digraphs, we mean directed paths and cycles, that is, the edges on these paths and cycles are oriented consistently. We write $P = x_1x_2 \cdots x_t$ to indicate that P is a path with edges $x_1x_2, x_2x_3, \dots, x_{t-1}x_t$, where x_1, \dots, x_t are distinct vertices. We occasionally denote such a path P by x_1Px_t to indicate that it starts at x_1 and ends at x_t . For two paths $P = a \cdots b$ and $Q = b \cdots c$, we write $aPbQc$ for the concatenation of the paths P and Q and this notation generalises to cycles in the obvious ways. In particular, for a cycle C and vertices a, b on the cycle, aCb denotes the paths from a to b along the cycle. We often use calligraphic letters, for example, \mathcal{P} for a set of paths $\mathcal{P} = \{P_1, \dots, P_r\}$. In that case $\cup \mathcal{P}$ refers to the digraph that is the union of the paths and $V(\mathcal{P})$ and $E(\mathcal{P})$ refer to the vertex and edge set of the union.

For a set X and $U \subseteq X$, we will write $I_U : X \rightarrow \{0, 1\}$ for the indicator function of U .

For $x, y \in (0, 1]$, we often use the notation $x \ll y$ to mean that x is sufficiently small as a function of y , that is, $x \leq f(y)$ for some implicitly given non-decreasing function $f : (0, 1] \rightarrow (0, 1]$.

Throughout, we omit floors and ceilings and treat large numbers as integers whenever this does not affect the argument.

2.2. Basic graph theory

We will very occasionally work with undirected graphs for which we use standard notation similar to that used for directed graphs (see, for example, [5]).

THEOREM 2.1 (variant of Hall’s Theorem). *Suppose G is a bipartite graph with vertex classes A and B and $k \in \mathbb{N}$. If $k|N_G(X)| \geq |X|$ for every $X \subseteq A$, then each $a \in A$ can be matched with some $b \in B$ such that each $b \in B$ is matched with at most k elements of A , that is, there exists a subgraph $G' \subseteq G$ in which every vertex in A has degree 1 and every vertex in B has degree at most k .*

COROLLARY 2.2. *Suppose G is a bipartite graph with vertex classes A and B both of size n and suppose $\delta(G) \geq n/2$. Then G has a perfect matching.*

For a directed graph D and $A, B \subseteq V(D)$, an *A, B -path* of D is a path of D that starts in A and ends in B . An *A, B -separator* of D is a vertex subset $S \subseteq V(D)$ such that there are no A, B -paths in $D - S$.

THEOREM 2.3 (Menger’s Theorem). *Suppose D is a directed graph and $A, B \subseteq V(D)$. If the smallest A, B -separator in D has size t , then there exist t internally vertex-disjoint A, B -paths in D .*

2.3. *Excess and partial decompositions*

We recall definitions from the introduction. Let D be a directed graph. For a vertex $v \in V(D)$, recall that $\text{ex}_D(v) := d_D^+(v) - d_D^-(v)$. We define $\text{ex}_D^+(v) := \max\{0, \text{ex}_D(v)\}$ and $\text{ex}_D^-(v) := \max\{0, -\text{ex}_D(v)\}$. Let

$$\text{ex}(D) := \frac{1}{2} \sum_{v \in V(D)} |\text{ex}_D(v)| = \sum_{v \in V(D)} \text{ex}_D^+(v) = \sum_{v \in V(D)} \text{ex}_D^-(v).$$

For $* \in \{+, -\}$, let $U^*(D) := \{v \in V(D) : \text{ex}_D^*(v) > 0\}$ and let $U^0(D) := \{v \in V(D) : \text{ex}_D(v) = 0\}$.

We state the following very simple observation, so that we can refer to it later.

PROPOSITION 2.4. *Suppose D is a directed graph and $H \subseteq D$ is a subdigraph in which $\text{ex}_H^*(v) \leq \text{ex}_D^*(v)$ for all $v \in V(D)$ and $* \in \{+, -\}$ (here we define $\text{ex}_H^*(v) = 0$ for $v \in V(D) \setminus V(H)$). Then $\text{ex}(D) = \text{ex}(H) + \text{ex}(D - H)$.*

Proof. To see this, note that either $\text{ex}_D(v)$, $\text{ex}_H(v)$, and $\text{ex}_{D-H}(v)$ are all at least zero or all at most zero for each $v \in V(D)$. Hence, $\text{ex}_D(v) = \text{ex}_H(v) + \text{ex}_{D-H}(v)$ for all $v \in V(D)$. We sum over all vertices to obtain the result. \square

The following definitions are convenient.

DEFINITION 2.5. A *perfect decomposition* of a digraph D is a set $\mathcal{P} = \{P_1, \dots, P_r\}$ of edge-disjoint paths of D that together cover $E(D)$ with $r = \text{ex}(D)$. (Thus, Conjecture 1.1 states that every even tournament has a perfect decomposition.)

A *partial decomposition* of a digraph D is a set $\mathcal{P} = \{P_1, \dots, P_k\}$ of edge-disjoint paths of D such that for every $v \in V(D)$ at most $\text{ex}_D^+(v)$ of the paths start at v and at most $\text{ex}_D^-(v)$ of the paths end at v .

It is easy to see that any subset of a perfect decomposition of G is a partial decomposition of G . We will need the following straightforward fact about perfect decompositions.

PROPOSITION 2.6. *If D is an acyclic digraph, then it has a perfect decomposition.*

Proof. Iteratively remove paths of maximum length. Note that removing such a path from an acyclic digraph reduces the excess by one (since such a path must begin at a vertex v , where $d^-(v) = 0$ (and hence, $\text{ex}(v) > 0$), and must end at a vertex, where $d^+(v) = 0$ (and hence, $\text{ex}(v) < 0$). So, the proposition holds by induction on $\text{ex}(D)$. \square

3. *Exact Decomposition for tournaments with high excess*

In this section, we prove Theorem 1.6. We start by showing that any Eulerian oriented graph can be decomposed into a small number of cycles. We will also need an extra technical condition on this cycle decomposition. We use the following result of Huang, Ma, Shapira, Sudakov, and Yuster [6, Proposition 1.5].

LEMMA 3.1. *Every Eulerian digraph D with n vertices and m edges has a cycle of length $1 + \max(m^2/24n^3, \lfloor \sqrt{m/n} \rfloor)$.*

LEMMA 3.2. *Let $n \in \mathbb{N}$. Let D be an Eulerian oriented graph with n vertices. Then we can decompose D into $t \leq 50n^{4/3} \log n$ cycles C_1, \dots, C_t and for each cycle C_i we can find distinct representatives $x_1^i, x_2^i, \dots, x_{r_i}^i \in V(C_i)$ (indexed in order) with the following properties.*

- (i) *Every cycle has at least two representatives, that is, $r_i \geq 2$ for all i .*
- (ii) *The interval between consecutive representative vertices on a cycle $x_j^i C_i x_{j+1}^i$ has length at most $n^{2/3}$.*
- (iii) *Every vertex $v \in V$ occurs as a representative at most $24n^{2/3} \log^{1/2} n$ times.*

Proof. We first show that D can be decomposed into at most $50n^{4/3} \log n$ cycles. Note that D has $m < n^2/2$ edges. We iteratively remove the longest cycle and let m_t be the number of edges remaining at step t . From Lemma 3.1, we have that $m_{t+1} \leq m_t - g(m_t)$, where

$$g(r) = \max \left\{ r^2/24n^3, \lfloor \sqrt{r/n} \rfloor \right\} > \frac{r}{24n^{4/3}}.$$

To see the inequality note that if $r \geq n^{5/3}$, then $r^2/24n^3 \geq r/24n^{4/3}$, and if $r < n^{5/3}$, then $\sqrt{r/n} > r/n^{1/2+5/6} = r/n^{4/3}$. Thus, we see that

$$m_{t+1} < m_t - \frac{m_t}{24n^{4/3}} = m_t \left(1 - \frac{1}{24n^{4/3}} \right) \leq m_t \exp \left(-\frac{1}{24n^{4/3}} \right).$$

Hence, $m_t < \exp(-t/24n^{4/3})n^2$ from which we see that $m_t < 1$ after at most $50n^{4/3} \log n$ steps, giving at most as many cycles in the greedy decomposition of D .

Next, we show how to obtain the representatives. Assume we have a decomposition of D into a minimum number of cycles C_1, \dots, C_t , where we know $t \leq 50n^{4/3} \log n$.

First we treat the long cycles. Assume without loss of generality that C_1, \dots, C_k are the cycles in our decomposition of length larger than $n^{2/3}$. Divide each such cycle C_i into intervals $I_1^i, \dots, I_{r_i}^i$ each of length between $n^{2/3}/4$ and $n^{2/3}/2$ with r_i minimal. Note that $r_i \leq 4|E(C_i)|n^{-2/3}$ for all $i \in [k]$. Thus, in total we have at most $\sum_{i \in [k]} 4|E(C_i)|n^{-2/3} \leq 4|E(D)|n^{-2/3} \leq 2n^{4/3}$ intervals each of length at least $n^{2/3}/4$. Therefore, we can greedily pick $x_j^i \in I_j^i$ such that no vertex in $V(D)$ appears as a representative more than $8n^{2/3}$ times.

Consider the remaining (short) cycles C_{k+1}, \dots, C_t , for which we need only find two representatives each. Let $\mathcal{C} = \{C_{k+1}, \dots, C_t\}$. First, we will find one representative in each cycle of \mathcal{C} such that no vertex is chosen more than $8n^{2/3} \log^{1/2} n$ times. Let H be the bipartite graph with vertex partitions \mathcal{C} and $V(D)$, where for $C \in \mathcal{C}$ and $v \in V(D)$ are joined if and only if $v \in V(C)$. We now apply (a version of) Hall's theorem (Theorem 2.1) to find one representative in each C such that no vertex is chosen more than $8n^{2/3} \log^{1/2} n$ times. If such a collection of representatives does not exist, then Theorem 2.1 implies that there exists a subset \mathcal{C}' of \mathcal{C} such that $8n^{2/3} \log^{1/2} n |N_H(\mathcal{C}')| < |\mathcal{C}'|$. On the other hand, we have

$$|N_H(\mathcal{C}')| = |V(\bigcup \mathcal{C}')| \geq |E(\bigcup \mathcal{C}')|^{1/2} \geq |\mathcal{C}'|^{1/2}.$$

This implies that $t \geq |\mathcal{C}'| > 64n^{4/3} \log n$, a contradiction. Thus, we have found one representative $x_1^i \in V(C_i)$ for each $k+1 \leq i \leq t$ such that each vertex $v \in V$ occurs as a representative at most $8n^{2/3} \log^{1/2} n$ times. Next, let $P_i := C_i \setminus x_1^i$ for each $k+1 \leq i \leq t$. Note that $|E(P_i)| \geq 1$. By a similar argument as above, we can find one representative $x_2^i \in V(P_i)$ for each $k+1 \leq i \leq t$ such that each vertex $v \in V$ occurs as a representative at most $8n^{2/3} \log^{1/2} n$ times. In summary, we have found two distinct representatives for each $C \in \mathcal{C}$ such that each $v \in V$ occurs as a representative at most $16n^{2/3} \log^{1/2} n$ times.

Now, combining the representatives of the long cycles and the short cycles, we see that each vertex is represented at most $8n^{2/3} + 16n^{2/3} \log^{1/2} n \leq 24n^{2/3} \log^{1/2} n$ times. \square

For the remainder of the section, assume $T = (V, E)$ is a tournament with $\text{ex}(T) \geq 8n^{17/9+\gamma}$, where $1/n \ll \gamma$. In the next two lemmas, we will construct paths in T that will form a partial decomposition of T when combined in the right way. Moreover, it will turn out that these paths can also be used to ‘absorb’ cycles; this is the crucial idea of the proof of Theorem 1.6.

For any digraph D , any $s \in \mathbb{R}$, and $* \in \{+, -\}$, we define $W_s^*(D) := \{v \in V(D) : \text{ex}_D^*(v) \geq s\}$.

LEMMA 3.3. *Let $n \in \mathbb{N}$ and $1/n \ll \gamma$. Suppose that $T = (V, E)$ is a tournament on n vertices with $\text{ex}(T) \geq 8n^{17/9+\gamma}$. Set $s = n^{8/9+\gamma}$. Let $H \subseteq T$ with $\Delta(H) \leq s$ and $S \subseteq V$ with $|S| \leq s$. For any $v \in V \setminus S$, there exist $n^{2/3+\gamma}$ paths in $T - H - S$ that start in $W_s^+(T)$, end at v , have length at most $4n^{1/9}$, and are vertex-disjoint except at their end point v .*

In the statement above, a path could be a single vertex v , and moreover this single vertex path could occur many times. Such a collection of paths should be thought of as being ‘vertex disjoint with each other except at their end point v ’.

Note that, by symmetry, the same result as above holds if we wish to find paths from v to $W_s^-(T)$.

Proof. We write W^+ for $W_s^+(T)$ and note that $\text{ex}(T) \leq |W^+|n + ns$, so that $|W^+| \geq 7n^{8/9+\gamma}$. If $v \in W^+$, then we are done (by the remark above), so assume not. Write $T' := T - H - S$. Let $A^+ := W^+ \setminus S$. Suppose that all (A^+, v) -separators in T' have size at least s . Thus, by Menger’s Theorem we can find at least s paths in T' that start in A^+ , end at v and are vertex disjoint except for their common end point v . If we pick the shortest $n^{2/3+\gamma}$ of these paths, they all have length at most $4n^{1/9}$ since otherwise we have at least $s - n^{2/3+\gamma} \geq \frac{1}{2}n^{8/9}$ paths of length at least $4n^{1/9}$ that are vertex-disjoint except for one common vertex; such paths cover at least $\frac{1}{2}n^{8/9} \cdot (4n^{1/9} - 1) > n$ vertices, a contradiction. Therefore, to prove the lemma, it suffices to show that all (A^+, v) -separators X in T' satisfy $|X| \geq n^{8/9+\gamma}$.

Let X be a (A^+, v) -separator in T' and let $\hat{T} := T' - X = T - H - (S \cup X)$. Define

$$B = \{x \in V(\hat{T}) : \exists \text{ a path from } x \text{ to } v \text{ in } \hat{T}\}.$$

Then $A^+ \cap B = \emptyset$ (since otherwise X is not a (A^+, v) -separator) and so $W^+ \cap B = \emptyset$. Furthermore, by the definition of B there are no directed edges in \hat{T} from $\bar{B} := V(\hat{T}) \setminus B$ to B . Using this and the fact that every vertex in a tournament is adjacent (in some direction) to every other vertex, we have for all $x \in B$ that

$$\begin{aligned} |N_T^+(x) \setminus B| &\geq V(T) - |B| - |X| - |S| - \Delta(H) \\ &\geq |W^+| - |X| - |S| - \Delta(H), \end{aligned}$$

and

$$|N_T^-(x) \setminus B| \leq |X| + |S| + \Delta(H).$$

Pick a vertex $x^* \in B$ with $\text{ex}_{T[B]}(x^*) \geq 0$ (note that every directed graph has a vertex with non-negative excess). Then we have

$$\begin{aligned} \text{ex}_T(x^*) &\geq \text{ex}_{T[B]}(x^*) + |N_T^+(x^*) \setminus B| - |N_T^-(x^*) \setminus B| \\ &\geq 0 + |W^+| - 2|X| - 2|S| - 2\Delta(H) \geq |W^+| - 4s - 2|X|. \end{aligned}$$

We know that $\text{ex}_T(x^*) \leq s$ (otherwise $x^* \in W^+$, a contradiction) and that $|W^+| \geq 7n^{8/9+\gamma} = 7s$. Hence, $|X| \geq s = n^{8/9+\gamma}$, as required. \square

By inductively applying the previous lemma, we obtain the following.

LEMMA 3.4. *Let $n \in \mathbb{N}$ and $1/n \ll \gamma$. Suppose that $T = (V, E)$ is a tournament on n vertices with $\text{ex}(T) \geq 8n^{17/9+\gamma}$. Let $\ell := n^{2/3+\gamma}$ and $m := 4n^{1/9}$ and $s := n^{8/9+\gamma}$. Then we can find edge-disjoint paths P_j^v, Q_j^v , where $v \in V, j = 1, \dots, \ell$ with the following properties.*

- (i) P_j^v is a path of length at most m from $W_s^+(T)$ to v and Q_j^v is a path of length at most m from v to $W_s^-(T)$.
- (ii) For each fixed $v \in V$, the paths P_1^v, \dots, P_ℓ^v are vertex-disjoint except that they all meet at v and the paths Q_1^v, \dots, Q_ℓ^v are vertex-disjoint except that they all meet at v .
- (iii) $\Delta(\bigcup_{v,j}(P_j^v \cup Q_j^v)) < n^{8/9+\gamma} = s$.

Proof. Fix an ordering v_1, \dots, v_n of the vertices of T and inductively construct the desired paths as follows. Suppose at the k th step, we have constructed the $P_j^{v_i}$ and $Q_j^{v_i}$ for all $i \leq k - 1$ and all $j \leq \ell$ satisfying the first two conditions of the lemma. Furthermore, we assume that the oriented graph H_{k-1} on V , which is union of the paths constructed so far, satisfies

$$d_{H_{k-1}}(v_i) \leq 2\ell + s/2 \quad \forall 1 \leq i \leq k - 1 \tag{3.1}$$

$$\text{and } d_{H_{k-1}}(v_i) \leq s/2 \quad \forall k \leq i \leq n. \tag{3.2}$$

By our choice of parameters, we have

$$\Delta(H_{k-1}) \leq 2\ell + s/2 < s. \tag{3.3}$$

Let S^* be the set of vertices $v \in V$ such that $d_{H_{k-1}}(v) \geq s/4$. Note that $\frac{1}{4}s|S^*| \leq 2|E(H_{k-1})| \leq 4nm\ell = 16n^{16/9+\gamma}$, so $|S^*| \leq 64n^{8/9} \leq s$.

Now, applying Lemma 3.3 (where (H_{k-1}, S^*) play the role of (H, S)), we obtain vertex-disjoint (except at v_k) paths $P_j^{v_k}$ for all $j \leq \ell$ from $W_s^+(T)$ to v_k each of length at most m . Applying Lemma 3.3 again (where $(H_{k-1} \cup \bigcup_j E(P_j^{v_k}), S^*)$ play the roles of (H, S) and noting $\Delta(\bigcup_j E(P_j^{v_k})) \leq \ell$), we obtain vertex-disjoint (except at v_k) paths $Q_j^{v_k}$ for all $j \leq \ell$ from v to $W_s^-(T)$, each of length at most m . Note that all the new paths are edge-disjoint from each other and from the old ones and satisfy properties (i) and (ii) of the lemma.

Letting H_k be the union of all the paths constructed so far, note that compared to H_{k-1} , the degree of v_k goes up by at most 2ℓ and the degree of every vertex $v \in V \setminus (S^* \cup v_k)$ goes up by at most 4. Thus, (3.1) and (3.2) hold. At the n th step, we are able to construct all the paths satisfying properties (i) and (ii), and property (iii) also holds by (3.3). \square

We now prove the following theorem which immediately implies Theorem 1.6 by taking $\varepsilon = 1/18$.

THEOREM 3.5. *Let $n \in \mathbb{N}$ and $1/n \ll \gamma$. Suppose that $T = (V, E)$ is a tournament on n vertices with $\text{ex}(T) \geq 8n^{17/9+\gamma}$. Then T has a perfect decomposition.*

Proof. Let $\ell := n^{2/3+\gamma}$ and $m := 4n^{1/9}$ and $s := n^{8/9+\gamma}$. Apply Lemma 3.4 to T , so that we obtain edge-disjoint paths P_j^v, Q_j^v , where $v \in V$ and $j = 1, \dots, \ell$ with the following properties.

- (i) P_j^v is a path of length at most m from $W_s^+(T)$ to v and Q_j^v is a path of length at most m from v to $W_s^-(T)$.
- (ii) For each fixed $v \in V$, the paths P_1^v, \dots, P_ℓ^v are vertex-disjoint except that they all meet at v and the paths Q_1^v, \dots, Q_ℓ^v are vertex-disjoint except that they all meet at v .
- (iii) $\Delta(\bigcup_{v,j}(P_j^v \cup Q_j^v)) < s$.

Call a path of the form P_j^v a *v-in-path* and a path of the form Q_j^v a *v-out-path*. Write H for the graph that is the union of these paths and let $T' = T - H$. For each $v \in V$ and $j \leq \ell$, each walk $P_j^v \cup Q_j^v$ starts in $W_s^+(T)$ and ends in $W_s^-(T)$ and no vertex occurs as a start or end point more than s times. Therefore, we have that $\text{ex}_H^\pm(v) \leq s \leq \text{ex}_T^\pm(v)$ for all $v \in W_s^\pm(T)$ and $\text{ex}_H(v) = 0$ for all other $v \in V \setminus (W_s^+(T) \cup W_s^-(T))$. This means in particular that $\text{ex}(T) = \text{ex}(H) + \text{ex}(T')$ (by Proposition 2.4). (In fact, $\text{ex}(H) = \ell n = n^{5/3+\gamma}$.) Let T_E be a maximal Eulerian subgraph of T' and let $T_R = T' - T_E$, where T_R is necessarily acyclic. Thus, we have that $T = H \cup T_R \cup T_E$ and $\text{ex}(T) = \text{ex}(H) + \text{ex}(T_R) + \text{ex}(T_E)$ (and of course $\text{ex}(T_E) = 0$).

Finally, we show how to decompose $T_E \cup H$ into $\text{ex}(H)$ paths. Apply Lemma 3.2 to T_E . Thus, we can decompose T_E into $t \leq 50n^{4/3} \log n$ cycles C_1, \dots, C_t and for each cycle C_i we can find distinct representatives $x_1^i, x_2^i, \dots, x_{r_i}^i \in V(C_i)$ (indexed in order) with the following properties.

- (i') Every cycle has at least two representatives, that is, $r_i \geq 2$ for all i .
- (ii') The interval between consecutive vertices on a cycle $x_j^i C_i x_{j+1}^i$ has length at most $n^{2/3}$.
- (iii') Every vertex $v \in V$ occurs as a representative at most $24n^{2/3} \log^{1/2} n$ times.

Write C_j^i for the interval $x_j^i C_i x_{j+1}^i$. By (i), (ii), (ii'), and (iii'), for each $i \leq t$ and $j \leq r_i$, we can greedily find distinct $P_j^{x_j^i}$ such that each $P_j^{x_j^i} C_j^i$ is a path from $W_s^+(T)$ to x_{j+1}^i . (Given C_j^i , since the paths $P_1^{x_j^i}, \dots, P_\ell^{x_j^i}$ are vertex disjoint (except at x_j^i), at least $\ell - |C_j^i|$ of these paths avoid C_j^i and since we never use more than $24n^{2/3} \log^{1/2} n$ of these paths, there is always one available.) Hence, we have shown that $\bigcup_{v,j} P_j^v \cup T_E$ can be edge-decomposed into ℓn paths $P_1, \dots, P_{\ell n}$ each of length at most $n^{2/3} + m$. Note crucially that each vertex v is an end point of exactly ℓ paths and at least $\ell - 24n^{2/3} \log^{1/2} n$ of such paths belong to $\{P_j^v : j \leq \ell\}$.

We now extend $P_1, \dots, P_{\ell n}$ using the paths $\{Q_j^v : v \in V, j \leq \ell\}$ as follows. Consider any $v \in V$. Let \mathcal{P}_v be the set of P_i with end point v and let $\mathcal{Q}_v = \{Q_j^v : j \leq \ell\}$. Clearly, $|\mathcal{P}_v| = \ell = |\mathcal{Q}_v|$. Let \mathcal{P}'_v (and \mathcal{Q}'_v) be the largest set of vertex-disjoint (except at v) paths of \mathcal{P}_v (and \mathcal{Q}_v , respectively). Thus, $|\mathcal{P}'_v| \geq \ell - 24n^{2/3} \log^{1/2} n$ and $|\mathcal{Q}'_v| = \ell$. Let B be the auxiliary bipartite graph with vertex partition \mathcal{P}_v and \mathcal{Q}_v , where $P \in \mathcal{P}_v$ is joined to $Q \in \mathcal{Q}_v$ if and only if $V(P) \cap V(Q) = \{v\}$. For each $Q \in \mathcal{Q}_v$, $|N_B(Q)| \geq |\mathcal{P}'_v| - |V(Q)| \geq \ell - (24n^{2/3} \log^{1/2} n) - m \geq \ell/2$. Similarly, we have $|N_B(P)| \geq \ell/2$ for each $P \in \mathcal{P}_v$. By Corollary 2.2, B has a perfect matching, which implies $\bigcup \mathcal{P}_v \cup \bigcup \mathcal{Q}_v$ can be decomposed into ℓ paths. Therefore, $T_E \cup H = \bigcup_{v \in V} (\bigcup \mathcal{P}_v \cup \bigcup \mathcal{Q}_v)$ can be decomposed into $\ell n = \text{ex}(H)$ paths.

Thus, we can now write $T = H \cup T' = (H \cup T_E) \cup T_R$, where $\text{ex}(T) = \text{ex}(H) + \text{ex}(T') = \text{ex}(H) + \text{ex}(T_R) = \ell n + \text{ex}(T_R)$ and where $H \cup T_E$ can be decomposed into ℓn paths and T_R can be decomposed into $\text{ex}(T_R)$ paths (by Proposition 2.6). Hence, T can be decomposed into $\text{ex}(T)$ paths. □

4. Further preliminaries and overview

In this section, we provide further preliminaries used in Sections 5 and 6 as well as an overview of the proof of Theorem 1.5.

4.1. Partial decompositions

We will use the following easy facts about partial decompositions repeatedly. The proofs are straightforward, but we give them for completeness.

PROPOSITION 4.1. *Let D be a directed graph and let $\mathcal{P} = \{P_1, \dots, P_k\}$ be a partial decomposition of D where P_i is a path from x_i to y_i . Then the following hold.*

- (a) *Any $\mathcal{Q} \subseteq \mathcal{P}$ is a partial decomposition of D and a partial decomposition of $D - E(\mathcal{P} \setminus \mathcal{Q})$.*
- (b) *If \mathcal{Q} is a partial decomposition of $D - E(\mathcal{P})$, then $\mathcal{P} \cup \mathcal{Q}$ is a partial decomposition of D (and so is \mathcal{Q}).*
- (c) *If π is a permutation of $[k]$ and $\mathcal{Q} = \{Q_1, \dots, Q_r\}$ is a set of edge-disjoint paths with $r \leq k$ and Q_i is a path from x_i to $y_{\pi(i)}$, then \mathcal{Q} is a partial decomposition of D .*
- (d) *If $D' \subseteq D$ is an Eulerian subdigraph of D and \mathcal{Q} is a partial decomposition of $D - D'$, then \mathcal{Q} is a partial decomposition of D .*

Proof. For any collection of paths $\mathcal{A} = \{A_1, \dots, A_t\}$ where A_i is a path in a digraph D and $x \in V(D)$, write $p_{\mathcal{A}}^+(x)$ for the number of paths in \mathcal{A} that start at x and $p_{\mathcal{A}}^-(x)$ for the number of paths in \mathcal{A} that end at x .

- (a) The fact that \mathcal{Q} (and $\mathcal{P} \setminus \mathcal{Q}$) is a partial decomposition of D is immediate. For the second part note that for any $x \in V(D)$, if $\text{ex}_D(x) \geq 0$, then

$$\begin{aligned} \text{ex}_{D-E(\mathcal{P} \setminus \mathcal{Q})}^+(x) &= \text{ex}_D^+(x) - p_{\mathcal{P} \setminus \mathcal{Q}}^+(x) + p_{\mathcal{P} \setminus \mathcal{Q}}^-(x) \\ &\geq p_{\mathcal{P}}^+(x) - p_{\mathcal{P} \setminus \mathcal{Q}}^+(x) + 0 = p_{\mathcal{Q}}^+(x), \end{aligned}$$

where the inequality holds since \mathcal{P} is a partial decomposition of D . A similar statement holds if $\text{ex}_D(x) \leq 0$.

- (b) Note that for any $x \in V(D)$, if $\text{ex}_D(x) \geq 0$, then

$$\text{ex}_D(x) - p_{\mathcal{P}}^+(x) + p_{\mathcal{P}}^-(x) = \text{ex}_{D-E(\mathcal{P})}^+(x) \geq p_{\mathcal{Q}}^+(x)$$

and a similar statement holds if $\text{ex}_D(x) \leq 0$. Rearranging gives $\text{ex}_D(x) \geq p_{\mathcal{P}}^+(x) + p_{\mathcal{Q}}^+(x) = p_{\mathcal{P} \cup \mathcal{Q}}^+(x)$.

- (c) Here we note that $p_{\mathcal{P}}^+(x) \geq p_{\mathcal{Q}}^+(x)$ and $p_{\mathcal{P}}^-(x) \geq p_{\mathcal{Q}}^-(x)$ for all $x \in V(D)$.

- (d) Here we note that $\text{ex}_D(x) = \text{ex}_{D-E(D')}$ for all $x \in V(D)$. □

PROPOSITION 4.2. *Let D be a directed graph and suppose there is a partition of $V(D)$ into sets A^+, A^-, R such that $E_D(R, A^+) = E_D(A^-, R) = E(D[A^+ \cup A^-]) = \emptyset$. Then the following holds.*

- (a) *If $\mathcal{P} = \{P_1, \dots, P_r\}$ is a partial decomposition of $D[R]$, then there is a partial decomposition $\mathcal{P}' = \{P'_1, \dots, P'_r\}$ of D such that $V(P'_i) \cap R = V(P_i)$ for all $i = 1, \dots, r$.*
- (b) *If there is a perfect decomposition of $D[R]$, then there is a perfect decomposition of D .*
- (c) *If in addition we assume $\text{ex}_D(v) \geq 0$ for every $v \in N_D^+(A^+)$ and $\text{ex}_D(v) \leq 0$ for every $v \in N_D^-(A^-)$ and $N_D^+(A^+) \cap N_D^-(A^-) = \emptyset$, then $\text{ex}(D[R]) = \text{ex}(D)$.*

Proof. (a) This is easily proved by induction on the number of paths; we give the details for completeness. By induction, we will find paths P'_1, \dots, P'_r such that each path $\{P'_i\}$ is a partial decomposition of $D_i := D - (P'_1 \cup \dots \cup P'_{i-1})$ for $i = 1, \dots, r$ and $V(P_i) = V(P'_i) \cap R$. By r applications of Proposition 4.1(b), $\{P'_1, \dots, P'_r\}$ is a partial decomposition of D with the desired properties.

Suppose we have found the paths P'_1, \dots, P'_{k-1} as described above. Then $\{P_k\}$ is a partial decomposition of $D[R] - (P'_1 \cup \dots \cup P'_{k-1}) = D_k[R]$. Write $P_k = xP_k y$. If there is some edge $a^+x \in E(D_k)$ with $a^+ \in A^+$, then append it to P_k and if there is some edge $ya^- \in E(D_k)$ with

$a^- \in A^-$, then append it to P_k and write P'_k for the resulting path. Let $P'_k = x'P_ky'$; we show that $\text{ex}_{D_k}(x') > 0 > \text{ex}_{D_k}(y')$ proving that $\{P'_k\}$ is a partial decomposition of D_k .

By symmetry, it is sufficient to show $\text{ex}_{D_k}(x') > 0$. If $x' = a^+ \in A^+$, then this is certainly the case. If $x' = x$, then

$$\text{ex}_{D_k}(x) \geq \text{ex}_{D_k[R]}(x) > 0.$$

The first inequality holds because there is no edge in D_k from A^+ to x (nor from A^- to x from the statement of the lemma). The second inequality holds because P_k starts at x and $\{P_k\}$ is a partial decomposition of $D_k[R]$.

(b) From (a) we can extend our perfect decomposition of $D[R]$ to a partial decomposition \mathcal{Q}_1 of D that uses every edge of $D[R]$. The remaining digraph $D - E(\mathcal{Q}_1) \subseteq D - E(D[R])$ is acyclic so has a perfect decomposition \mathcal{Q}_2 by Proposition 2.6. Therefore, $\mathcal{Q}_1 \cup \mathcal{Q}_2$ is a perfect decomposition of D by Proposition 4.1(b).

(c) This is proved by induction on the number of edges between $A^+ \cup A^-$ and R . If D has no edges between $A^+ \cup A^-$ and R , then we are done. For any edge $e = a^+r$ with $a^+ \in A^+$ and $r \in R$, $\text{ex}_{D-e}(r) > \text{ex}_D(r) \geq 0$. Furthermore, $r \notin N_{D-e}^-(A^-)$ because $N_D^+(A^+) \cap N_D^-(A^-) = \emptyset$. It is easy to check that the conditions in (c) are satisfied for $D - e$ so we can assume by induction that $\text{ex}(D - e) = \text{ex}(D[R])$. Also, we see that adding the edge e back to $D - e$ reduces $\text{ex}(r)$ by 1 and increases $\text{ex}(a^+)$ by 1, so that $\text{ex}(D) = \text{ex}(D - e) = \text{ex}(D[R])$. The case when $e = ra^-$ for some $r \in R$ and some $a^- \in A^-$ holds similarly. \square

4.2. Robust expanders

The notation of robust expanders was first introduced by Kühn, Osthus, and Treglown [11]. Here we introduce the basic notions of robust expansion and their consequences, which we will use in Sections 5 and 6. Most of this can be found in [9, 10]

We give the definition of robust expander here for completeness. We will not use the definition directly, but only use some of the consequences given below.

DEFINITION 4.3. An n -vertex digraph D is a robust (ν, τ) -outexpander if for every $S \subseteq V(D)$ with $\tau n \leq |S| \leq (1 - \tau)n$ there is some set $T \subseteq V(D)$ with $|T| \geq |S| + \nu n$ such that every vertex in T has at least νn in-neighbours in $|S|$.

It turns out that sufficiently dense oriented graphs are robust expanders.

LEMMA 4.4 [9, Lemma 13.1]. *Let $0 < 1/n \ll \nu \ll \tau \ll \varepsilon$. Suppose that D is an oriented graph on n vertices with $\delta^0(D) \geq (3/8 + \varepsilon)n$. Then D is a robust (ν, τ) -outexpander.*

The notion of robust expansion was developed to help solve problems on Hamilton cycles. Here are two of the main results.

THEOREM 4.5 [8, Corollary 6.9]. *Let $0 < 1/n \ll \nu \ll \tau \ll \delta$. Suppose that D is a robust (ν, τ) -outexpander on n vertices with $\delta^0(D) \geq \delta n$. Let $a, b \in V(D)$. Then D contains a Hamilton path from a to b .*

THEOREM 4.6 [9]. *Let $0 < 1/n \ll \nu \ll \tau \ll \delta$. Suppose that D is an r -regular oriented graph with $r \geq \delta n$ and a robust (ν, τ) -outexpander. Then $E(D)$ can be decomposed into r edge-disjoint Hamilton cycles.*

An immediate consequence of the above is the following path decomposition result, which we use right at the end of the paper.

THEOREM 4.7. *Let $0 < 1/n \ll 1$ and let D be an oriented graph with a vertex partition $V(D) = X^+ \cup X^- \cup X^0$ with $|X^+| = |X^-| = d \geq 3n/7$ such that*

$$\text{ex}_D(v) = \begin{cases} 1 & \text{if } v \in X^+; \\ 0 & \text{if } v \in X^0; \\ -1 & \text{if } v \in X^-; \end{cases} \quad \text{and} \quad d_D(v) = \begin{cases} 2d - 1 & \text{if } v \in X^+; \\ 2d & \text{if } v \in X^0; \\ 2d - 1 & \text{if } v \in X^-. \end{cases}$$

Then D has a perfect decomposition.

Proof. Fix ν, τ such that $1/n \ll \nu \ll \tau \ll 1$. We form D' by adding a vertex y such that $N_{D'}^+(y) = X^+$ and $N_{D'}^-(y) = X^-$. Then D' is a regular oriented graph with in- and outdegree $d > 3/7n$ and so is a robust (ν, τ) -outexpander by Lemma 4.4. Thus, it has an edge decomposition into Hamilton cycles H_1, \dots, H_d by Theorem 4.6. Taking P_i to be the path $H_i - y$, $\mathcal{P} = \{P_1, \dots, P_d\}$ gives a perfect decomposition of D . \square

Robust expanders are highly connected as one would expect and so we can find (many) short paths between any pair of vertices. This is made precise in the following three lemmas.

LEMMA 4.8 (see, for example, [12, Lemma 9]). *Let $n \in \mathbb{N}$ and $0 < 1/n \ll \nu \ll \tau \ll \delta \leq 1$. Suppose that D is a robust (ν, τ) -outexpander on n vertices with $\delta^0(D) \geq \delta n$. Then, given any distinct vertices $x, y \in V(D)$, there exists a path P in D from x to y such that $|V(P)| \leq \nu^{-1}$.*

The following lemma and its corollary will be used many times in our proof.

LEMMA 4.9. *Let $n \in \mathbb{N}$ and $0 < 1/n \ll \gamma \ll 1$. Suppose that D is an oriented graph on n vertices with $\delta^0(D) \geq 3n/7$. Let H_1, \dots, H_m be directed multigraphs on $V(D)$ with $\Delta(H_i) \leq 2$, $|E(H_i)| \leq \gamma n$ and $m \leq \gamma n$. Let $S_1, \dots, S_m \subseteq V(D)$ with $|S_i| \leq n/25$ and $S_i \cap V(E(H_i)) = \emptyset$. Then there exists a set of edge-disjoint paths $\mathcal{P} = \{P_{i,e} : i \in [m] \text{ and } e \in E(H_i)\}$ in D such that:*

- (i) $P_{i,e}$ has the same starting and ending points as e ;
- (ii) the paths in $\mathcal{P}_i := \{P_{i,e} : e \in E(H_i)\}$ are internally vertex-disjoint;
- (iii) $V(\cup \mathcal{P}_i) = V(D) \setminus S_i$;
- (iv) $\Delta(\cup \mathcal{P}) \leq 2m$.

Proof. We proceed by induction on m , the number of multigraphs. Suppose that we have already found $\mathcal{P}' := \{P_{i,e} : i \in [m-1] \text{ and } e \in E(H_i)\}$ with the desired properties.

Let ν, τ, ε be such that $\gamma \ll \nu \ll \tau \ll \varepsilon \ll 1$. Pick an arbitrary ordering e_1, \dots, e_r of the edges in $E(H_m)$. Further, assume that for some $j \in [r]$, we have already constructed paths P_1, \dots, P_{j-1} such that, for each $j' \in [j-1]$,

- (i) $P_{j'}$ has the same starting and ending points as $e_{j'}$ and has length at most ν^{-1} ;
- (ii) $V(E(H_i)), S_i, V(P_1) \setminus V(e_1), \dots, V(P_{j-1}) \setminus V(e_{j-1})$ are disjoint.

We now find P_j as follows. Let $e_j = xy$. Let $D' := D - E(\cup \mathcal{P}') - S_i - (V(P_1 \cup \dots \cup P_{j-1}) \setminus \{x, y\})$. Since $|S_i \cup V(P_1 \cup \dots \cup P_{j-1})| \leq (\frac{1}{25} + \nu^{-1}\gamma)n$, then $|D'| \geq (1 - \frac{1}{25} - \nu^{-1}\gamma)n$ and $\delta^0(D') \geq \delta^0(D) - (\frac{1}{25} + \nu^{-1}\gamma)n \geq (3/8 + \varepsilon)|D'|$. By Lemma 4.4, D' is a robust (ν, τ) -outexpander. If $j < r$, then D' has a path P_j from x to y of length at most ν^{-1} by Lemma 4.8. If $j = r$, then D' has a Hamilton path P_j from x to y by Theorem 4.5. We are done by setting $P_{m,e_j} := P_j$ for all $j \in [r]$. \square

Let H be a directed multigraph on n vertices with $\Delta(H) \leq \gamma n$. Note that H can be decomposed into digraphs H_1, \dots, H_m with $m \leq 2\sqrt{\gamma n}$ and $\Delta(H_i) \leq 1$ and $|E(H_i)| \leq 2\sqrt{\gamma n}$.

(By Vizing’s theorem, H can be partitioned into $\gamma n + 1$ matchings and each matching can then be further split into $\gamma^{-1/2}$ almost equal parts to give us the H_i .) Applying the previous lemma to these H_i (taking the S_i to be empty), we obtain the following corollary.

COROLLARY 4.10. *Let $n \in \mathbb{N}$ and $0 < 1/n \ll \gamma \ll 1$. Suppose that D is an oriented graph on n vertices with $\delta^0(D) \geq 3n/7$. Let H be a directed multigraph on $V(D)$ with $\Delta(H) \leq \gamma n$. Then there exists a set of edge-disjoint paths $\mathcal{P} = \{P_e : e \in E(H)\}$ in D such that:*

- (i) P_e has the same starting and ending points as e ;
- (ii) $\Delta(\cup \mathcal{P}) \leq 4\sqrt{\gamma}n$.

4.3. Overview

In this subsection, we give an overview of the proof of Theorem 1.5 (which is proved in Sections 5 and 6). We wish to show that every even n -vertex tournament T satisfying $\text{ex}(T) > Cn$ and n sufficiently large has a perfect decomposition (that is, is consistent). Let us fix such a tournament T ; we may further assume by Theorem 1.6 that $\text{ex}(T) < n^{2-\varepsilon}$. We will accomplish this in three steps. In each step, we reduce the problem of finding a perfect decomposition of T to the problem of finding a perfect decomposition of a digraph that looks more and more like the digraph described in Theorem 4.7.

Step 1: Remove vertices of high excess. Let $W = \{v \in V(T) : |\text{ex}(v)| > \alpha n\}$ for some suitable α . Note that since $\text{ex}(T)$ is small, W is also small. Let W^\pm be, respectively, the vertices of W with positive / negative excess and let $R = V(T) \setminus W$. We will construct a partial decomposition \mathcal{P}_0 of T with a small number of paths that uses all edges in $E_T(R, W^+) \cup E_T(W^-, R) \cup E_T(W)$ but does not interfere much with $E_T(R)$. Set $D_1 = T - \cup \mathcal{P}_0 - W$. Now, we can apply Proposition 4.2(b) to $T - \cup \mathcal{P}_0$ to conclude that if D_1 has a perfect decomposition, then so does $T - \cup \mathcal{P}_0$ and hence, so does T . Thus, we have reduced the problem of finding a perfect decomposition of T to that of finding one for D_1 , but where D_1 has no vertices of high excess and

$$\text{ex}(D_1) = \text{ex}(T - \cup \mathcal{P}_0) = \text{ex}(T) - |\mathcal{P}_0| \geq C'n.$$

Since there are no vertices of high excess, D_1 is close to regular and so one can apply the methods of robust expansion. This step takes place in Theorem 6.1 and the key tool for finding \mathcal{P}_0 is Lemma 5.1 from Section 5.

Step 2: Equalise the number of vertices of positive and negative excess. Given D_1 from the previous step, it may be the case that almost all vertices of D_1 have say negative excess that is $U^-(D_1)$ is significantly larger than $U^+(D_1)$, where $U^\pm(D)$ denote the set of vertices of positive / negative excess in D .

For some fixed $z \in U^-(D_1)$, consider how we might change the sign of its excess. The idea would be to find $x \in U^+(D_1)$ with $xz \in E(D_1)$ and a partial decomposition \mathcal{Q} that:

- has a path Q^* that starts at x , uses the edge xz but does not end at z ;
- uses all edges incident with x ;
- has exactly $\text{ex}_{D_1}^-(z)$ paths ending at z .

If we can find such a \mathcal{Q} , then consider $D'_1 = D_1 - E(\mathcal{Q} \setminus \{Q^*\}) - x$. We have $\text{ex}_{D'_1}(z) = 1$ and moreover if D'_1 has a perfect decomposition, so does D_1 (the path that starts at z in a perfect decomposition of D'_1 would be extended by the edge xz in D_1).

We refine this idea to switch the sign of the excess for many vertices in $U^-(D_1)$ in Theorem 6.7. We carefully choose a small set of vertices $X \subseteq U^+(D_1)$ and a suitably larger set $Z \subseteq U^-(D_1)$ and a partial decomposition \mathcal{P}_1 of D_1 such that writing $D_2 = D_1 - E(\mathcal{P}_1) - X$,

D_1 has a perfect decomposition if D_2 does, and $U^+(D_2) = U^-(D_2) \cup Z \setminus X$ and $U^-(D_2) = U^-(D_1) \setminus Z$. Again we use Lemma 5.1 from Section 5 as a tool.

Step 3: Control the degrees. In this final step (Theorem 6.12), starting with D_2 we carefully construct a partial decomposition \mathcal{P}_2 of D_2 such that $D_3 = D_2 - E(\mathcal{P}_2)$ is a digraph satisfying the properties of Theorem 4.7. Hence, D_3 has a perfect decomposition, and thus so do D_2 , D_1 , and T .

We make use of the robust expansion properties of D_2 to construct \mathcal{P}_2 ; this is why we need step 1. Also, essentially by definition, the excess of a vertex can never change sign when we remove a partial decomposition from a digraph; this is why we need step 2. Each of steps 1 and 2 will require us to remove a partial decomposition of size linear in n , and this is why we must start with $\text{ex}(T) > Cn$ for a suitably large C .

5. Removing small vertex subsets

In Section 3, we showed how to find a perfect decomposition of n -vertex tournaments T (n even) whenever $\text{ex}(T) > n^{2-\varepsilon}$. For the remaining cases of Theorem 1.5, we will require a preliminary result which we prove in this section. For almost complete oriented graphs D satisfying $Cn \leq \text{ex}(D) \leq n^{2-\varepsilon}$, we show in Lemma 5.1 that for certain choices of small $W \subseteq V(D)$, we can find a partial decomposition \mathcal{P} of D that uses all the edges incident with W going in the ‘wrong’ direction. We will also guarantee that \mathcal{P} uses only a small number of edges from $D - W$ and that $|\mathcal{P}|$ is small. This will be useful later as, in combination with Proposition 4.2, it allows us to remove a small number of problematic vertices from our digraph D at the expense of a small reduction in $\text{ex}(D)$. This is the content of Lemma 5.1 and our goal in this section is to prove it.

LEMMA 5.1. *Let $n \in \mathbb{N}$ and $0 < 1/n \ll \alpha, \beta \ll \gamma \ll 1$ and $0 < 1/n \ll \varepsilon \ll 1$ and $C \geq 32$. Let D be an oriented graph on n vertices such that $\delta(D) \geq (1 - \varepsilon)n$ and $\text{ex}(D) \geq Cn$. Let $W \subseteq V(D)$ of size $|W| \leq \beta n$. Suppose that $|\text{ex}_D(v)| \leq \alpha n$ for all $v \in V(D) \setminus W$. Then there exists a partial decomposition \mathcal{P} of D such that writing $H = \cup \mathcal{P}$ we have:*

- (i) for all $v \in V(D) \setminus W$, $d_H(v) = 2d$ for some $d \leq (18\beta + 4\gamma)n$;
- (ii) $H[W] = D[W]$;
- (iii) for all $w \in W$, if $\pm \text{ex}_D(w) \geq 0$, then $d_{D-H}^\mp(w) = 0$;
- (iv) $\text{ex}(D - H) = \text{ex}(D) - \text{ex}(H) \geq Cn/4$.

Note that (iii) guarantees that for every $w \in W$ with $\text{ex}(w) \geq 0$ (respectively, $\text{ex}(w) \leq 0$), every edge of the form vw (respectively, wv) is in H and we informally refer to such edges as going in the ‘wrong’ direction. The proof of Lemma 5.1 is split into two lemmas, Lemmas 5.4 and 5.7. In Lemma 5.4, we deal with all edges inside W and in Lemma 5.7, we deal with the edges between W and $V(D) \setminus W$ going in the ‘wrong’ direction. The basic idea in each case is as follows. Write F for the set of edges incident with W which we wish to remove from D (and thus to add to H). Each of these edges can be thought of as a path and we start by extending these paths (if necessary), so that their end points lie in $V(D) \setminus W$ to give a set of paths \mathcal{Q} . The reason for doing this is that $D - W$ is a robust expander and so has good connectivity properties; this allows us to connect the large number of paths in \mathcal{Q} into a small number of long paths \mathcal{Q}' (see Corollary 5.3). At the same time, we can ensure the paths in \mathcal{Q}' have suitable start and end points, so that \mathcal{Q}' is a partial decomposition with a small number of paths that contains all edges in F . While this is conceptually quite simple, the process of extending the paths into $V(D) \setminus W$ and choosing appropriate start and end points becomes technical if we wish to ensure that the paths we create do not interfere with each other.

Before we can prove these two lemmas, we will need a technical definition and one preliminary result.

Consider a digraph D and a vertex subset $W \subseteq V(D)$. Let $V = V(D) \setminus W$. Suppose we have two internally vertex-disjoint paths P, P' that both start at some $x \in V(D)$ and end at some different vertex $y \in V(D)$. Now, starting with $P \cup P'$ delete any edges of $P \cup P'$ that occur inside V ; this is essentially what we refer to as a (W, V) -path system, which is formally defined below.

DEFINITION 5.2. Let W and V' be disjoint vertex sets and let X, Y , and J be sets of paths on $W \cup V'$. We write, for example, $V(J)$ to mean the set of all vertices of all paths in J .

We say that (X, Y, J) is a (W, V') -path system if there exist distinct vertices x and y such that:

- (P1) $X = \{x\}$ if $x \in V'$; otherwise X is a set of two edge-disjoint paths that both start at x and end in V' ;
- (P2) $Y = \{y\}$ if $y \in V'$; otherwise Y is a set of two edge-disjoint paths that both start in V' and end at y ;
- (P3) J is a set of vertex-disjoint paths such that each path in J has both end points in V' ;
- (P4) $d_{X \cup Y \cup J}(v) \leq 1$ for all $v \in V'$;
- (P5) $V(X), V(Y)$, and $V(J)$ are disjoint.

We will often take $X = \{xx', xx''\}$ for some $x', x'' \in V'$ if $x \in W$ and similarly for Y . We will interchangeably think of X, Y , and J both as a set of paths and as the graph which is the union of those paths, but it will always be clear from the context.

We say that the two paths P_1 and P_2 extend (X, Y, J) , if $X \cup Y \cup J \subseteq P_1 \cup P_2$ and each P_i starts at x and ends at y . We refer to x and y as the source and sink, respectively.

The following corollary (of Lemma 4.9) shows how to simultaneously extend a collection of vertex-disjoint (W, V') -path systems, so that the resulting paths are internally vertex-disjoint.

COROLLARY 5.3. Let $n, s \in \mathbb{N}$ and $0 < 1/n \ll \varepsilon, \varepsilon' \ll 1$ and $1/n \ll 1/s$. Let D be an oriented graph with vertex partition $V(D) = W \cup V'$ such that $|V'| = n$ and $\delta^0(D[V']) \geq (1/2 - \varepsilon)n$. For $i \in [s]$, let (X_i, Y_i, J_i) be (W, V') -path systems. Suppose that the sets $V_i := V(X_i \cup Y_i \cup J_i) \cap V'$ for $i \in [s]$ are disjoint and that $|\bigcup_{i \in [s]} V_i| \leq \varepsilon'n$. Then $D \cup \bigcup_{i \in [s]} (X_i \cup Y_i \cup J_i)$ contains paths $P_1, P'_1, \dots, P_s, P'_s$ such that:

- (a) for each $i \in [s]$, P_i and P'_i extend (X_i, Y_i, J_i) ;
- (b) $d_{\bigcup_{i \in [s]} (P_i \cup P'_i)}(v) = 2$ for all $v \in V'$.

Proof. Let $\tilde{V} = \bigcup_{i \in [s]} V_i$, so that $s \leq |\tilde{V}| \leq \varepsilon'n$. Let \tilde{E} be the set of edges used in all the paths in all the path systems (X_i, Y_i, J_i) for all $i \in [s]$. Write $D' = D[V'] - \tilde{E}$.

For each $i \in [s]$ let $P_{i1}, \dots, P_{it(i)}$ be the paths of J_i . We will apply Lemma 4.9 to join the paths of our path systems together. Let a_{ij} and b_{ij} be starting and ending points of P_{ij} , respectively, so $a_{ij}, b_{ij} \in \tilde{V}$. Also, let x_i, x'_i be the two end points in V' of the paths in X and let y_i, y'_i be the two end points in V' of the paths in Y (where possibly $x_i = x'_i$ and/or $y_i = y'_i$).

Let $H := \bigcup_{i \in [s]} T_i$ be a multigraph on $\tilde{V} \subseteq V'$, where

$$T_i := \{x_i a_{i1}, b_{i1} a_{i2}, b_{i2} a_{i3}, \dots, b_{i(t(i)-1)} a_{it(i)}, b_{it(i)} y_i, x'_i y'_i\}.$$

By property (P4) of path systems T_i is a matching and since the V_i are disjoint, then H is a matching on \tilde{V} so $|E(H)| \leq |\tilde{V}| \leq \varepsilon'n \leq |D'|$. Note that $\delta^0(D') \geq \delta^0(D[V']) - |\tilde{V}| \geq (1/2 -$

$\varepsilon - \varepsilon'n \geq 3n/7$. We apply Lemma 4.9 with $D', H, \emptyset, 2\varepsilon'$ playing the roles of D, H_1, S_1, γ and obtain a set of edge-disjoint paths $\mathcal{Q} := \{Q_e : e \in E(H)\}$ such that:

- for each $e = xy \in E(H)$, Q_e is a path from x to y ;
- the paths $Q_e : e \in E(H)$ are vertex-disjoint (since H is a matching);
- $V(\cup \mathcal{Q}) = V$ and $\Delta(\cup \mathcal{Q}) \leq 2$.

For each $i \in [s]$, set

$$P_i := (X_i \cup Y_i \cup J_i) \cup \bigcup_{e \in T_i \setminus \{x'_i y'_i\}} Q_e \quad \text{and} \quad P'_i = Q_{x'_i y'_i}.$$

Note that P_i forms a path by our choice of T_i and that P_i, P'_i extends (X_i, Y_i, J_i) ; thus conditions (a) and (b) of the corollary are satisfied. \square

Our first step towards proving Lemma 5.1 is Lemma 5.4 where we construct a partial decomposition that uses all the edges inside W .

LEMMA 5.4. *Let $n \in \mathbb{N}$ and $0 < 1/n \ll \alpha, \beta, \varepsilon \ll 1$. Let $C \geq 32$. Let D be an oriented graph on n vertices such that $\delta(D) \geq (1 - \varepsilon)n$ and $\text{ex}(D) \geq Cn$. Let $W \subseteq V(D)$ of size $|W| \leq \beta n$. Suppose that $|\text{ex}_D(v)| \leq \alpha n$ for all $v \in V(D) \setminus W$. Then there exists a partial decomposition \mathcal{P} of D such that writing $H = \cup \mathcal{P}$ we have:*

- (i) $H[W] = D[W]$;
- (ii) $\Delta(H) \leq 21|W|$ and $d_H(v) = 18|W|$ for all $v \in V(D) \setminus W$;
- (iii) $\text{ex}(D - H) \geq Cn/2$.

Proof. Let $\gamma > 0$ be such that $\alpha, \beta, \varepsilon \ll \gamma \ll 1$. Let $\ell := |W|$ and let $V' := V(D) \setminus W$. Note that

$$\delta^0(D[V']) \geq (1/2 - \varepsilon - \alpha - \beta)n. \tag{5.1}$$

Let $W_\gamma^\pm := \{w \in W : \text{ex}^\pm(w) \geq (1 - \gamma)n\}$ and $W_0 := W \setminus (W_\gamma^+ \cup W_\gamma^-)$. By Vizing's theorem, $D[W]$ can be decomposed into ℓ (possibly empty) matchings M_1, \dots, M_ℓ . For each $i \in [\ell]$, we partition M_i into matchings

$$\begin{aligned} M_i^0 &:= \{ab \in M_i : a \in W_\gamma^+, b \in W_\gamma^-\}, \\ M_i^+ &:= \{ab \in M_i : a \in W_\gamma^+, b \notin W_\gamma^-\}, \\ M_i^- &:= \{ab \in M_i : a \notin W_\gamma^+, b \in W_\gamma^-\}, \\ M_i' &:= \{ab \in M_i : a \notin W_\gamma^+, b \notin W_\gamma^-\} = M_i \setminus (M_i^0 \cup M_i^+ \cup M_i^-). \end{aligned}$$

Let $m_i^* := |M_i^*|$ for all $* \in \{0, \iota, +, -\}$. Note that for each i

$$m_i^0 + m_i^+ + m_i^- \leq |W_\gamma^+| + |W_\gamma^-|, \tag{5.2}$$

since $M_i^0 \cup M_i^+ \cup M_i^-$ is a matching in which all edges are incident with $W_\gamma^+ \cup W_\gamma^-$.

Suppose that we have found partial decompositions $\mathcal{P}_1, \dots, \mathcal{P}_\ell$ such that writing $H_j = \cup \mathcal{P}_j$, we have for each $j \in [\ell]$,

- (i') \mathcal{P}_j is a partial decomposition of $D_{j-1} := D - (H_1 \cup \dots \cup H_{j-1})$ (and hence, $\cup_{j \in [\ell]} \mathcal{P}_j$ is a partial decomposition of D by Proposition 4.1(b));
- (ii') $H_j[W] = M_j$;
- (iii') $\Delta(H_j) \leq 21$ and $d_{H_j}(v) = 18$ for all $v \in V'$;
- (iv') $|\mathcal{P}_j| = \text{ex}(H_j) = m_i^0 + 2m_i^+ + 2m_i^- + 4$;
- (v') $H_1, \dots, H_j, M_{j+1}, \dots, M_\ell$ are edge-disjoint.

Set $\mathcal{P} := \cup_{j \in [\ell]} \mathcal{P}_j$ and $H := \cup_{j \in [\ell]} H_j$. Clearly, (i) and (ii) hold by (ii') and (iii'), respectively. To see (iii), note that (iv') and (i') imply

$$\text{ex}(D - H) = \text{ex}(D) - \text{ex}(H) \stackrel{(5.2)}{\geq} \text{ex}(D) - 2(|W_\gamma^+| + |W_\gamma^-| + 2)\ell.$$

If $|W_\gamma^+| + |W_\gamma^-| \leq 2C$, then $\text{ex}(D - H) \geq Cn - 4(C + 1)\ell \geq Cn/2$. If $|W_\gamma^+| + |W_\gamma^-| > 2C$, then

$$\begin{aligned} \text{ex}(D - H) &\geq \text{ex}(D) - 3(|W_\gamma^+| + |W_\gamma^-|)\ell \\ &= \frac{1}{2} \sum_{v \in V(D)} |\text{ex}_D(v)| - 3(|W_\gamma^+| + |W_\gamma^-|)\ell \\ &\geq \frac{(1 - \gamma)n}{2} (|W_\gamma^+| + |W_\gamma^-|) - 3(|W_\gamma^+| + |W_\gamma^-|)\ell \\ &\geq (1 - \gamma - 6\beta)n(|W_\gamma^+| + |W_\gamma^-|)/2 \geq Cn/2. \end{aligned}$$

Therefore, to prove the lemma, it suffices to show that such $\mathcal{P}_1, \dots, \mathcal{P}_\ell$ exist.

Suppose for some $i \in [\ell]$, we have already found partial decompositions $\mathcal{P}_1, \dots, \mathcal{P}_{i-1}$ satisfying (i')–(v'). We now construct $\mathcal{P}_i = \mathcal{P}'_i \cup \mathcal{P}_i^+ \cup \mathcal{P}_i^- \cup \mathcal{P}^0$, where \mathcal{P}'_i is a partial decomposition containing the edges of M_i^* for $*$ in $\{+, -, ', 0\}$. We immediately define $\mathcal{P}_i^0 = M_i^0$. We will write $H_i, H'_i, H_i^+, H_i^-, H_i^0$, respectively, for the union of paths in $\mathcal{P}_i, \mathcal{P}'_i, \mathcal{P}_i^+, \mathcal{P}_i^-, \mathcal{P}_i^0$. Let $D_{i-1}^0 := D_{i-1} - H_i^0 - (M_{i+1} \cup \dots \cup M_\ell)$. Note that by (5.1) and (iii'),

$$\delta^0(D_{i-1}^0[V']) \geq \delta^0(D[V']) - 21(i-1) - \ell \geq (1/2 - \gamma)n \quad (5.3)$$

and (by a similar argument as used to bound $\text{ex}(D - H)$) we have

$$\text{ex}(D_{i-1}^0) \geq Cn/2. \quad (5.4)$$

We first construct the partial decomposition \mathcal{P}'_i of D_{i-1}^0 containing M'_i in the following claim.

CLAIM 5.5. *There exists a partial decomposition \mathcal{P}'_i of D_{i-1}^0 such that, recalling $H'_i = \cup \mathcal{P}'_i$, we have:*

- (a₁) $|\mathcal{P}'_i| = 4$ (and there exist vertices x_1, x_2, y_1, y_2 such that two of the paths start at x_1 and end at y_1 and the other two start at x_2 and end at y_2);
- (a₂) $H'_i[W] = M'_i$, $\Delta(H'_i) \leq 4$ and $d_{H'_i}(v) = 2$ for all $v \in V'$.

Proof of Claim. Let x_1, x_2, y_1, y_2 be any four distinct vertices such that $\text{ex}_{D_{i-1}^0}^+(x_j) \geq 2$ and $\text{ex}_{D_{i-1}^0}^-(y_j) \geq 2$ for all $j \in [2]$. Note that such vertices exist by (5.4). Consider $j \in [2]$. If $x_j \in V'$, then set $X_j = \{x_j\}$; if $x_j \in W$, then $x_j \notin W_\gamma^-$ since $\text{ex}_{D_{i-1}^0}(x_j) > 0$. So,

$$d_{D_{i-1}^0}^+(x_j, V') \geq \frac{d_{D_{i-1}^0}(x_j)}{2} - |W| \stackrel{(iii')}{\geq} \frac{(1 - \varepsilon)n - 21(i-1) - 2|W|}{2} \geq 2|W| + 4$$

and we can set $X_j = \{x_j x'_j, x_j x''_j\}$ for some distinct $x'_j, x''_j \in N_{D_{i-1}^0}^+(x_j) \cap V'$. Similarly, if $y_j \in V'$, then set $Y_j = \{y_j\}$; if $y_j \in W$, then set $Y_j = \{y'_j y_j, y''_j y_j\}$ for some distinct $y'_j, y''_j \in N_{D_{i-1}^0}^-(y_j) \cap V'$. Moreover, we may further assume that X_1, X_2, Y_1, Y_2 are vertex-disjoint. Let $U := V(X_1) \cup V(X_2) \cup V(Y_1) \cup V(Y_2)$. Partition M'_i into M^1 and M^2 such that

(by relabelling X_1, X_2, Y_1, Y_2 if necessary) $V(M^j) \cap (X_j \cup Y_j) = \emptyset$ for $j \in [2]$. Let

$$M^1 := \{a_j b_j : j \in [r]\} \text{ and } M^2 := \{a_{r+j} b_{r+j} : j \in [s]\}.$$

For each $j \in [r + s]$, note that $a_j \in W \setminus W_\gamma^+$ and so

$$\begin{aligned} d_{D_{i-1}^0}^-(a_j, V') &\stackrel{\text{(iii')}}{\geq} d_D^-(a_j) - |W| - 21|W| = \frac{d_D(a_j) - \text{ex}_D^+(a_j)}{2} - 22|W| \\ &\geq \frac{(1 - \varepsilon)n - (1 - \gamma)n}{2} - 22\beta n \geq 2\beta n + 4 \geq 2|W| + 4. \end{aligned}$$

By a similar argument, we have $d_{D_{i-1}^0}^+(b_j, V') \geq 2|W| + 4$. So, there exist distinct $a'_1, \dots, a'_{r+s}, b'_1, \dots, b'_{r+s} \in V' \setminus U$ such that $a'_j \in N_{D_{i-1}^-}(a_j)$ and $b'_j \in N_{D_{i-1}^+}(b_j)$ for all $j \in [r + s]$. Let

$$J_1 := \{a'_j a_j b_j b'_j : j \in [r]\}, \quad J_2 := \{a'_{r+j} a_{r+j} b_{r+j} b'_{r+j} : j \in [s]\}.$$

Observe that (X_1, Y_1, J_1) and (X_2, Y_2, J_2) are (W, V') -path systems. Note further that $X_1, Y_1, J_1, X_2, Y_2, J_2$ are vertex-disjoint and their union has size at most $2|W| + 4 \leq 3\beta n$. By considering $(X_1, Y_1, J_1), (X_2, Y_2, J_2)$ and (5.3), Corollary 5.3 implies that $D_{i-1}^0[V'] \cup J_1 \cup J_2$ contains paths P_1, P'_1, P_2, P'_2 such that, for $j \in [2]$, P_j and P'_j extend (X_j, Y_j, J_j) and $d_{P_1 \cup P'_1 \cup P_2 \cup P'_2}(v) = 2$ for all $v \in V'$. Let $\mathcal{P}'_i := \{P_1, P'_1, P_2, P'_2\}$. It is easy to check that \mathcal{P}'_i has the desired properties. \square

In the next claim, we construct the partial decompositions \mathcal{P}_i^+ and \mathcal{P}_i^- of $D'_{i-1} := D_{i-1}^0 - H'_i$ containing M_i^+ and M_i^- , respectively, as follows.

CLAIM 5.6. *There is a partial decomposition $\mathcal{P}_i^+ \cup \mathcal{P}_i^-$ of D'_{i-1} such that, recalling $H_i^\pm = \cup \mathcal{P}_i^\pm$, we have:*

- (b₁) $|\mathcal{P}_i^\pm| = 2m_i^\pm$;
- (b₂) $H_i^\pm[W] = M_i^\pm$, $\Delta(H_i^\pm) = 8$ and $d_{H_i^\pm}(v) = 8$ for all $v \in V'$.

Proof of Claim. First we arbitrarily partition M_i^+ into four matchings, which we denote by N_1, N_2, N_3, N_4 , each of size $\lfloor m_i^+/4 \rfloor$ or $\lceil m_i^+/4 \rceil$. Let $m := |N_1|$ and $N_1 = \{w_j w'_j : j \in [m]\}$.

We show that there exist distinct vertices $z_1, \dots, z_m \in V(D) \setminus V(N_1)$ such that, for all $j \in [m]$, $\text{ex}_{D'_{i-1}}^-(z_j) \geq 2$. Indeed, if $m \leq C/8$, then $\text{ex}(D'_{i-1}) \geq \text{ex}(D_{i-1}^0) - 4 \geq 3Cn/8 \geq 3mn = (|V(N_1)| + m)n$ by (5.4), so we can find such z_j in this case. On the other hand, if $m \geq C/8 \geq 4$, then

$$\begin{aligned} \text{ex}(D'_{i-1}) &\geq \sum_{w \in W_\gamma^+} \text{ex}_{D'_{i-1}}^+(w) \geq |W_\gamma^+| \left((1 - \gamma)n - \sum_{j \in [i-1]} \Delta(H_j) - \Delta(H'_i) - 1 \right) \\ &\stackrel{\text{(iii')}, \text{(a}_2\text{)}}{\geq} |M_i^+| (1 - 2\gamma)n \geq (4m - 3)(1 - 2\gamma)n \geq 3mn = (|V(N_1)| + m)n, \end{aligned}$$

so again we can find the desired z_j .

By a similar argument as used in the proof of Claim 5.5, there exist (W, V') -path systems (W_j, Z_j, \emptyset) for $j \in [m]$ such that, for all $j \in [m]$,

- $W_j = \{w_j w'_j w''_j, w_j w''_j\}$ with $w''_j, w'''_j \in V'$;
- if $z_j \in V'$, then $Z_j = \{z_j\}$; otherwise $Z_j = \{z'_j z_j, z''_j z_j\}$ for some $z'_j, z''_j \in V'$;
- the $2m$ graphs W_j, Z_j with $j \in [m]$ are vertex-disjoint and are subgraphs of D'_{i-1} .

By considering (W, V') -path systems (W_j, Z_j, \emptyset) (and using (5.3) and (a₂)), Corollary 5.3 implies that $D'_{i-1}[V']$ contains paths $P_1, P'_1, \dots, P_m, P'_m$ such that:

- for all $j \in [m]$, P_j and P'_j extend (W_j, Z_j, \emptyset) ;
- $d_{\cup_{j \in [m]}(P_j \cup P'_j)}(v) = 2$ for all $v \in V'$.

Let $\mathcal{P}_{i,1}^+ := \{P_j, P'_j : j \in [m]\}$ and $H_{i,1}^+ := \cup \mathcal{P}_{i,1}^+$. Note that $|\mathcal{P}_{i,1}| = 2m$ (where two paths start at w_j and end at z_j for every $j \in [m]$). Moreover, $H_{i,1}^+[W] = N_1$, $\Delta(H_{i,1}^+) = 2$, $d_{H_{i,1}^+}(v) = 2$ for all $v \in V'$ and $\mathcal{P}_{i,1}$ is a partial decomposition D'_{i-1} by the choice of z_j, w_j . By a similar argument, $D'_{i-1} - H_{i,1}^+$ has edge-disjoint partial decompositions $\mathcal{P}_{i,2}^+, \mathcal{P}_{i,3}^+, \mathcal{P}_{i,4}^+$ such that $\mathcal{P}_i^+ := \cup_{k \in [4]} \mathcal{P}_{i,k}^+$ satisfies (b₁) and (b₂). By a similar argument, we can construct a partial decomposition \mathcal{P}_i^- of $D_i - H'_i - H_i^+$ satisfying (b₁) and (b₂). □

Finally, we let $\mathcal{P}_i = \mathcal{P}'_i \cup \mathcal{P}_i^+ \cup \mathcal{P}_i^- \cup \mathcal{P}^0$. Our sequential construction of partial decompositions in the digraphs with earlier partial decompositions removed means that (i') holds by Proposition 4.1(b). Clearly, (ii') holds. Also (v') holds by our definition of D_{i-1}^0 . Note that (iv') is implied by (a₁), (b₁). Finally, (iii') holds by (a₂) and (b₂). This completes the proof of the lemma. □

In the next lemma, we show how to construct a partial decomposition with few paths that uses all those edges incident with W in the ‘wrong’ direction; this will help us to isolate the vertices of W in later sections.

LEMMA 5.7. *Let $n \in \mathbb{N}$ and $0 < 1/n \ll \alpha, \beta \ll \gamma \ll 1$ and $1/n \ll \varepsilon \ll 1$. Let $C \geq 5$. Let D be an oriented graph on n vertices such that $\delta(D) \geq (1 - \varepsilon)n$ and $\text{ex}(D) \geq Cn$. Let $W \subseteq V(D)$ of size $|W| \leq \beta n$ such that $D[W]$ is empty. Suppose that $|\text{ex}_D(v)| \leq \alpha n$ for all $v \in V(D) \setminus W$. Then there exists a partial decomposition \mathcal{P} of D such that writing $H = \cup \mathcal{P}$ we have:*

- (i) $|\mathcal{P}| \leq 2(2 + 3\beta)n$ (equivalently $\text{ex}(H) \leq 2(2 + 3\beta)n$);
- (ii) if $w \in W$ with $\pm \text{ex}_D(w) \geq 0$, then $N_{D-H}^\mp(w) = \emptyset$;
- (iii) for all $v \in V(D) \setminus W$, $d_H(v) = 2d$ for some $d \leq 4\gamma n$.

Proof. Let $q \in \mathbb{N}$ be such that $\alpha, \beta \ll 1/q \ll \gamma$. Let $\ell := \lceil \gamma n \rceil$ and $p := \lceil n/q \rceil \geq \alpha n, \beta n$. Let $V' := V(D) \setminus W$. Note that

$$\delta^0(D[V']) \geq (1/2 - \varepsilon - \alpha - \beta)n. \tag{5.5}$$

Let

$$\begin{aligned} W^\pm &:= \{w \in W : \text{ex}_D^\pm(w) \geq 0\}, \\ W_\gamma^\pm &:= \{w \in W : \text{ex}^\pm(w) \geq (1 - \gamma)n\}, \\ W_0^\pm &:= W^\pm \setminus W_\gamma^\pm. \end{aligned}$$

Let $D_0 := D[V', W^+] \cup D[W^-, V']$. We start by showing that if we can find a family \mathcal{S} of edge-disjoint (W, V') -path systems satisfying the following properties, then the lemma holds:

- (i') $\mathcal{S} = \{(X_{i,j}, Y_{i,j}, J_{i,j}) : i \in [p], j \in [q]\} \cup \{(X_i, Y_i, J_i) : i \in [p + 1, p + 3\ell]\}$;
- (ii') $\bigcup_{i \in [p], j \in [q]} J_{i,j} \cup \bigcup_{i' \in [p+1, p+3\ell]} J_{i'}$ contains D_0 ;
- (iii') each $J_{i,j}$ and $J_{i'}$ consists of vertex-disjoint paths of length 2 of the form awb for some $a, b \in V'$ and $w \in W$;

- (iv') for all $v \in V'$ we have $2f(v) \leq \text{ex}^+(v)$ and $2g(v) \leq \text{ex}^-(v)$, where $f(v)$ (respectively, $g(v)$) denotes the number of times v appears as a source (respectively, a sink) in \mathcal{S} (recall the definition of source and sink for a (W, V') -path system);
- (v') for all $i \in [p]$, $\{V' \cap V(X_{i,j} \cup Y_{i,j} \cup J_{i,j}) : j \in [q]\}$ are disjoint and $|V' \cap \bigcup_{j \in [q]} V(X_{i,j} \cup Y_{i,j} \cup J_{i,j})| \leq 2\beta n + 4$;
- (vi') for all $i \in [p+1, p+3\ell]$, $|V' \cap V(X_i \cup Y_i \cup J_i)| \leq 2\beta n + 4$.

Let $\mathcal{S}_i := \{(X_{i,j}, Y_{i,j}, J_{i,j}) : j \in [q]\}$ for $i \in [p]$ and $\mathcal{S}_{i'} := \{(X_{i'}, Y_{i'}, J_{i'})\}$ for $i' \in [p+1, p+3\ell]$. It is easy to verify that by repeated application of Corollary 5.3 (once for each \mathcal{S}_i), there exists a set \mathcal{P} of edge-disjoint paths of D with $\mathcal{P} = \mathcal{P}_1 \cup \dots \cup \mathcal{P}_{p+3\ell}$ and $H_i := \cup \mathcal{P}_i$ such that:

- (a) for all $i \in [p]$, we have $|\mathcal{P}_i| = 2q$ with $\mathcal{P}_i = \{P_{i,1}, P'_{i,1}, \dots, P_{i,q}, P'_{i,q}\}$;
- (b) for each $i \in [p]$ and $j \in [q]$, $P_{i,j}$ and $P'_{i,j}$ extend $(X_{i,j}, Y_{i,j}, J_{i,j})$;
- (c) for each $i' \in [p+1, p+3\ell]$, $\mathcal{P}_{i'} = \{P_{i'}, P'_{i'}\}$, where $P_{i'}$ and $P'_{i'}$ extends $(X_{i'}, Y_{i'}, J_{i'})$;
- (d) for all $i \in [p+3\ell]$ and all $v \in V'$, $d_{H_i}(v) = 2$.

Now, we check that the conclusion of the lemma holds for \mathcal{P} as defined above. Note that by the choice of sources and sinks for the path systems, that is, (iv'), \mathcal{P} is a partial decomposition of D . Note also that (i) holds since $|\mathcal{P}| = \sum_{i \in [p+3\ell]} |\mathcal{P}_i| = 2pq + 6\ell \leq 2(2 + 3\beta)n$. Also (ii) holds by (ii'). Finally, (iii) holds by (d) as $p + 3\ell \leq 4\gamma n$. Thus, to prove the lemma, it suffices to show that such \mathcal{S} exists.

Here we give a brief outline of the remainder of the proof. First, we will find all sources and sinks that are required. We split D_0 into

$$D_1 := D[V', W_\gamma^+] \cup D[W_\gamma^-, V'] \quad \text{and} \quad D_2 := D[V', W_0^+] \cup D[W_0^-, V'].$$

The edges of D_1 will be covered by $\mathcal{S}_{p+i'}$ for $i' \in [3\ell]$ and the edges of D_2 will be covered by \mathcal{S}_i for $i \in [p]$.

Finding sources and sinks

First, we define the sources and sinks for the (W, V') -path systems. Choose a multiset $X := \{x_{i,j} : i \in [p], j \in [q]\} \cup \{x_{p+i'} : i' \in [3\ell]\}$ of vertices such that $\text{ex}_D^+(v) \geq 2f(v)$ for all $v \in V(D)$, where $f(v)$ denotes the number of times v appears in X . Note that such X exists since

$$\text{ex}(D) \geq Cn \geq 2(pq + 3\ell) + n.$$

Similarly, choose a multiset $Y := \{y_{i,j} : i \in [p], j \in [q]\} \cup \{y_{p+i'} : i' \in [3\ell]\}$ of vertices such that $\text{ex}_D^-(v) \geq 2g(v)$ for all $v \in V(D)$, where $g(v)$ denotes the number of times v appears in Y . Note that for all $v \in V(D)$,

$$f(v) + g(v) \leq |\text{ex}_D(v)|/2. \tag{5.6}$$

Since $\text{ex}_D(v) \leq \alpha n$ for all $v \in V'$ and $\alpha n \leq p, \ell$, we may assume that, by relabelling if necessary,

- for all $i \in [p]$, the multiset $V' \cap \{x_{i,1}, \dots, x_{i,q}, y_{i,1}, \dots, y_{i,q}\}$ contains no repeated vertices;
- for all $i' \in [3\ell]$, the multiset

$$V' \cap \{x_{p+3i'-2}, x_{p+3i'-1}, x_{p+3i'}, y_{p+3i'-2}, y_{p+3i'-1}, y_{p+3i'}\}$$

contains no repeated vertices.

Note that $x_{i,j}$ and $y_{i,j}$ will be the source and sink for $(X_{i,j}, Y_{i,j}, J_{i,j})$ and $x_{p+i'}$ and $y_{p+i'}$ will be the source and sink for $(X_{p+i'}, Y_{p+i'}, J_{p+i'})$. For $i \in [p]$, let $f_i, g_i : W \rightarrow [0, q]$ be functions such that $f_i(w)$ (and $g_i(w)$) is the number of $j \in [q]$ satisfying $w = x_{i,j}$ (and $w = y_{i,j}$, respectively). Our choices here guarantee that (iv') holds.

Covering edges in D_1

Consider any $i' \in [3\ell]$. If $x_{p+i'} \in V'$, then set $X_{p+i'} = \{x_{p+i'}\}$. If $x_{p+i'} \in W$, then (by our choice of $x_{p+i'}$) we have $\text{ex}(x_{p+i'}) > 0$ and so $d_D^+(x_{p+i'}, V') = d_D^+(x_{p+i'}) \geq n/4$. We can set $X_{p+i'} = \{x_{p+i'}x'_{p+i'}, x_{p+i'}x''_{p+i'}\}$ for some distinct $x'_{p+i'}, x''_{p+i'} \in N_D^+(x_{p+i'}) \subseteq V'$. Similarly, if $y_{p+i'} \in V'$, then set $Y_{p+i'} = \{y_{p+i'}\}$. If $y_{p+i'} \in W$, then we can set $Y_{p+i'} = \{y'_{p+i'}y_{p+i'}, y''_{p+i'}y_{p+i'}\}$ for some distinct $y'_{p+i'}, y''_{p+i'} \in N_D^-(y_{p+i'}) \subseteq V'$. Furthermore, we can assume that all $X_{p+i'}, Y_{p+i'}$ are edge-disjoint and

- for all $i' \in [\ell]$, $X'_{p+3i'-2}, X'_{p+3i'-1}, X'_{p+3i'}, Y'_{p+3i'-2}, Y'_{p+3i'-1}, Y'_{p+3i'}$ are vertex-disjoint, where $X'_j = X_j \cap V'$ and $Y'_j = Y_j \cap V'$.

Let $\hat{X} := \bigcup_{i' \in [3\ell]} X_{p+i'}$ and $\hat{Y} := \bigcup_{i' \in [3\ell]} Y_{p+i'}$. Note that D_0, \hat{X}, \hat{Y} are edge-disjoint and

$$\Delta(\hat{X} \cup \hat{Y}) \leq 6\ell. \tag{5.7}$$

For all $w \in W_\gamma^+$,

$$d_{D_1}^-(w) = d_D^-(w) = \frac{d_D(w) - \text{ex}_D^+(w)}{2} \leq \gamma n/2 \leq \ell$$

and, similarly, $d_{D_1}^+(w) \leq \ell$ for all $w \in W_\gamma^-$. Since $|W| \leq \beta n < \ell$, we deduce $\Delta(D_1) < \ell$. By Vizing's theorem, D_1 can be decomposed into ℓ matchings M'_1, \dots, M'_ℓ . Consider any $i \in [\ell]$. Partition M'_i into three matchings $M_{p+3i-2}, M_{p+3i-1}, M_{p+3i}$ such that for each $j \in [3]$,

$$V(X_{p+3i-3+j} \cup Y_{p+3i-3+j}) \cap V(M_{p+3i-3+j}) = \emptyset.$$

For each $i' \in [3\ell]$ and each vertex $w \in V(M_{p+i'}) \cap W$, if $w \in W_\gamma^\pm$, then by (5.7)

$$d_{D-\hat{X}-\hat{Y}}^\pm(w) \geq d_D^\pm(w) - 6\ell \geq n/4.$$

Hence, by a simple greedy argument, we can extend each $M_{p+i'}$ (with $i' \in [3\ell]$) into a graph $J_{p+i'}$ such that:

- $J_{p+i'}$ consists of precisely $|M_{p+i'}|$ vertex-disjoint paths of length 2 with starting points and end points in V' (and midpoint in W) and $J_{p+i'}$ is vertex-disjoint from $X_{p+i'}$ and $Y_{p+i'}$;
- the 9ℓ different graphs $X_{p+1}, Y_{p+1}, J_{p+1}, \dots, X_{p+3\ell}, Y_{p+3\ell}, J_{p+3\ell}$ are edge-disjoint.

Note that each (X_i, Y_i, J_i) is a (W, V') -path system satisfying (iii') and (vi'). Let $S' := \bigcup_{i' \in [3\ell]} (X_{p+i'} \cup Y_{p+i'} \cup J_{p+i'})$ and note that S' covers all the edges in D_1 .

Covering edges in D_2

We now construct $(X_{i,j}, Y_{i,j}, J_{i,j})$, which cover all the edges in D_2 . Initially, set $X_{i,j} = \{x_{i,j}\}$, $Y_{i,j} = \{y_{i,j}\}$ and let $J_{i,j}$ be empty for all $i \in [p]$ and $j \in [q]$. If $x_{i,j} \in W_\gamma^+$, then $d_{D-S'}^+(x_{i,j}) \geq n/4$ and we can set $X_{i,j} = \{x_{i,j}x'_{i,j}, x_{i,j}x''_{i,j}\}$ for some distinct $x'_{i,j}, x''_{i,j} \in N_{D-S'}^+(x_{i,j}) \subseteq V'$. Similarly, if $y_{i,j} \in W_\gamma^-$, then set $Y_{i,j} = \{y'_{i,j}y_{i,j}, y''_{i,j}y_{i,j}\}$ for some distinct $y'_{i,j}, y''_{i,j} \in N_{D-S'}^-(y_{i,j}) \subseteq V'$. (Later, in Claim 5.8 we will modify those $X_{i,j}$ (respectively, $Y_{i,j}$) for which $x_{i,j} \in W_0^+$ (respectively, $y_{i,j} \in W_0^-$.) We can furthermore assume that all $X_{i,j}, Y_{i,j}$ are edge-disjoint and, for all $i \in [p]$, $X'_{i,1}, \dots, X'_{i,q}, Y'_{i,1}, \dots, Y'_{i,q}$ are vertex-disjoint, where $X'_{i,j} := V' \cap V(X_{i,j})$ and $Y'_{i,j} := V' \cap V(Y_{i,j})$.

Instead of constructing $(X_{i,j}, Y_{i,j}, J_{i,j})$ one at a time, we build them up in rounds, in each round simultaneously adding a little extra to every $(X_{i,j}, Y_{i,j}, J_{i,j})$. Before proving this, we describe somewhat informally how to construct the $J_{i,j}$. Let w_1, \dots, w_s be an enumeration of $W_0^+ \cup W_0^-$. For simplicity, we further assume that none of the w_i is a source or sink, that is, $f(w_i) = 0 = g(w_i)$. For each $i \in [p]$ and $k \in [s]$, we will construct sets $A_{i,k} \subseteq N_{D-S'}^-(w_k)$

and $B_{i,k} \subseteq N_{D-S'}^+(w_k)$ of suitable size (with $|A_{i,k}| = |B_{i,k}| \leq q$) such that for each i the sets $A_{i,1}, \dots, A_{i,s}, B_{i,1}, \dots, B_{i,s} \subseteq V'$ are disjoint. We further guarantee that

$$\bigcup_{i \in [p]} A_{i,k} = N_{D-S'}^-(w_k) \quad \text{if } w_k \in W_0^+ \tag{5.8}$$

$$\text{and } \bigcup_{i \in [p]} B_{i,k} = N_{D-S'}^+(w_k) \quad \text{if } w_k \in W_0^-. \tag{5.9}$$

These sets will be built up in rounds using matchings, but assuming we have these sets, for each i , we define F_i to be the graph with edges

$$E(F_i) := \bigcup_{k \in [s]} \{aw_k, w_kb : a \in A_{i,k}, b \in B_{i,k}\},$$

so that $\bigcup_{i \in [p]} F_i \supseteq D_2$ by (5.8) and (5.9). Note that F_i is the union of vertex-disjoint oriented stars with centres w_1, \dots, w_s , where the star at w_i has an equal number of edges (at most q) entering and exiting w_i . So, it is easy to see that each F_i can be decomposed into $J_{i,1}, \dots, J_{i,q}$ where each $J_{i,j}$ satisfies (iii'). Let us prove all of this formally noting that the fact that some of the w_i are sources or sinks will mean we will have to be more careful about the sizes of our $A_{i,j}$ and $B_{i,j}$.

Let $h : W_0^+ \cup W_0^- \rightarrow [n/2]$ be the function such that if $w \in W_0^\pm$, then

$$h(w) = d_D^\mp(w) = d_{D-S'}^\mp(w) \leq n/2.$$

So, $h(w)$ will correspond to the number of $\{J_{i,j} : i \in [p], j \in [q]\}$ that will contain w . Let $h_1, \dots, h_p : W_0^+ \cup W_0^- \rightarrow [0, q]$ be functions such that, for each $w \in W_0^+ \cup W_0^-$, $\sum_{i \in [p]} h_i(w) = h(w)$ and

$$f_i(w) + g_i(w) + h_i(w) \leq q. \tag{5.10}$$

Indeed, this is possible by considering $h'_i(w) := q - f_i(w) - g_i(w) \geq 0$, so that

$$\sum_{i \in [p]} h'_i(w) \geq pq - f(w) - g(w) \stackrel{(5.6)}{\geq} n - \frac{1}{2}|\text{ex}_D(w)| \geq n/2 \geq h(w)$$

and making a suitable choice of $h_i(w) \leq h'_i(w)$. Here $h_i(w)$ will help determine the number of $\{J_{i,j} : j \in [q]\}$ that will contain w .

Recall w_1, \dots, w_s is an enumeration of $W_0^+ \cup W_0^-$. For $i \in [p]$, let $X_i := \bigcup_{j \in [q]} X_{i,j}$ and $Y_i := \bigcup_{j \in [q]} Y_{i,j}$. Suppose for some $k \in [0, s]$, we have already found $\{A_{i,k'}, B_{i,k'}\}_{i \in [p], k' \in [k]}$ such that:

- (a') for each $i \in [p]$, $V' \cap V(X_i), V' \cap V(Y_i), A_{i,1}, \dots, A_{i,k}, B_{i,1}, \dots, B_{i,k}$ are disjoint;
- (b') for each $i \in [p]$ and $k' \in [k]$, $|A_{i,k'}| = h_i(w_{k'}) + 2g_i(w_{k'})$ and $|B_{i,k'}| = h_i(w_{k'}) + 2f_i(w_{k'})$;
- (c') for each $k' \in [k]$, $A_{1,k'}, \dots, A_{p,k'} \subseteq N_{D-S'}^-(w_{k'})$ are disjoint;
- (d') for each $k' \in [k]$, $B_{1,k'}, \dots, B_{p,k'} \subseteq N_{D-S'}^+(w_{k'})$ are disjoint.

CLAIM 5.8. *If $k = s$, then we can construct \mathcal{S} satisfying (i')–(vi') (so completing the proof of the lemma).*

Proof of Claim. Define F_i to be the graph with edge set

$$E(F_i) := \bigcup_{k' \in [s]} \{aw_{k'}, w_{k'}b : a \in A_{i,k'}, b \in B_{i,k'}\}.$$

Note that by (a')–(d') and our choice of h, h_i we have that $F_1 \cup \dots \cup F_p \supseteq D_2$.

By (c') and (d'), we know that F_i can be decomposed into (W, V') -path systems $(X_{i,j}, Y_{i,j}, J_{i,j})$ (one for each $j \in [q]$) such that $(X_{i,j}, Y_{i,j}, J_{i,j})$ has source $x_{i,j}$ and sink $y_{i,j}$. To see this we colour the edges of F_i with colours from $[q]$ as follows. For each $j \in [q]$ if $x_{i,j} \in W_0^+$ assign colour j to any two out-edges $x_{i,j}x'_{i,j}$ and $x_{i,j}x''_{i,j}$ in F_i at $x_{i,j}$ and (re)set $X_{i,j} = \{x_{i,j}x'_{i,j}, x_{i,j}x''_{i,j}\}$. If $y_{i,j} \in W_0^-$ assign colour j to any two in-edges $y'_{i,j}y_{i,j}$ and $y''_{i,j}y_{i,j}$ in F_i at $y_{i,j}$ and (re)set $Y_{i,j} = \{y'_{i,j}y_{i,j}, y''_{i,j}y_{i,j}\}$. Such edges exist by (b'). Given $w \in W_0^+ \cup W_0^-$ write $c(w)$ for the colour assigned (if any) to edges at w . Let F'_i be the remaining (that is, uncoloured) edges of F_i , noting that there are precisely $h_i(w) \leq q - f_i(w) - g_i(w)$ in-edges and the same number of out-edges at w in F'_i . For each w , pick any set of colours $S_w \subseteq [q] \setminus \{c(w)\}$ with $|S_w| = h_i(w)$. Assign distinct colours of S_w first to the in-edges of F'_i at w and then to the out-edges of F'_i at w . Now, take $J_{i,j}$ to be the edges of F'_i coloured j . In particular,

$$\bigcup_{i \in [p]} \bigcup_{j \in [q]} (X_{i,j} \cup Y_{i,j} \cup J_{i,j}) = \bigcup_{i \in [p]} F_i \supseteq D_2. \quad (5.11)$$

Now, taking $\mathcal{S} = \{(X_{i,j}, Y_{i,j}, J_{i,j}) : i \in [p], j \in [q]\} \cup \{(X_i, Y_i, J_i) : i \in [p+1, p+3\ell]\}$, we see that (i')–(vi') hold. Indeed, (i') and (iii') hold by construction. (ii') holds because we showed \mathcal{S}' covers all edges in D_1 and (5.11) shows all edges in D_2 are covered. We showed (iv') holds when choosing sources and sinks. The disjointness condition in (v') and the edge-disjointness of \mathcal{S} hold by construction. The bounds in (v') and (vi') hold by (iii'). \square

Therefore, we may assume $k \in [0, s-1]$. We show how to find $\{A_{i,k+1}\}_{i \in [p]}$; finding $\{B_{i,k+1}\}_{i \in [p]}$ is similar. Without loss of generality, assume $w_{k+1} \in W_0^+$. We have

$$h(w_{k+1}) = d_D^-(w_{k+1}) \geq \frac{d_D(w_{k+1}) - \text{ex}_D(w_{k+1})}{2} \geq (\gamma - \varepsilon)n/2 \geq \gamma n/4. \quad (5.12)$$

For each $i \in [p]$, set

$$U_i := (V' \cap V(X_i \cup Y_i)) \cup \bigcup_{k' \in [k]} (A_{i,k'} \cup B_{i,k'}).$$

Note that U_i is the set of ‘forbidden’ vertices for $A_{i,k+1}$ and $B_{i,k+1}$ (in order to maintain (a'), (c'), and (d')).

Define an auxiliary bipartite graph F_A with vertex classes A and I as follows. Let $A \subseteq N_{D-S'}^-(w_{k+1})$ be of size $h(w_{k+1}) + 2 \sum_{i \in [p]} g_i(w_{k+1})$; this is possible since

$$h(w_{k+1}) + 2 \sum_{i \in [p]} g_i(w_{k+1}) = d_{D-S'}^-(w_{k+1}) + 0 = d_{D-S'}^-(w_{k+1}).$$

(Note that in the case when we try to find $\{B_{i,k+1}\}_{i \in [p]}$ we use a slightly different calculation[†].) Let I be a multiset consisting of exactly $h_i(w_{k+1}) + 2g_i(w_{k+1})$ copies of $i \in [p]$. Clearly,

$$\begin{aligned} \dagger h(w_{k+1}) + 2 \sum_{i \in [p]} f_i(w_{k+1}) &= d_{D-S'}^-(w_{k+1}) + 2f(w_{k+1}) - 2 \sum_{i' \in [p+1, p+3\ell]} f_{i'}(w_{k+1}) \\ &\leq d_{D-S'}^-(w_{k+1}) + \text{ex}_D^+(w_{k+1}) - \text{ex}_{S'}^+(w_{k+1}) \\ &= d_{D-S'}^-(w_{k+1}) + \text{ex}_{D-S'}^+(w_{k+1}) \\ &= d_{D-S'}^+(w_{k+1}). \end{aligned}$$

$|A| = |I|$. A vertex $v \in A$ is joined to $i \in I$ in F_A if and only if $v \notin U_i$. Note that, for all $v \in A \subseteq V'$, v is in at most,

$$f(v) + g(v) + d_D(v, W) \leq (\alpha/2 + \beta)n \leq \gamma n/8q \stackrel{(5.12)}{\leq} |I|/4q,$$

many of the U_i . Since each $i \in I$ has multiplicity at most q , we deduce

$$d_{F_A}(v) \geq |I| - q \cdot |I|/4q = |I|/2.$$

For each $i \in I$, note that

$$\begin{aligned} |U_i| &\leq 2q + \sum_{k' \in [s]} (2h_i(w_{k'}) + 2f_i(w_{k'}) + 2g_i(w_{k'})) \stackrel{(5.10)}{\leq} 2q(1 + s) \\ &\leq 2q(1 + \beta n) \leq \gamma n/8 \stackrel{(5.12)}{\leq} |A|/2 \end{aligned}$$

implying

$$d_{F_A}(i) \geq |A| - |A|/2 = |A|/2.$$

Therefore, F_A contains a perfect matching M by Hall's Theorem. For each $i \in [p]$, define $A_{i,k+1} := \{v \in A : vi \in M\}$. By a similar argument, there exist $B_{1,k+1}, \dots, B_{p,k+1} \subseteq N_{D-S'}^+(w_{k+1})$ and by construction the sets satisfy (a')–(d'). Indeed, (a') holds by the choice of U_i , (b') holds by the choice of I , and (c') and (d') hold because M is a matching. This completes the proof of the lemma. \square

We now prove Lemma 5.1 using Lemmas 5.4 and 5.7.

Proof of Lemma 5.1. By Lemma 5.4, there exists a partial decomposition \mathcal{P}_1 of D such that writing $H_1 = \cup \mathcal{P}_1$ we have:

- (i') $H_1[W] = D[W]$;
- (ii') $\Delta(H_1) \leq 21|W|$ and $d_{H_1}(v) = 18|W|$ for all $v \in V(D) \setminus W$;
- (iii') $\text{ex}(D - H_1) \geq Cn/2$.

Let $D_1 := D - H_1$. Note that $\delta(D_1) \geq (1 - \varepsilon)n - 21|W| \geq (1 - \varepsilon - 21\beta)n$ and $|\text{ex}_{D_1}(v)| \leq |\text{ex}_D(v)| \leq \alpha n$ for all $v \in V(D) \setminus W$. By Lemma 5.7, there exists a partial decomposition \mathcal{P}_2 of D_1 such that writing $H_2 = \cup \mathcal{P}_2$ we have:

- (i'') $|\mathcal{P}_2| = \text{ex}(H_2) \leq 2(2 + 3\beta)n \leq Cn/4$;
- (ii'') if $w \in W$ with $\text{ex}_D^\pm(w) \geq 0$, then $N_{D_1-H_2}^\mp(w) = \emptyset$;
- (iii'') for all $v \in V(D) \setminus W$, $d_{H_2}(v) = 2d$ for some $d \leq 4\gamma n$.

The lemma holds by setting $\mathcal{P} = \mathcal{P}_1 \cup \mathcal{P}_2$, which is a partial decomposition of D by Proposition 4.1(b). \square

6. The final decomposition

In this section, we prove Theorem 1.5. We prove it in three main steps as discussed in the overview (Section 4.3). We begin with a tournament T that satisfies the hypothesis of Theorem 1.5 but assume that it does not have a perfect decomposition. Gradually, we show that certain subdigraphs of T with various additional properties also do not have a perfect decomposition. Finally, we show that these additional properties are in fact sufficient to guarantee a perfect decomposition, giving the desired contradiction.

6.1. Removing vertices with high excess

The following theorem allows us to remove vertices of high excess from our tournament to leave an almost complete oriented graph D with slightly smaller excess and with the property that a perfect decomposition of D would give a perfect decomposition of T .

THEOREM 6.1. *Let $1/n \ll \beta \ll \alpha \ll \varepsilon$ with n even and let $C > 32$. Let T be an n -vertex tournament with $\text{ex}(T) \geq Cn$. Suppose that T does not have a perfect decomposition. Then there exists a subdigraph D of T with the following properties:*

- (i) D does not have a perfect decomposition;
- (ii) $|D| \geq (1 - \beta)n$ is even;
- (iii) $d_D(v) \geq (1 - \varepsilon)|D|$ for all $v \in V(D)$;
- (iv) $1 \leq |\text{ex}_D(v)| \leq 3\alpha|D|$ for all $v \in V(D)$;
- (v) $\text{ex}(D) \geq (C/4 - 5)n$.

We will need the following three relatively straightforward results before we can prove Theorem 6.1. The first proposition says that any almost regular, almost complete oriented graph has an Eulerian subgraph that uses most of the edges at every vertex and whose removal leaves an acyclic subgraph.

PROPOSITION 6.2. *Let $1/n \ll \varepsilon \ll \varepsilon' \ll 1$. Suppose that D is an n -vertex digraph with $\delta^0(D) \geq \frac{1}{2}(1 - \varepsilon)n$. Then there is an Eulerian digraph $D' \subseteq D$ with $\delta^0(D') \geq \frac{1}{2}(1 - \varepsilon')n$ and such that $D - D'$ is acyclic.*

Proof. Note that $|\text{ex}_D(v)| \leq 2\varepsilon n$ for every $v \in V(D)$. Let K^+ be the multiset of vertices such that each vertex occurs exactly $\text{ex}^+(v)$ times and let K^- be the multiset of vertices such that each vertex occurs exactly $\text{ex}^-(v)$ times. Thus, $|K^+| = |K^-|$ and write $K^+ = \{k_1^+, \dots, k_d^+\}$ and $K^- = \{k_1^-, \dots, k_d^-\}$, where $d = \text{ex}(D)$. Let H be the directed multigraph on $V(D)$ with $E(H) = \{k_i^+ k_j^- : i \in [d]\}$. Note that $\Delta(H) \leq 2\varepsilon n$. We apply Corollary 4.10 and obtain a set of edge-disjoint paths $\mathcal{P} = \{P_e : e \in E(H)\}$ in D such that P_e has the same starting and ending points as e and $\Delta(\cup \mathcal{P}) \leq 4\sqrt{2\varepsilon}n$. By our choice of K^+, K^- , we have that \mathcal{P} is a partial decomposition of D and that $D' := D - \cup \mathcal{P}$ is Eulerian. Also $\delta^0(D') \geq \delta^0(D) - \Delta(\cup \mathcal{P}) \geq \frac{1}{2}(1 - \varepsilon')n$. To ensure that $D - D'$ is acyclic, any cycle in $D - D'$ is added to D' . \square

Given an oriented graph D for which the underlying undirected graph is slightly irregular, the proposition below will be useful in trying to find a small partial decomposition \mathcal{P} of D such that the underlying undirected graph of $D - \cup \mathcal{P}$ is regular. The function f will record the irregularities in the underlying undirected graph of D and the sets T_1, \dots, T_{2tm} obtained will identify the vertex sets of the paths in \mathcal{P} . Some further technical conditions are present that will be useful later.

Recall that, for $U \subseteq X$, we write $I_U : X \rightarrow \{0, 1\}$ for the indicator function of U .

PROPOSITION 6.3. *Let $n, t, m \in \mathbb{N}$ with $tm, 2t \leq n$. Let V be a set with n elements. Let $f : V \rightarrow [m]$ be a function with $m := \max_{v \in V} f(v)$. Suppose $x_1, \dots, x_{2tm}, y_1, \dots, y_{2tm}$ are elements of V (with repetitions) such that $x_i, y_i, x_{tm+i}, y_{tm+i}$ are distinct for each $i \in [tm]$. Then we can find a collection of sets $T_1, \dots, T_{2tm} \subseteq V$ such that:*

- (i) for all $v \in V$, $\sum_{i \in [2tm]} I_{T_i}(v) = f(v) + (2t - 1)m$;
- (ii) $|T_i| \geq (1 - 1/t)n$ for all $i \in [2tm]$;
- (iii) $x_i, y_i \in T_i$ for all $i \in [2tm]$.

Proof. Given any U , take an arbitrary partition of $V \setminus U$ into sets A_1, \dots, A_t with $|A_i| \leq n/t$ for all $i \in [t]$ (we allow empty sets in the partition). Then writing $B_i := V \setminus A_i$, set $\mathcal{S}_U := \{B_1, \dots, B_t\}$. Note that for each $v \in V$,

$$\sum_{S \in \mathcal{S}_U} I_S(v) = I_U(v) + t - 1.$$

Since $f(v) \leq m$ for all $v \in V$, we can find sets U_1, \dots, U_m such that $f \equiv I_{U_1} + \dots + I_{U_m}$. Taking $\mathcal{S} = \bigcup_{i \in [m]} \mathcal{S}_{U_i}$, we have $|\mathcal{S}| = tm$ and

$$\sum_{S \in \mathcal{S}} I_S(v) = \sum_{i \in [m]} (I_{U_i}(v) + t - 1) = f(v) + (t - 1)m.$$

Write S_1, \dots, S_{tm} for the sets in \mathcal{S} . For $i \in [tm]$, let

$$T_i := S_i \cup \{x_i, y_i\} \text{ and } T_{tm+i} := V \setminus (\{x_i, y_i\} \setminus S_i).$$

Let $\mathcal{T} := \{T_i : i \in [2tm]\}$. Note $|T_i| \geq (1 - 1/t)n$ and $x_i, y_i \in T_i$ for all $i \in [2tm]$. For all $v \in V$,

$$\begin{aligned} \sum_{i \in [2tm]} I_{T_i}(v) &= \sum_{i \in [tm]} (I_{T_i}(v) + I_{T_{tm+i}}(v)) \\ &= \sum_{i \in [tm]} (I_{S_i}(v) + 1) = f(v) + (2t - 1)m. \end{aligned} \quad \square$$

The following Lemma shows how to decompose any almost complete Eulerian oriented graph into a small number of cycles. Some extra technical conditions are placed on the cycles that will be useful later.

LEMMA 6.4. *Let $n \in \mathbb{N}$ with $1/n \ll \varepsilon \ll 1$. Suppose D is an n -vertex Eulerian oriented graph with $\delta^0(D) \geq \frac{1}{2}(1 - \varepsilon)n$. Suppose $\phi : V(D) \rightarrow [n]$ satisfies $\sum_{v \in V(D)} \phi(v) \geq 4n$. Then we can decompose D into $t \leq n$ cycles C_1, \dots, C_t where each cycle is assigned two distinct representatives $x_i, y_i \in V(C_i)$ such that no vertex $v \in V(D)$ occurs as a representative more than $\phi(v)$ times.*

Proof. We assume $\frac{1}{2}(1 - \varepsilon)n$ is an integer. For $x \in V(D)$, write $f(x) = \frac{1}{2}(d_D(x) - (1 - \varepsilon)n) \geq 0$. Let $t = \lceil \varepsilon^{-1/2} \rceil$ and $m = \max_{x \in V(D)} f(x)$, so $m \leq \varepsilon n$ and $tm \leq 2\sqrt{\varepsilon}n \leq n$.

Let M be the multiset of vertices in which $v \in V(D)$ occurs $\phi(v)$ times, so that $|M| \geq 4n$ and no vertex occurs more than n times. Let m_1, m_2, \dots be an ordering of the elements of M (with multiplicity) from most frequent to least frequent. For each $i \in [tm]$, write $(x_i, y_i, x_{tm+i}, y_{tm+i}) = (m_i, m_{n+i}, m_{2n+i}, m_{3n+i})$. Note that, as vertices, $x_i, y_i, x_{tm+i}, y_{tm+i}$ are distinct (because no vertex v occurs more than n times in M).

By Proposition 6.3, we can find sets $T_1, \dots, T_{2tm} \subseteq V(D)$ and vertices $x_1, \dots, x_{2tm}, y_1, \dots, y_{2tm} \in V(D)$ such that:

- (i') for all $v \in V$, $\sum_{i \in [2tm]} I_{T_i}(v) = f(v) + (2t - 1)m$;
- (ii') $|T_i| \geq (1 - 1/t)n \geq (1 - \sqrt{\varepsilon})n$ for all $i \in [2tm]$;
- (iii') each T_i is assigned two distinct representatives $x_i, y_i \in T_i$;
- (iv') no vertex $v \in V(D)$ occurs as a representative more than $\phi(v)$ times.

For $i \in [2tm]$, let $S_i := V(D) \setminus T_i$ and H_i be the multidigraph on $V(D)$ with $E(H_i) = \{x_i y_i, y_i x_i\}$. Let $H = \bigcup_{i \in [2tm]} H_i$. Note that $\Delta(H) \leq 4tm \leq 8\sqrt{\varepsilon}n$ and $|S_i| \leq \sqrt{\varepsilon}n$. Apply Lemma 4.9 with $(D, H_i, S_i, 4\sqrt{\varepsilon})$ playing the role of (D, H_i, S_i, γ) to obtain edge-disjoint cycles C_1, \dots, C_{2tm} such that $V(C_i) = T_i$ for each i .

Now, by our choices of T_1, \dots, T_{2tm} , we have that $C := C_1 \cup \dots \cup C_{2tm}$ satisfies $d_C(x) = 2f(x) + 2(2t - 1)m$ by (i') and so $D - C$ is a regular Eulerian digraph with $\delta(D - C) \geq$

$(1 - \varepsilon)n - 4tm \geq 3n/7$. By Lemma 4.4 and Theorem 4.6, $D - C$ can be decomposed into $s \leq n/2$ Hamilton cycles. Each of these cycles is assigned two distinct representatives from $M' = M \setminus \{x_1, \dots, x_{2tm}, y_1, \dots, y_{2tm}\}$ arbitrarily (this is possible since $|M'| \geq 2n$ and no vertex occurs more than n times in M'). Thus, altogether we obtain a decomposition of D into $t \leq n/2 + 2tm \leq n$ cycles with representatives as desired. \square

We now prove Theorem 6.1.

Proof of Theorem 6.1. Fix parameters $\varepsilon_0, \varepsilon_2, \varepsilon'_2, \varepsilon_3$ such that $\beta \ll \alpha \ll \varepsilon_0 \ll \varepsilon_2 \ll \varepsilon'_2 \ll \varepsilon_3 \ll \varepsilon$. Let

$$W^\pm := \{v \in V(T) : \text{ex}^\pm(v) \geq \alpha n\}, \quad W := W^+ \cup W^- \quad \text{and} \quad \overline{W} := V(T) \setminus W.$$

We further guarantee $|W|$ and hence, $|\overline{W}|$ is even by moving an arbitrary vertex $v \in \overline{W}$ to W if $|W|$ is odd; in this case v is added to W^+ if $\text{ex}(v) > 0$ and to W^- if $\text{ex}(v) < 0$. Since T does not have a perfect decomposition, Theorem 3.5 implies $\text{ex}(T) < n^{19/10}$. In particular,

$$|W| \leq 1 + 2\text{ex}(T)/\alpha n \leq \beta n.$$

So, we can apply Lemma 5.1 where $(\alpha, \beta, \varepsilon_0/10, \varepsilon_0/10, C)$ play the role of $(\alpha, \beta, \gamma, \varepsilon, C)$ to obtain a partial decomposition \mathcal{P}_0 of T such that, writing $D_0 := T - \cup \mathcal{P}_0$, we have:

- (a₁) D_0 does not have a perfect decomposition;
- (a₂) $d_{D_0}(v) = d$ for all $v \in \overline{W}$ and some odd $d \geq (1 - \varepsilon_0)n$;
- (a₃) $E(D_0[W]) = E_{D_0}(\overline{W}, W^+) = E_{D_0}(W^-, \overline{W}) = \emptyset$;
- (a₄) $\text{ex}(D_0) \geq Cn/4$;
- (a₅) $|\text{ex}_{D_0}(v)| \leq \alpha n$ for all $v \in \overline{W}$.

Since T does not have a perfect decomposition, (a₁) holds. Note that (a₂), (a₃), and (a₄) follow from Lemma 5.1(i), (ii) and (iii), and (iv), respectively. Finally, (a₅) follows by our choice of W and the fact that \mathcal{P} is a partial decomposition of T .

Let \mathcal{P} be a partial decomposition of D_0 such that every path in \mathcal{P} is of the form w^+v, vw^- , or w^+vw^- for some $w^+ \in W^+$, $w^- \in W^-$, $v \in \overline{W}$. We further assume that first the number of paths in \mathcal{P} of type w^+vw^- is maximal and, subject to this, that \mathcal{P} has maximal size.

Let

$$D_1 := D_0 - \cup \mathcal{P}, \quad D_2 := D_1 - W = D_0 - W.$$

Note that:

- (b₁) $\delta(D_2) \geq d - |W| \geq (1 - \varepsilon_0 - \beta)n \geq (1 - \varepsilon_2)n$;
- (b₂) for every $v \in \overline{W}$, $|\text{ex}_{D_2}(v)| \leq |\text{ex}_{D_1}(v)| + |W| \leq 2\alpha n$;
- (b₃) $\delta^0(D_2) \geq \frac{1}{2}[\delta(D_2) - \max_v |\text{ex}_{D_2}(v)|] \geq \frac{1}{2}(1 - \varepsilon'_2)n$.

CLAIM 6.5. $|\mathcal{P}| < 4n$.

Proof of Claim. Suppose the contrary that $|\mathcal{P}| \geq 4n$. By Proposition 6.2, we can find an Eulerian subgraph D_3 of D_2 such that $\delta^0(D_3) \geq \frac{1}{2}(1 - \varepsilon_3)n$ and $D_2 - D_3$ is acyclic. Let $R := D_1 - D_3$. By (a₃), any cycle in R lies in $R[\overline{W}] = D_2 - D_3$. Hence, R is acyclic. By Proposition 2.6, R has a perfect decomposition \mathcal{P}_1 , which is a partial decomposition of D_0 by Proposition 4.1(d) and (b).

We now show that $D_0 - R = \cup \mathcal{P} \cup D_3$ has a perfect decomposition \mathcal{P}' , which will contradict (a₁) (since then $\mathcal{P}_1 \cup \mathcal{P}'$ is a partial decomposition of D_0 by Proposition 4.1(b)). Note that each path in \mathcal{P} has a unique vertex in \overline{W} . For each $v \in \overline{W}$, write $\phi(v)$ for the number of paths in \mathcal{P} that contain v . Then $\sum_{v \in \overline{W}} \phi(v) = |\mathcal{P}| \geq 4n$.

By Lemma 6.4 (with ε_3 playing the role of ε), we can decompose D_3 into $t \leq n$ cycles C'_1, \dots, C'_t such that each cycle is assigned two distinct representative vertices $x_i, y_i \in C_i$ such that each vertex v occurs as a representative at most $\phi(v)$ times. In particular, we can assign two distinct paths $P_i, Q_i \in \mathcal{P}$ to C_i such that $V(P_i) \cap V(C_i) = x_i$ and $V(Q_i) \cap V(C_i) = y_i$ and $P_1, \dots, P_t, Q_1, \dots, Q_t$ are distinct paths of \mathcal{P} . Now, construct \mathcal{P}' from \mathcal{P} by replacing for each $i = 1, \dots, t$ the paths P_i and Q_i by the paths $P_i x_i C_i y_i Q_i$ and $Q_i y_i C_i x_i P_i$. Now, we see $|\mathcal{P}'| = |\mathcal{P}|$ and that the paths in \mathcal{P}' have the same start and end points as those in \mathcal{P} , so that \mathcal{P}' is a partial decomposition of D_0 by Proposition 4.1(c). Finally, by construction

$$\cup \mathcal{P}' = \cup \mathcal{P} \cup C'_1 \cup \dots \cup C'_t = \cup \mathcal{P} \cup D_3 = D_0 - R,$$

as required. □

It turns out that if $\text{ex}_{D_2}(v) \neq 0$ for all $v \in \overline{W}$, then one can relatively easily prove the theorem by taking $D = D_2$. However, in order to fulfil condition (iv), we must deal with vertices for which $\text{ex}_{D_2}(v) = 0$: this is not hard but is technically cumbersome. We will modify \mathcal{P} by extending some of its paths. Let

$$U^\pm := \{v \in \overline{W} : \text{ex}_{D_2}^\pm(v) > 0\}, \quad U^0 := \overline{W} \setminus (U^+ \cup U^-),$$

$$U_\pm^0 := U^0 \cap \{v \in \overline{W} : \text{ex}_{D_0}^\mp(v) > 0\}.$$

Note that U_+^0 and U_-^0 partition U^0 (since $\text{ex}_{D_0}(u) \neq 0$ by (a₂)). For each $u \in U_+^0$ (and $u \in U_-^0$), let $P_u \in \mathcal{P}$ be a path ending (and starting) at u (such a path exists since $\text{ex}_{D_0}(u) \neq 0$ by (a₂)). Let $\mathcal{P}'_\pm := \{P_u : u \in U_\pm^0\} \subseteq \mathcal{P}$ and let $\mathcal{P}' := \mathcal{P}'_- \cup \mathcal{P}'_+$. Our aim is to extend each path in \mathcal{P}' , so that its starting and ending points avoid U^0 .

We show for later that $\text{ex}(D_2)$ is large. By the maximality of \mathcal{P} , we have

$$N_{D_1}^\pm(W^\pm) \subseteq U^\pm \cup U_\pm^0. \tag{6.1}$$

Together with Proposition 4.2(c), we have

$$\text{ex}(D_2) = \text{ex}(D_1[\overline{W}]) = \text{ex}(D_1) = \text{ex}(D_0) - |\mathcal{P}| \stackrel{(a_4), \text{Claim 6.5}}{>} (Cn/4) - 4n > 4n. \tag{6.2}$$

Our aim is to extend each path in \mathcal{P}' , so that its starting and ending points avoid U^0 . In fact, we replace \mathcal{P}' by \mathcal{Q}' using the following claim.

CLAIM 6.6. *There exists a partial decomposition \mathcal{Q}' of $\cup \mathcal{P}' \cup D_1 = D_0 - \cup(\mathcal{P} \setminus \mathcal{P}')$ such that:*

- (c₁) $\text{ex}(\cup \mathcal{Q}' - W) \leq |U^0| \leq n$;
- (c₂) $\cup \mathcal{P}' \subseteq \cup \mathcal{Q}'$;
- (c₃) $\Delta(\cup \mathcal{Q}' - W) \leq 2\varepsilon_3 n$;
- (c₄) $1 \leq \text{ex}_{D_2 - \cup \mathcal{Q}'}^\pm(u) \leq 2\alpha n$ if $w \in U_\mp^0 \cup U^\pm$.

Proof of Claim. We will show how to extend the paths in \mathcal{P}'_\pm to obtain sets of paths \mathcal{Q}_\pm and we will take $\mathcal{Q} = \mathcal{Q}_+ \cup \mathcal{Q}_-$. We show how to construct \mathcal{Q}_+ ; the construction of \mathcal{Q}_- follows similarly.

For each $u \in U_+^0$, pick a vertex $b_u \in U^-$ such that no $v \in U^-$ is chosen more than $\text{ex}_{D_2}(v) - 1$ times (which is possible as $|U_+^0| \leq n \leq \text{ex}(D_2) - n$ by (6.2)) and let $e_u = ub_u$. Define a digraph H on $V(D)$ with edge set $\{e_u : u \in U_+^0\}$. Note that $\Delta(H) \leq 2\alpha n$ by (b₂). We apply Corollary 4.10 with $D_2, H, 2\alpha$ playing the roles of D, H, γ to obtain a set of edge-disjoint paths $\mathcal{P}''_+ := \{P'_u : u \in U_+^0\}$ in D_2 such that each P'_u starts at u and ends at b_u and $\Delta(\cup \mathcal{P}'') \leq \varepsilon_3 n$. Recalling that for $u \in U_+^0$, the path P_u is a single edge starting at W^+ and ending at u , we

see that the path $P_u P'_u$ starts at W^+ and ends at b_u . Let $D_1^+ := \cup \mathcal{P}'_+ \cup D_1$. By our choices of \mathcal{P}'_+ , b_u and (6.1), $\mathcal{Q}_+ := \{P_u P'_u : u \in U_+^0\}$ is a partial decomposition of $D_1^+ - W^-$. Moreover, we have

$$\text{ex}_{D_2 - \cup \mathcal{Q}_+}(w) \begin{cases} \in [\min\{\text{ex}_{D_2}(w), -1\}, -1] \stackrel{(b_2)}{\subseteq} [-2\alpha n, -1] & \text{if } w \in U_+^0 \cup U^-; \\ = \text{ex}_{D_2}(w) & \text{if } w \in U_-^0 \cup U^+, \end{cases}$$

where the first case follows since $\text{ex}_{\cup \mathcal{Q}_+}(u) = \text{ex}_{\cup \mathcal{P}'_+}(u) = 1$ for all $u \in U_+^0$, and by our choice of $b_u \in U^-$. By (a₃) and Proposition 4.2(a) (with $(D_1^+, \emptyset, W^-, V(D) \setminus W^-)$ playing the role of (D, A^+, A^-, R)), we can extend \mathcal{Q}_+ to a partial decomposition $\mathcal{Q}'_+ = \{Q'_u : u \in U_+^0\}$ of D_1^+ such that for all $u \in U_+^0$ we have:

- (d₁) $Q'_u - W^- = P_u P'_u$;
- (d₂) Q'_u is a path from W^+ to $U^- \cup W^-$;
- (d₃) $Q'_u - Q'_u[V(D) \setminus W^+] = P_u$;
- (d₄) $\Delta(\cup \mathcal{Q}'_+ - W) \leq \varepsilon_3 n$;
- (d₅) for all $w \in \overline{W}$,

$$\text{ex}_{D_2 - \cup \mathcal{Q}'_+}(w) \begin{cases} \in [-2\alpha n, -1] & \text{if } w \in U_+^0 \cup U^-; \\ = \text{ex}_{D_2}(w) & \text{if } w \in U_-^0 \cup U^+. \end{cases}$$

By a similar argument, we can find a corresponding partial decomposition $\mathcal{Q}'_- = \{Q'_u : u \in U_-^0\}$ of $\cup \mathcal{P}'_- \cup D_1$ edge disjoint from $\cup \mathcal{Q}'_+$. By setting $\mathcal{Q}' := \mathcal{Q}'_+ \cup \mathcal{Q}'_-$, our claim follows. Note that (c₂), (c₃), and (c₄) follow from (d₃), (d₄), and (d₅), respectively, while (c₁) follows from (d₁) and the fact that $|\mathcal{Q}'_{\pm}| = |U_{\pm}^0|$. □

Let

$$D_3 := \cup \mathcal{P}' \cup D_1 - \cup \mathcal{Q}' = D_0 - \cup (\mathcal{P} \setminus \mathcal{P}') - \cup \mathcal{Q}',$$

$$D := D_3 - W = D_2 - \cup \mathcal{Q}'.$$

We show that D satisfies the conclusion of the theorem. In order to prove (i), if D has a perfect decomposition, then so does D_3 by (a₃) and Proposition 4.2(b), and hence, so does D_0 since $(\mathcal{P} \setminus \mathcal{P}') \cup \mathcal{Q}'$ is a partial decomposition of D_0 . This contradicts (a₁), so D has no perfect decomposition and so, (i) holds. Our choice of W implies (ii). Note that (iii) follows from (b₁) and (c₃), and (iv) follows from (c₄). Finally, to see (v),

$$\text{ex}(D) \geq \text{ex}(D_2) - \text{ex}(\cup \mathcal{Q}' - W) \stackrel{(6.2), (c_1)}{\geq} Cn/4 - 5n$$

as required. □

6.2. Balancing the number of positive and negative excess vertices

Given the oriented graph D produced by Theorem 6.1, the following theorem produces a digraph D' that has the same properties as D (with slightly weaker parameters) but with the additional property that the number of vertices of positive excess is almost the same as the number of vertices with negative excess. Recall that for a digraph D , $U^+(D)$ (respectively, $U^-(D)$) denotes the set of vertices of D with positive (respectively, negative) excess.

THEOREM 6.7. *Let $1/n \ll 1/C \ll \alpha, \beta \ll \varepsilon \ll \lambda, \varepsilon' \ll 1$ with n even. Suppose that D is an n -vertex oriented graph, where $\text{ex}(D) \geq Cn$, $\delta(D) \geq (1 - \varepsilon)n$, and $1 \leq |\text{ex}_D(v)| \leq \alpha n$ for all $v \in V(D)$. Suppose that D does not have a perfect decomposition. Then there exists a subgraph D' of D with the following properties:*

- (i) D' does not have a perfect decomposition;
- (ii) $|D'| \geq (1 - \beta)n$ with $|D'|$ even;
- (iii) $\delta(D') \geq (1 - \varepsilon'/2)n \geq (1 - \varepsilon')|D'|$;
- (iv) $1 \leq |\text{ex}_{D'}(v)| \leq \alpha n \leq 2\alpha|D'|$ for all $v \in V(D')$;
- (v) $\text{ex}(D') \geq \lambda Cn/32 \geq \lambda C|D'|/32$;
- (vi) $||U^-(D')| - |U^+(D')|| \leq 2\lambda|D'|$.

Proof. We introduce a parameter ε_1 satisfying $\varepsilon \ll \varepsilon_1 \ll \varepsilon' \ll 1$. Let us write $U^\pm := U^\pm(D)$. If $||U^-| - |U^+|| \leq \lambda n$ then we can take $D' = D$ and we are done, so assume without loss of generality that $|U^-| > |U^+| + \lambda n$. We make the following claim.

CLAIM 6.8. *There exist sets $X \subseteq U^+$ and $Z \subseteq U^-$ satisfying the following.*

- (a₁) $|X| \leq \beta n$ and $|X|$ is even.
- (a₂) Z can be partitioned into sets $Z_x : x \in X$ with $|Z_x| \leq \text{ex}_D(x)$ and $Z_x \subseteq N_D^+(x)$.
- (a₃) $n < \sum_{x \in X} \text{ex}_D(x) \leq (1 + 2\alpha)n \leq (1 - \lambda/4)\text{ex}(D)$ and $\sum_{z \in Z} \text{ex}_D^-(z) \leq (1 - \lambda/4)\text{ex}(D)$.
- (a₄) $||U^-| - |U^+| - 2|Z|| \leq \lambda n$ or equivalently $||U^- \setminus Z| - |U^+ \cup Z|| \leq \lambda n$.

Proof of Claim. Assume βn is an even integer and let X' be the set of βn vertices of U^+ of highest excess. Then

$$\sum_{x \in X'} \text{ex}_D(x) \geq \beta n \text{ex}(D) / |U^+| > C\beta n^2 / (n/2) = 2C\beta n > n.$$

Now, we remove suitable vertices from X' to obtain a set X such that

$$n < \sum_{x \in X} \text{ex}_D(x) \leq (1 + 2\alpha)n \quad \text{and} \quad |X| \text{ even.}$$

This is possible as $|\text{ex}_D(v)| \leq \alpha n$ for all $v \in V(D)$. For each $x \in X$, we have

$$\begin{aligned} |N_D^+(x) \cap U^-| &\geq \frac{1}{2}d_D^+(x) - |U^+| \geq \frac{1}{2}(1 - \varepsilon)n - |U^+| \\ &= \frac{1}{2}(|U^-| - |U^+|) - \frac{1}{2}\varepsilon n \geq \frac{1}{2}(|U^-| - |U^+|) - \frac{1}{4}\lambda n. \end{aligned}$$

Thus, for each $x \in X$, we can greedily pick disjoint $Z'_x \subseteq N_D^+(x) \cap U^-$ with $|Z'_x| \leq \text{ex}_D(x)$ and $|\cup_{x \in X} Z'_x| = \frac{1}{2}(|U^-| - |U^+|) - \frac{1}{4}\lambda n$. Let Y be the $\frac{1}{4}\lambda n$ vertices of lowest excess (that is, of highest negative excess) in $Z' := \cup_{x \in X} Z'_x$. Set $Z_x := Z'_x \setminus Y$ and $Z := Z' \setminus Y$. Hence, $Z := \cup_{x \in X} Z_x$ and $|Z| = \frac{1}{2}||U^+| - |U^-|| - \frac{1}{2}\lambda n$. Also

$$\begin{aligned} \sum_{z \in Z} \text{ex}^-(z) &= \sum_{z' \in Z'} \text{ex}^-(z') - \sum_{y \in Y} \text{ex}^-(y) \\ &\leq (1 - \lambda/4) \sum_{z' \in Z'} \text{ex}^-(z') \leq (1 - \lambda/4)\text{ex}(D). \end{aligned} \quad \square$$

We will construct the final graph D' such that $V(D') = V(D) \setminus X$, where $U^+(D') = (U^+ \setminus X) \cup Z$ and $U^-(D') = U^- \setminus Z$, and where $\text{ex}_{D'}(z) = 1$ for all $z \in Z$.

For each $z \in Z$, we write x_z for the vertex x such that $z \in Z_x$. Note that $x_z z \in E(D)$.

CLAIM 6.9. *There exists a partial decomposition $\mathcal{P}_Z := \{x_z Q_z : z \in Z\}$ of D such that each Q_z is a non-empty path in $D - X$ starting at z and ending in $U^- \setminus Z$. Moreover, $\text{ex}_{D - \cup \mathcal{P}_Z}(v) \neq 0$ for all $v \in V(D) \setminus X$ and $\Delta(\cup \mathcal{P}_Z) \leq \varepsilon_1 n$.*

Proof of Claim. For each $z \in Z$, pick a vertex $b_z \in U^- \setminus Z$ such that no $v \in U^-$ is chosen more than $\text{ex}_D^-(v) - 1$ times (which is possible as $|Z| \leq n \leq \lambda \text{ex}(D)/4 \leq \text{ex}(D) - \sum_{z \in Z} \text{ex}_D^-(z)$ by (a₃)) and let $e_z = zb_z$. Define a digraph H on $V(D) \setminus X$ with edge set $\{e_z : z \in Z\}$. Note that $\Delta(H) \leq \alpha n \leq 2\alpha|D - X|$. We apply Corollary 4.10 with $D - X, H, 2\alpha$ playing the roles of D, H, γ and obtain a set of edge-disjoint paths $\mathcal{Q} := \{Q_z : z \in Z\}$ such that each Q_z starts at z and ends at b_z and $\Delta(\cup \mathcal{Q}) \leq \varepsilon_1 n/2$. Our claim follows by our choice of \mathcal{Q} . \square

Let $D_1 := D - \cup \mathcal{P}_Z$ and write $\mathcal{Q}_Z := \{Q_z : z \in Z\}$.

CLAIM 6.10. *There exists a partial decomposition \mathcal{P}_1 of D_1 such that $\cup \mathcal{P}_1 \subseteq D_1 - X$, $\text{ex}(\cup \mathcal{P}_1) \leq n$, $\Delta(\cup \mathcal{P}_1) \leq \varepsilon_1 n$ and $\text{ex}_{\cup \mathcal{P}_1}(v) \neq 0$ if $v \notin X \cup Z$ and $\text{ex}_{\cup \mathcal{P}_1}(v) = 0$ otherwise.*

Proof of Claim. Let H be any digraph on $V(D) \setminus (X \cup Z)$ with edges from U^+ to U^- such that $1 \leq d_H(v) \leq |\text{ex}_{D_1}(v)|$ for all $v \in V(D) \setminus (X \cup Z)$. Note that $\Delta(H) \leq \alpha n \leq 2\alpha|D - X|$. By deleting edges of H if necessary, we may assume that H has at most n edges. We apply Corollary 4.10 with $D_1 - X, H, 2\alpha$ playing the roles of D, H, γ and obtain the desired partial decomposition \mathcal{P}_1 . \square

Let $D_2 := D_1 - \cup \mathcal{P}_1$. Note that $\delta(D_2) \geq (1 - 3\varepsilon_1)n$.

CLAIM 6.11. *There exists a partial decomposition \mathcal{P}_2 of D_2 such that, writing $D_3 := D_2 - \cup \mathcal{P}_2$, we have $d_{D_3}(x) = 0$ for all $x \in X$, $d_{D_3}(v) \geq (1 - \varepsilon'/2)n$ for all $v \notin X$, $\text{ex}_{D_3}(z) = 0$ for all $z \in Z$, and $\text{ex}(D_3) \geq \lambda Cn/32$*

Proof of Claim. Let

$$m := \max \left\{ \sum_{x \in X} \text{ex}_D(x), \sum_{z \in Z} \text{ex}_D^-(z) \right\}.$$

Recall that $\delta(D_2) \geq (1 - 3\varepsilon_1)n$ and

$$\text{ex}(D_2) \geq \text{ex}(D) - 2n \geq (1 - \lambda/8)\text{ex}(D) \stackrel{(a_3)}{\geq} \lambda \text{ex}(D)/8 + m \geq \lambda Cn/8 + m.$$

Let H be a digraph on $V(D)$ with m edges from U^+ to U^- such that $d_H(v) = |\text{ex}_{D_2}(v)|$ for all $v \in X \cup Z$ and $d_H(v) \leq |\text{ex}_{D_2}(v)|$ otherwise. (Such an H exists by the calculation above.) Note that $\Delta(H) \leq \alpha n$. We apply Corollary 4.10 with D_2, H, α playing the roles of D, H, γ and obtain a partial decomposition \mathcal{P}'_2 of D_2 such that, writing $D'_2 := D_2 - \cup \mathcal{P}'_2$, we have

$$\text{ex}_{D'_2}(v) = 0 \text{ for all } v \in X \cup Z, \text{ex}(D'_2) \geq \lambda Cn/8 \text{ and } \Delta(\cup \mathcal{P}'_2) \leq \varepsilon_1 n.$$

We now apply Lemma 5.1 with $(D'_2, X, \alpha, \beta, \varepsilon_1, 4\varepsilon_1, \lambda C/8)$ playing the roles of $(D, W, \alpha, \beta, \gamma, \varepsilon, C)$. We obtain a partial decomposition \mathcal{P}''_2 of D'_2 such that, writing $D_3 := D'_2 - \cup \mathcal{P}''_2$, we have

$$d_{D_3}(x) = 0 \text{ for all } x \in X, \text{ex}(D_3) \geq \lambda Cn/32 \text{ and } \Delta(\cup \mathcal{P}''_2 - X) \leq 5\varepsilon_1 n.$$

The claim holds by setting $\mathcal{P}_2 := \mathcal{P}'_2 \cup \mathcal{P}''_2$. \square

Finally, we show how to prove the theorem using Claim 6.11. Note that \mathcal{P}_2 is a partial decomposition of D by Proposition 4.1(b). Let

$$\begin{aligned} D'' &:= D - \cup \mathcal{P}_2 = \cup(\mathcal{P}_Z \cup \mathcal{P}_1) \cup D_3, \\ D' &:= D'' - X = \cup(\mathcal{Q}_Z \cup \mathcal{P}_1) \cup (D_3 - X). \end{aligned}$$

Since vertices of X are isolated in D_3 , we have $E(D'' - D') = \{x_z z : z \in Z\}$. Therefore, by Proposition 4.2(b), (with $(D'', X, \emptyset, V(D) \setminus X)$ playing the roles of (D, A^+, A^-, R)) we see that if D' has a perfect decomposition, then so does D'' and hence, so does D , a contradiction; hence, D' does not have a perfect decomposition, proving (i). Note that (ii) follows from (a₁). Since $E(D_3 - X) = E(D_3)$, (iii) holds by Claim 6.11. For all $z \in Z$, we have $\text{ex}_{\mathcal{P}_1}(z) = 0$ and $\text{ex}_{D_3 - X}(z) = 0$, and so $\text{ex}_{D'}(z) = \text{ex}_{\cup \mathcal{Q}_Z}(z) = 1$ by Claim 6.9. Since \mathcal{P}_2 is a partial decomposition of D , $\text{ex}_{D'}^\pm(u) \leq \alpha n$ for all $u \in U^\pm$. Moreover, for $u \in U^\pm \setminus (X \cup Z)$, $\text{ex}_{D'}^\pm(u) = \text{ex}_{\cup \mathcal{P}_1}^\pm(u) \geq 1$. Hence, (iv) holds. Furthermore, we have $U^+(D') = (U^+ \setminus X) \cup Z$ and $U^-(D') = U^- \setminus Z$. Thus, (vi) holds by (a₄). Note that \mathcal{Q}_Z and \mathcal{P}_1 are partial decompositions[†] of D' , so $\text{ex}(D') \geq \text{ex}(D_3 - X) = \text{ex}(D_3) \geq \lambda Cn/32$ implying (v). \square

We now show that the digraph produced by Theorem 6.7 has a perfect decomposition. Together with Theorems 6.1 and 6.7, this will give us all the ingredients to prove Theorem 1.5.

THEOREM 6.12. *Let $1/n \ll \alpha, \lambda, \varepsilon \ll 1$. Suppose that D is an n -vertex oriented graph, where:*

- $\text{ex}(D) \geq 2n$;
- $\delta(D) \geq (1 - \varepsilon)n$;
- $1 \leq |\text{ex}_D(v)| \leq \alpha n$ for all $v \in V(D)$;
- $||U^-(D)| - |U^+(D)|| \leq 2\lambda|D|$.

Then D has a perfect decomposition.

Proof. Fix a parameter ε' satisfying $1/n \ll \alpha, \lambda, \varepsilon \ll \varepsilon' \ll 1$ such that $\varepsilon'n$ is an integer. Let

$$d = (1/2 - 52\varepsilon')n \text{ and } t := n - 2d = 104\varepsilon'n.$$

Arbitrarily partition $V(D)$ into X^+, X^-, X^0 such that

$$|X^+| = |X^-| = d, \quad |X^0| = t, \quad X^\pm \subseteq U^\pm := U^\pm(D).$$

(Note that such partition exists as $|U^\pm| \geq d$.) Our goal is to remove a partial decomposition \mathcal{P} of D such that the resulting digraph $D' := D - \cup \mathcal{P}$ satisfies

$$\text{ex}_{D'}(v) = \begin{cases} 1 & \text{if } v \in X^+; \\ 0 & \text{if } v \in X^0; \\ -1 & \text{if } v \in X^-; \end{cases} \text{ and } d_{D'}(v) = \begin{cases} 2d - 1 & \text{if } v \in X^+; \\ 2d & \text{if } v \in X^0; \\ 2d - 1 & \text{if } v \in X^-. \end{cases} \quad (6.3)$$

Then D' has a perfect decomposition \mathcal{P}' by Theorem 4.7 and so $\mathcal{P} \cup \mathcal{P}'$ is a perfect decomposition of D (by Proposition 4.1(b)). Thus, it remains to find such a \mathcal{P} .

We will construct \mathcal{P} as a union of three partial decompositions $\mathcal{P}_1, \mathcal{P}_2, \mathcal{P}_3$. Let $D_0 := D$ and write $D_i := D_{i-1} - \cup \mathcal{P}_i$ for $i = 1, 2, 3$. First, we reserve two multisets A_2 and A_3 , which will be sets of starting and ending points of \mathcal{P}_2 and \mathcal{P}_3 , respectively. Second, we find a partial decomposition \mathcal{P}_1 such that $\text{ex}_{D_1}(v)$ has the correct value provided $v \notin A_2 \cup A_3$ (see Claim 6.13). The partial decomposition \mathcal{P}_2 will ensure that $d_{D_2}(v) = 2d' - I_{X^+ \cup X^-}(v)$ for some $d' > d$. Finally, we adjust d' to d using \mathcal{P}_3 .

Since $\text{ex}(D) \geq 2n$ and $|\text{ex}_D(v)| \leq \alpha n$, we know we can find vertices $x_1, \dots, x_{26\varepsilon'n}, x'_1, \dots, x'_{26\varepsilon'n} \in U^+$ such that $x_i \neq x'_i$ and no vertex $v \in U^+$ is chosen more than

[†]To see this note that \mathcal{P}_Z and \mathcal{P}_1 are partial decompositions of D'' . We obtain, respectively, $D', \mathcal{Q}_Z, \mathcal{P}_1$ by deleting X from $D'', \mathcal{P}_Z, \mathcal{P}_1$. Then noting that $\text{ex}_{D''}(z) = \text{ex}_{\cup \mathcal{P}_Z \cup \mathcal{P}_1}(z) = 0$ for all $z \in Z$ and that the only edges incident to X in D'' are the initial edges of paths in \mathcal{P}_Z , we can conclude \mathcal{Q}_Z and \mathcal{P}_1 are partial decomposition of D' .

$(\text{ex}_D(v) - 1)/2$ times. Similarly, we are able to pick vertices $y_1, \dots, y_{26\varepsilon'n}, y'_1, \dots, y'_{26\varepsilon'n} \in U^-$ such that $y_i \neq y'_i$ and no vertex $v \in U^-$ is chosen more than $(|\text{ex}_D(v)| - 1)/2$ times. Clearly, x_i, x'_i, y_i, y'_i are distinct for all i . Let

$$A_2 := \{x_i, x'_i, y_i, y'_i : i \in [25\varepsilon'n]\},$$

$$A_3 := \{x_i, x'_i, y_i, y'_i : i \in [25\varepsilon'n + 1, 26\varepsilon'n]\}.$$

For $j \in \{2, 3\}$, let $\phi_j^+(v)$ (and $\phi_j^-(v)$) be the number of times that v is chosen as x_i or x'_i (and y_i or y'_i) in A_j . Let $\phi_j(v) := \phi_j^+(v) - \phi_j^-(v)$. Note that $\sum_{v \in V(D)} \phi_j(v) = 0$ and $2|\phi_2(v) + \phi_3(v)| < |\text{ex}_D(v)|$.

CLAIM 6.13. *There exists a partial decomposition \mathcal{P}_1 of D such that, writing $D_1 := D - \cup \mathcal{P}_1$, we have $\delta(D_1) \geq (1 - \varepsilon')n$ and for all $v \in V(D)$,*

$$\text{ex}_{D_1}(v) = \begin{cases} 2\phi_2(v) + 2\phi_3(v) + \pm 1 & \text{if } v \in X^\pm; \\ 2\phi_2(v) + 2\phi_3(v) & \text{if } v \in X^0. \end{cases}$$

Proof of Claim. Let $f : V(D) \rightarrow [n]$ be such that

$$f(v) = \begin{cases} \text{ex}(v) - 2\phi_2(v) - 2\phi_3(v) \mp 1 & \text{if } v \in X^\pm; \\ \text{ex}(v) - 2\phi_2(v) - 2\phi_3(v) & \text{if } v \in X^0. \end{cases}$$

Note that $\sum_{v \in V(D)} f(v) = 0$ and $|f(v)| \leq \alpha n$ for all $v \in V(D)$. Define a directed multigraph H on $V(D)$ such that $d_H^+(v) = \max\{f(v), 0\}$ and $d_H^-(v) = \max\{-f(v), 0\}$. Note that $\Delta(H) \leq \alpha n$. We apply Corollary 4.10 with D, H, α playing the roles of D, H, γ and obtain the desired partial decomposition \mathcal{P}_1 . \square

Let

$$s := \max_{v \in V} \{d_{D_1}(v) + I_{X^+ \cup X^-}(v)\}.$$

Note that $d_{D_1}(v)$ is even if $v \in X^0$ and odd otherwise. So, s is even and $(1 - \varepsilon')n \leq s \leq n$. Let $d' := s/2 - 50\varepsilon'n$, so

$$d + \varepsilon'n = (1/2 - 51\varepsilon')n \leq d' \leq (1/2 - 50\varepsilon')n = d + 2\varepsilon'n. \tag{6.4}$$

CLAIM 6.14. *There exists a partial decomposition \mathcal{P}_2 of D_1 such that, for all $v \in V(D)$, $\text{ex}_{\cup \mathcal{P}_2}(v) = 2\phi_2(v)$ and $d_{D_2}(v) = 2d' - I_{X^+ \cup X^-}(v)$, where we write $D_2 := D_1 - \cup \mathcal{P}_2$.*

Proof of Claim. Define $f : V(D) \rightarrow [\varepsilon'n]$ to be such that

$$f(v) = \varepsilon'n - \frac{s - (d_{D_1}(v) + I_{X^+ \cup X^-}(v))}{2}.$$

Note that $\max_{v \in V(D)} f(v) = \varepsilon'n$. Recall that $A_2 = \{x_i, x'_i, y_i, y'_i : i \in [25\varepsilon'n]\}$. Write $(x_i^*, x_{25\varepsilon'n+i}^*, y_i^*, y_{25\varepsilon'n+i}^*) = (x_i, x'_i, y_i, y'_i)$.

By Proposition 6.3, where we take $(V(D), 25, \varepsilon'n, x_i^*, y_i^*)$ to play the roles of (V, t, m, x_i, y_i) , we can find a collection of sets $T_1, \dots, T_{50\varepsilon'n} \subseteq V(D)$:

- (i') for all $v \in V(D)$, $\sum_{i \in [50\varepsilon'n]} I_{T_i}(v) = f(v) + 49\varepsilon'n$;
- (ii') $|T_i| \geq 24n/25$ for all $i \in [50\varepsilon'n]$;
- (iii') $x_i^*, y_i^* \in T_i$ for all $i \in [50\varepsilon'n]$.

For $i \in [50\varepsilon'n]$, let $S_i := V(D) \setminus T_i$ and H_i be the multidigraph on $V(D)$ with $E(H_i) = \{x_i^*y_i^*, x_i^*y_i^*\}$. Let $H = \bigcup_{i \in [50\varepsilon'n]} H_i$. Note that $|E(H_i)| = \Delta(H_i) = 2$ and $|S_i| \leq n/25$. We apply Lemma 4.9 where we take $(D_1, H_i, S_i, 50\varepsilon')$ to play the role of (D, H_i, S_i, γ) and obtain edge-disjoint paths $P_1, P'_1, \dots, P_{50\varepsilon'n}, P'_{50\varepsilon'n}$ such that both P_i and P'_i start at x_i^* and end at y_i^* and $d_{P_i \cup P'_i}(v) = 2I_{T_i}(v)$ for all $v \in V(D)$.

Set $\mathcal{P}_2 = \{P_i, P'_i : i \in [50\varepsilon'n]\}$. Note that $\text{ex}_{\cup \mathcal{P}_2}(v) = 2\phi_2(v)$ for all $v \in V(D)$. For all $v \in V(D)$,

$$d_{D_2}(v) = d_{D_1}(v) - 2 \sum_{i \in [50\varepsilon'n]} I_{T_i}(v) \stackrel{(i')}{=} d_{D_1}(v) - 2(f(v) + 49\varepsilon'n)$$

$$= s - 100\varepsilon'n - I_{X+\cup X^-}(v) = 2d' - I_{X+\cup X^-}(v),$$

as required. □

CLAIM 6.15. *There exists a partial decomposition \mathcal{P}_3 of D_2 such that for all $v \in V(D)$, $\text{ex}_{\cup \mathcal{P}_3}(v) = 2\phi_3(v)$ and $d_{D_3}(v) = 2d - I_{X+\cup X^-}(v)$, where we write $D_3 := D_2 - \cup \mathcal{P}_3$.*

Proof of Claim. Recall that $A_3 := \{x_i, x'_i, y_i, y'_i : i \in [25\varepsilon'n + 1, 26\varepsilon'n]\}$. Let $m := d' - d - \varepsilon'n$, so $0 \leq m \leq \varepsilon'n$ by (6.4).

We now define multidigraphs $H_1, \dots, H_{m+\varepsilon'n}$ on $V(D)$ as follows. Define $f(i) = 25\varepsilon'n + i$. For $i \in [m]$,

$$E(H_i) := \{x_{f(i)}y_{f(i)}, x_{f(i)}y_{f(i)}\}, \quad E(H_{\varepsilon'n+i}) := \{x'_{f(i)}y'_{f(i)}, x'_{f(i)}y'_{f(i)}\}.$$

For $i \in [m + 1, \varepsilon'n]$, set

$$E(H_i) := \{x_{f(i)}y_{f(i)}, x_{f(i)}y_{f(i)}, x'_{f(i)}y'_{f(i)}, x'_{f(i)}y'_{f(i)}\}.$$

Note that $|E(H_i)| \leq 4$ and $\Delta(H_i) = 2$. Let $H := \bigcup_{i \in [\varepsilon'n+m]} H_i$. We apply Lemma 4.9 with $(D_2, H_i, \emptyset, 2\varepsilon')$ playing the roles of (D, H_i, S_i, γ) to obtain a set of edge-disjoint paths $\mathcal{P}_3 = \{P_e : e \in E(H)\}$ such that $d_{\cup \mathcal{P}_3}(v) = 2(\varepsilon'n + m) = 2(d' - d)$ and $\text{ex}_{\cup \mathcal{P}_3}(v) = 2\phi_3(v)$ for all $v \in V(D)$. Note that \mathcal{P}_3 is a partial decomposition of D_2 by our choice of H_i and that $D_3 = D_2 - \cup \mathcal{P}_3$ satisfies the desired properties. □

For all $v \in V(D)$,

$$\text{ex}_{D_3}(v) = \text{ex}_{D_1}(v) - \text{ex}_{\cup \mathcal{P}_2}(v) - \text{ex}_{\cup \mathcal{P}_3}(v) = \begin{cases} \pm 1 & \text{if } v \in X^\pm; \\ 0 & \text{if } v \in X^0. \end{cases}$$

Also $d_{D_3}(v) = 2d - I_{X+\cup X^-}(v)$ for all $v \in V(D)$ by Claim 6.15. We are done by setting $D' := D_3$. □

6.3. Final proof

Now, we can finally prove our main theorem.

Proof of Theorem 1.5. Assume $1/n_0 \ll 1/C \ll 1$ and that T is an even tournament with $n \geq n_0$ vertices and $\text{ex}(T) \geq Cn$.

We pick parameters $\alpha_1, \beta_1, \varepsilon_1, \alpha_2, \beta_2, \varepsilon_2, \lambda$ satisfying:

$$1/n \ll 1/C \ll \beta_1 \ll \alpha_1, \beta_2 \ll \varepsilon_1 \ll \varepsilon_2, \lambda \ll 1 \quad \text{and} \quad 1/C \ll \lambda. \tag{6.5}$$

By Theorem 6.1, either T has a perfect decomposition or there is a digraph D_1 satisfying the following properties.

- (a₁) If D_1 has a perfect decomposition, then T has a perfect decomposition.
- (a₂) $n_1 := |D_1| \geq (1 - \beta_1)n$ with n_1 even.
- (a₃) $\delta(D_1) \geq (1 - \varepsilon_1)n_1$.
- (a₄) $1 \leq |\text{ex}_{D_1}(v)| \leq \alpha_1 n_1$ for all $v \in V(D_1)$.
- (a₅) $\text{ex}(D_1) \geq (C/4 - 5)n_1 =: C_1 n_1$.

By Theorem 6.7, there exists a digraph D_2 satisfying the following properties.

- (b₁) If D_2 has a perfect decomposition, then D_1 has a perfect decomposition.
- (b₂) $n_2 := |D_2| \geq (1 - \beta_2)n_1$ with n_2 even.
- (b₃) $\delta(D_2) \geq (1 - \varepsilon_2)n_2$.
- (b₄) $1 \leq |\text{ex}_{D_2}(v)| \leq \alpha_1 n_1 \leq 2\alpha_1 n_2$ for all $v \in V(D_2)$.
- (b₅) $\text{ex}(D_2) \geq \lambda C_1 n_2 / 32 =: C_2 n_2 \geq 2n_2$.
- (b₆) $||U^-(D_2)| - |U^+(D_2)|| \leq 2\lambda n_2$.

Note that by (6.5), we have $1/n_2 \ll 2\alpha_1, \lambda, \varepsilon_2 \ll 1$ since $n_2 \geq (1 - \beta_1)(1 - \beta_2)n \geq n/2$. By Theorem 6.12, D_2 has a perfect decomposition; hence, so does D_1 (by (b₁)) and so does T (by (a₁)) as required. \square

7. Conclusion

We have proved many cases of Conjecture 1.1. The obvious open problem remaining is to fill the remaining gap, that is to prove that $\text{pn}(T) = \text{ex}(T)$ for all even tournaments satisfying $n/2 < \text{ex}(T) \leq Cn$ for some sufficiently large C . We believe that with a little work, one should be able to apply the results of Kühn and Osthus [9] to prove the conjecture when $\text{ex}(T)$ is very close to $n/2$ but that probably some new ideas are needed, say when $n \leq \text{ex}(T) \leq Cn$.

Another direction, which is currently work in progress, is to investigate analogues of Conjecture 1.1 for directed graphs that are not tournaments. In forthcoming work, we consider dense directed graphs as well as random and quasi-random directed graphs.

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