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Ancestral Gumbel-Top-$k$ Sampling
for Sampling Without Replacement

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Abstract
We develop ancestral Gumbel-Top-$k$ sampling: a generic and efficient method for sampling without replacement from discrete-valued Bayesian networks, which includes multivariate discrete distributions, Markov chains and sequence models. The method uses an extension of the Gumbel-Max trick to sample without replacement by finding the top $k$ of perturbed log-probabilities among all possible configurations of a Bayesian network.

Despite the exponentially large domain, the algorithm has a complexity linear in the number of variables and sample size $k$. Our algorithm allows to set the number of parallel processors $m$, to trade off the number of iterations versus the total cost (iterations times $m$) of running the algorithm. For $m = 1$ the algorithm has minimum total cost, whereas for $m = k$ the number of iterations is minimized, and the resulting algorithm is known as Stochastic Beam Search. We provide extensions of the algorithm and discuss a number of related algorithms.

We analyze the properties of Gumbel-Top-$k$ sampling and compare against alternatives on randomly generated Bayesian networks with different levels of connectivity. In the context of (deep) sequence models, we show its use as a method to generate diverse but high-quality translations and statistical estimates of translation quality and entropy.

Keywords: sampling without replacement, ancestral sampling, bayesian networks, stochastic beam search, gumbel-max trick

1. This paper is an extended version of Kool et al. (2019b) which introduced Stochastic Beam Search.

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1. Introduction

Sampling from graphical models is a widely studied problem in machine learning. In many applications, such as neural machine translation, one may desire to obtain multiple samples, but wish to avoid duplicates, i.e., sample without replacement. In general, this is non-trivial: for example rejection sampling may take long if entropy is low.

In this paper, we extend the idea of sampling through optimization (Papandreou and Yuille, 2011; Hazan and Jaakkola, 2012; Tarlow et al., 2012; Hazan et al., 2013; Ermon et al., 2013; Maddison et al., 2014; Chen and Ghahramani, 2016; Balog et al., 2017) to generate multiple samples without replacement. The most well-known example of sampling by optimization is the Gumbel-Max trick (Gumbel, 1954; Maddison et al., 2014), which samples from the categorical distribution by optimizing (i.e. taking the argmax of) log-probabilities perturbed with independent Gumbel noise. Whereas the Gumbel-Max trick only considers the argmax (top 1), taking the top $k$ of perturbed log-probabilities actually results in a sample without replacement (Yellott, 1977; Vieira, 2014).

We consider sampling from discrete-valued Bayesian networks (since sampling without replacement from continuous domains is trivial), which means that we sample from a discrete multivariate distribution which is represented by a probabilistic directed acyclic graphical model (for examples, see Figure 1). By treating each possible configuration of the variables in the network as a category in a single (flat) categorical distribution, we can use ‘Gumbel-Top-$k$ sampling’ to sample $k$ configurations without replacement. To efficiently sample from the exponentially large domain, we use a top-down sampling procedure (Maddison et al., 2014) combined with an efficient branch-and-bound algorithm, which runs in time linear in the number of variables and number of samples.

The algorithm presented in this paper, which we refer to as ancestral Gumbel-Top-$k$ sampling, is a generalization of Stochastic Beam Search (originally presented in Kool et al. 2019b;), which allows to trade off ‘parallelizability’ against total required computation and is applicable to general Bayesian networks. As such, it serves the same purposes and can be used to generate representative and unique sequences from sequence models, for example for task such as neural machine translation (Sutskever et al., 2014; Bahdanau et al., 2015) and image captioning (Vinyals et al., 2015), where the diversity can be controlled by the sampling temperature. Additionally, being a sampling method, it can be used to construct statistical estimators as we show in Section 5. The ability to sample without replacement enables the construction of lower variance gradient estimators (Kool et al. 2019a, 2020).

Figure 1: Examples of Bayesian networks.

(a) Independent
\[ p(y) = \prod_v p(y_v) \]

(b) Markov chain
\[ p(y) = p(y_1) \prod_{t>1} p(y_t | y_{t-1}) \]

(c) General network
\[ p(y) = \prod_v p(y_v | y_{pa(v)}) \]

(d) Sequence model
\[ p(y) = \prod_t p(y_t | y_{1:t-1}) \]
2. Preliminaries

This section introduces Bayesian Networks, Deep Learning and the Gumbel-Max trick.

2.1. Bayesian Networks

A Bayesian network, also known as belief network, is a probabilistic directed acyclic graphical model that represents a joint probability distribution over a set of variables, which are nodes in a directed acyclic graph. We index the variables by $v \in \mathcal{V}$, where a node $v$ can take values $y_v \in D_v$. For an arbitrary subset $S \subseteq \mathcal{V}$, we write the corresponding set of values as $y_S = (y_v : v \in S)$, with domain $D_S = \prod_{v \in S} D_v$. For the complete network we write $y = y_{\mathcal{V}}$ with domain $D = D_{\mathcal{V}}$. The probability distribution for $y_v$ is defined conditionally on the parents $\text{pa}(v) \subseteq \mathcal{V} \setminus \{v\}$, with values $y_{\text{pa}(v)}$. This way, the distribution $p(y)$ is given by

$$p(y) = \prod_{v \in \mathcal{V}} p(y_v | y_{\text{pa}(v)}) \quad (1)$$

Any directed acyclic graph has at least one topological order (see e.g. Kahn 1962), which is an ordering in which each node is preceded by its parent nodes. For an example of a Bayesian network in topological order, see Figure 1c. Examples of Bayesian networks include multivariate categorical distributions where $y_v, v \in \mathcal{V}$ are independent (Figure 1a), Markov chains where $y_t$ depends only on $y_{t-1}$ (assuming $\mathcal{V} = \{1, ..., T\}$), and (finite length) sequence models where $y_t$ depends on the complete ‘prefix’ $y_{1:t-1}$ and the topological order is natural.

To keep notation compact, in general, we will not make the distinction between variables and their realizations. As each variable $y_v$ has a finite domain $D_v$, there is a finite number of possible realizations, or configurations, for the complete Bayesian network, specified by the domain $D = \prod_{v \in \mathcal{V}} D_v$. A configuration $y \in D$ has probability given by (1) and therefore, ignoring the graphical structure, we can treat $y$ as a ‘flat’ categorical variable over the domain $D$, where each category corresponds to a possible configuration $y$.

2.2. Deep Learning

Our focus is on modern Deep Learning (LeCun et al., 2015) applications, especially sampling from models represented as (discrete-valued) Stochastic Computation Graphs (Schulman et al., 2015), which can be considered Bayesian networks. Such models specify conditional distributions for variables using neural networks, dependent on parameters $\theta$ and an input or context $x$. In a discrete setting, such models usually output probabilities for all $y_v \in D_{\mathcal{V}}$ in a single model evaluation, by computing a softmax (with temperature $\tau \geq 0$) over unnormalized log-probabilities $\phi_{\theta}(y_v | y_{\text{pa}(v)}, x)$:

$$p_{\theta}(y_v | y_{\text{pa}(v)}, x) = \frac{\exp \left( \phi_{\theta}(y_v | y_{\text{pa}(v)}) / \tau \right)}{\sum_{y_{v} \in D_v} \exp \left( \phi_{\theta}(y_{v} | y_{\text{pa}(v)}) / \tau \right)} \quad \forall y_v \in D_v. \quad (2)$$

In Deep Learning, model evaluations involve millions of computations, which is why we seek to minimize them. Additionally, to make efficient use of modern hardware, algorithms should be efficiently parallelizable. In the remainder of this paper, we are concerned with sampling (without replacement) given fixed values of $\theta$ and $x$ so we will omit these in the notation.
2.3. The Gumbel-Max Trick

Treating the Bayesian network as specifying a categorical distribution, we can use the Gumbel-Max trick (Gumbel, 1954; Maddison et al., 2014) to sample from it by finding the configuration $\mathbf{y}$ with the largest perturbed log-probability. We define $\phi_\mathbf{y}$ as the log-probability of $\mathbf{y} \in \mathcal{D}$:

$$\phi_\mathbf{y} = \log p(\mathbf{y}) = \sum_{\mathbf{v} \in \mathcal{V}} \log p(y_v | \mathbf{pa}(v)).$$  

(3)

Next we define $G_{\phi_\mathbf{y}}$ as the perturbed log-probability of $\mathbf{y}$, which is obtained by adding (independent!) Gumbel noise $G_{\mathbf{y}} \sim \text{Gumbel}(0)$ to $\phi_\mathbf{y}$, where Gumbel($\phi$) is the Gumbel distribution with CDF

$$F_\phi(g) = \exp(-\exp(\phi-g)).$$  

(4)

Using inverse transform sampling, this noise is generated as $G_{\phi_\mathbf{y}} := G_{\mathbf{y}} + \phi_\mathbf{y} = \phi_\mathbf{y} - \log(-\log U_\mathbf{y}) \sim \text{Gumbel}(\phi_\mathbf{y})$.  

(5)

For any subset $B \subseteq \mathcal{D}$ it holds that (Maddison et al., 2014)

$$\max_{\mathbf{y} \in B} G_{\phi_\mathbf{y}} \sim \text{Gumbel}\left(\log \sum_{\mathbf{y} \in B} \exp \phi_\mathbf{y}\right),$$  

(6)

$$\arg \max_{\mathbf{y} \in B} G_{\phi_\mathbf{y}} \sim \text{Categorical}\left(\frac{\exp \phi_\mathbf{y}}{\sum_{\mathbf{y}' \in B} \exp \phi_{\mathbf{y}'}}; \mathbf{y} \in B\right).$$  

(7)

Equation (6) is a useful property that states that the maximum of a set of Gumbel variables is a Gumbel variable with as location the LOGSUMEXP of the individual Gumbel locations. Equation (7) states the most important result: the configuration $\mathbf{y}$ corresponding to the largest perturbed log-probability is a sample from the desired categorical distribution (since $\exp \phi_\mathbf{y} = p(\mathbf{y})$). Additionally, the max (6) and argmax (7) are independent, which is an important property that is used in this paper. For details, see Maddison et al. (2014). Figure 2 gives a visual illustration of the Gumbel-Max trick.

![Gumbel-Max Trick Diagram](image)

Figure 2: The Gumbel-Max trick: the argmax of perturbed log-probabilities has a categorical distribution. The maximum has an independent Gumbel distribution.
2.4. Gumbel-Top-k Sampling

An extension of the Gumbel-Max trick can be used to sample from the categorical distribution without replacement (Yellott, 1977; Vieira, 2014). To this end, let $y_1^*, ..., y_k^* = \arg \text{top}_k G_{\phi y}$, i.e. $y_1^*, ..., y_k^*$ are the configurations with the $k$ largest perturbed log-probabilities in decreasing order (see Figure 3). Denoting with $D^*_j = D\{y_1^*, ..., y_{j-1}^*\}$ the domain (without replacement) for the $j$-th sampled element, the probability for this ordered sample without replacement is given by

$$p(y_1^*, ..., y_k^*) = \prod_{j=1}^{k} p(y_j^* | y_1^*, ..., y_{j-1}^*)$$

$$= \prod_{j=1}^{k} P\left(y_j^* = \arg \max_{y \in D^*_j} G_{\phi y} \mid \max_{y \in D^*_j} G_{\phi y} < G_{\phi y_{j-1}^*}\right)$$

$$= \prod_{j=1}^{k} P\left(y_j^* = \arg \max_{y \in D^*_j} G_{\phi y}\right)$$

$$= \prod_{j=1}^{k} \frac{\exp \phi_{y_j^*}}{\sum_{y \in D^*_j} \exp \phi_y}$$

$$= \prod_{j=1}^{k} \frac{p(y_j^*)}{1 - \sum_{i=1}^{j-1} p(y_i^*)}. \quad (12)$$

To understand the derivation, note that conditioning on $y_1^*, ..., y_{j-1}^*$ means that $y_1^*, ..., y_{j-1}^*$ are the configurations with the $j - 1$ largest perturbed log-probabilities, so $y_j^*$, the configuration with the $j$-th largest perturbed log-probability, must be the arg max of the remaining log-probabilities. Additionally, we know that $\max_{y \in D^*_j} G_{\phi y}$, the maximum of the remaining log-probabilities, must be smaller than $G_{\phi y_{j-1}^*}$, which is the smallest of the $j - 1$ largest perturbed log-probabilities. This allows us to rewrite (8) as (9). The step from (9) to (10) follows from the independence of the max and arg max (Section 2.3) and the step from (10) to (11) uses the Gumbel-Max trick.

The form (11) is also known as the Plackett-Luce model (Plackett, 1975; Luce, 1959) and the form (12) highlights its interpretation as sampling without replacement, where the probabilities get renormalized after each sampled configuration. We refer to sampling without replacement by taking the top $k$ of Gumbel perturbed log-probabilities as Gumbel-Top-$k$ sampling. This is mathematically equivalent to weighted reservoir sampling (Efraimidis and Spirakis, 2006) which can be seen as a streaming implementation of Gumbel-Top-$k$ sampling.
3. Ancestral Gumbel-Top-k Sampling

Ancestral Gumbel-Top-k sampling is an efficient implementation of Gumbel-Top-k sampling for sampling from a probability distribution specified as a Bayesian network. It exploits the graphical structure to sample variables one at a time, conditionally on their parent variables, similar to ancestral sampling. The algorithm finds the top \( k \) configurations with largest perturbed log-probabilities (which determine the sample without replacement) implicitly, i.e. without sampling perturbed log-probabilities for all possible configurations of the Bayesian network. It does so by bounding the perturbed log-probabilities for parts of the domain, such that they can be pruned from the search if they are guaranteed to not contain a top \( k \) perturbed log-probability. Therefore, it can be considered a branch-and-bound algorithm (see e.g. Lawler and Wood 1966).

To derive ancestral Gumbel-Top-k sampling, we start with explicit Gumbel-Top-k sampling to sample \( k \) configurations \( y \) without replacement from the Bayesian network. This requires instantiating all configurations \( y \in D \), sampling their perturbed log-probabilities and finding the \( k \) largest to obtain the sample. See for an example Figure 4, where the leaf nodes represent all 8 possible configurations for a Bayesian network with 3 binary variables. As the size of the domain \( D \) is exponential in the number of variables \( |V| \), this is not practically feasible in general, and we will derive ancestral Gumbel-Top-k sampling as an alternative, equivalent, sampling procedure that does not require to instantiate the complete domain. Instead, it uses a top-down sampling procedure to sample perturbed log-probabilities for partial configurations \( y_S \) (internal nodes in Figure 4) first, which is equivalent but allows to obtain a bound on the perturbed log-probabilities of completions of \( y_S \) (descendants of \( y_S \) in the tree). This allows the dashed parts of the tree in Figure 4 to be pruned from the search, while obtaining the same result.

3.1. The Probability Tree Diagram

To help develop our theory, we will first assume a fixed topological order of the nodes \( V = \{1, ..., T\} \). For \( t \leq T \), we will use the notation \( y_{1:t} = y_{\{1,...,t\}} = (y_1, ..., y_t) \) to indicate a prefix of \( y \), which is a partial configuration of the Bayesian network with domain \( D_{1:t} = \prod_{v \in \{1,...,t\}} D_v \). This topological order allows us to represent possible configurations of the Bayesian network as the tree diagram in Figure 4, where level \( t \) has \(|D_{1:t}|\) nodes: one node for each possible partial configuration \( y_{1:t} \in D_{1:t} \). Given a partial configuration \( y_{1:t} \) of the nodes \( \{1, ..., t\} \), let \( S \) be a superset of those nodes (i.e. \( \{1, ..., t\} \subseteq S \)) and let \( D_{S|y_{1:t}} \) be the domain of configurations \( y_S \) given the partial assignment \( y_{1:t} \):

\[
D_{S|y_{1:t}} = \{y'_S \in D_S : y'_{1:t} = y_{1:t}\} = \prod_{v \in \{1,...,t\}} \{y_v\} \times D_S \setminus \{1,...,t\}.
\]

In particular, for \( S = \{1, ..., t + 1\} \), \( D_{1:t+1|y_{1:t}} = \prod_{v \in \{1,...,t\}} \{y_v\} \times D_{t+1} \) is the set of possible extensions of \( y_{1:t} \) by the variable \( y_{t+1} \), which defines the set of direct child nodes of \( y_{1:t} \) in the tree diagram. \( D_{1:t} = D_V|y_{1:t} \) is the set of possible complete configurations compatible with (or given) \( y_{1:t} \), which corresponds to the set of leaf nodes in the subtree rooted at \( y_{1:t} \) in the tree diagram.
We define the marginal probability for the partial configuration $y_{1:t}$ by marginalizing over all possible configurations $y \in D_{|y_{1:t}}$ that match the partial configuration:

$$p(y_{1:t}) = \sum_{y \in D_{|y_{1:t}}} p(y)$$

$$= \sum_{y \in D_{|y_{1:t}}} \prod_{v \in \{1, \ldots, t\}} p\left(y_v | y_{pa(v)}\right) \prod_{v \in \{t+1, \ldots, T\}} p\left(y_v | y_{pa(v)}\right)$$

$$= \prod_{v \in \{1, \ldots, t\}} p\left(y_v | y_{pa(v)}\right) \sum_{y \in D_{|y_{1:t}}} \prod_{v \in \{t+1, \ldots, T\}} p\left(y_v | y_{pa(v)}\right)$$

$$= \prod_{v \in \{1, \ldots, t\}} p\left(y_v | y_{pa(v)}\right)$$

$$= p(y_{1:t-1})p\left(y_t | y_{pa(t)}\right).$$

(14)

Note that $y_{1:t-1}$ is the direct parent of the node $y_{1:t}$ so (14) allows efficient computation of the probability $p(y_{1:t-1})$ by multiplying the conditional probability $p\left(y_t | y_{pa(t)}\right)$ (note that $pa(t) \subseteq \{1, \ldots, t-1\}$) with the marginal probability $p(y_{1:t-1})$ of the parent in the tree diagram.
3.2. Explicit Gumbel-Top-$k$ Sampling

By using Gumbel-Top-$k$ sampling (see Section 2.4) explicitly, we can (inefficiently!) obtain a sample without replacement of size $k$ from the Bayesian network as follows:

- Compute $\phi_y = \log p_\theta(y)$ for all configurations $y \in D$. This means that the complete tree diagram is instantiated, as in Figure 4.
- For $y \in D$, sample $G_{\phi_y} \sim \text{Gumbel}(\phi_y)$, so $G_{\phi_y}$ is the perturbed log-probability of $y$.
- Let $y_1^*, ..., y_k^* = \text{arg top}_k G_{\phi_y}$, then $y_1^*, ..., y_k^*$ is a sample of configurations from (1) without replacement.

Note that the result corresponds to a subset of leaf nodes in the tree diagram. As explicit Gumbel-Top-$k$ sampling requires to instantiate the complete tree diagram to compute perturbed log-probabilities for all leaf nodes, this is computationally prohibitive. Therefore, we construct an equivalent process that we call ancestral Gumbel-Top-$k$ sampling, which implicitly finds the configurations $y_1^*, ..., y_k^*$ with computation that is only linear in the number of samples $k$ and the number of variables $T$.

3.3. Perturbed Log-Probabilities of Partial Configurations

So far, we have only defined the perturbed log-probabilities $G_{\phi_y}$ for complete configurations $y \in D$, corresponding to leaf nodes in the tree diagram. To derive ancestral Gumbel-Top-$k$ sampling, we find it convenient to also define the perturbed log-probabilities for partial configurations $y_{1:t} \in D_{1:t}$, which correspond to internal nodes in the tree diagram. From Equation (13), it follows that a partial configuration $y_{1:t}$ has log-probability

$$\phi_{y_{1:t}} = \log p(y_{1:t}) = \log \sum_{y \in D_{1:t}} p(y) = \log \sum_{y \in D_{1:t}} \exp \phi_y.$$ 

For a partial configuration $y_{1:t}$, we now consider the maximum of the perturbed log-probabilities $G_{\phi_y}$ of compatible complete configurations $y \in D_{1:t}$. This corresponds to the maximum of perturbed log-probabilities in the subtree below $y_{1:t}$ in Figure 4. By the Gumbel-Max trick (Equation 6), it has a Gumbel distribution with location $\phi_{y_{1:t}}$. Therefore, we use the notation $G_{\phi_{y_{1:t}}}$:

$$G_{\phi_{y_{1:t}}} = \max_{y \in D_{1:t}} G_{\phi_y} \sim \text{Gumbel}(\phi_{y_{1:t}}).$$

(15)

Since $G_{\phi_{y_{1:t}}} \sim \text{Gumbel}(\phi_{y_{1:t}})$, we can rewrite it as

$$G_{\phi_{y_{1:t}}} = \phi_{y_{1:t}} + G_{y_{1:t}},$$

(16)

where we have defined $G_{y_{1:t}} \sim \text{Gumbel}(0)$. Equation (16) reveals that we can interpret $G_{\phi_{y_{1:t}}}$ as the perturbed log-probability of partial configuration $y_{1:t}$. Thus, the perturbed log-probability of a partial configuration $y_{1:t}$ is the maximum of perturbed log-probabilities of its possible completions $D_{1:t}$. This relation is key to the top-down sampling algorithm that we will derive.

Looking at the tree diagram in Figure 4, the maximum below a node $y_{1:t}$ must be attained in one of the subtrees, such that the perturbed log-probability of $y_{1:t}$ must be the maximum of perturbed log-probabilities of its possible one-variable extensions (children)
As an alternative to sampling the perturbed log-probabilities ‘bottom-up’, we can reverse \( \phi \) The CDF of a Gumbel(0), which is independent of the maximum (see Section 2.3). This means that we can sample its perturbed log-probability \( G \) we denote by \( y \) \( \phi \) perturbed \( \phi \) for the leaf nodes \( y \in D \) and computing Equation (17) recursively up the tree. We refer to this procedure as bottom-up sampling of the perturbed log-probabilities. Here we consider computing \( G_{\phi y_{1:t}} \) using Equation (17) as a sampling step because \( G_{\phi y_{1:t}} \) is a random variable which has a degenerate (constant) distribution by conditioning on the children.

### 3.4. Top-down Sampling of Perturbed Log-Probabilities

As an alternative to sampling the perturbed log-probabilities ‘bottom-up’, we can reverse the process. Starting from the root of the tree diagram with an empty configuration, which we denote by \( y_C \) with length \( t = 0 \), we conditionally sample the perturbed log-probabilities for the children recursively.

For the root \( y_C \) at level \( t = 0 \) it holds that \( D(y_C) = D \) and the log-probability is

\[
\phi_{y_C} = \log \sum_{y \in D} \exp \phi_y = \log \sum_{y \in D} p(y) = \log 1 = 0.
\]

This means that we can sample its perturbed log-probability \( G_{\phi y_C} \sim \text{Gumbel}(0) \), or we can simply set \( G_{\phi y_C} = 0 \), which conditions the sampling process on the event that \( \max_{y \in D} G_{\phi y} = 0 \), which does not affect the result since the sample is determined by the \( \arg \max \) / \( \arg \top k \), which is independent of the maximum (see Section 2.3).

Given the log-probability \( \phi_{y_{1:t}} \) of a node \( y_{1:t} \), we can efficiently compute the log-probabilities \( \phi_{y_{1:t+1}} \) for its children as \( \phi_{y_{1:t+1}} = \phi_{y_{1:t}} + \log p(y_t | y_{pa(t)}) \) (this follows from Equation 14). Similarly, given the perturbed log-probability \( G_{\phi y_{1:t}} \) and the log-probabilities \( \phi_{y_{1:t+1}} \), we can top-down sample the perturbed log-probabilities \( G_{\phi y_{1:t+1}} \) for each possible \( y_{t+1} \in D_{t+1} \). However, there is a dependence between \( G_{\phi y_{1:t}} \) and \( G_{\phi y_{1:t+1}} \) given by \( G_{\phi y_{1:t}} = \max_{y_{t+1} \in D_{t+1}} G_{\phi y_{1:t+1}} \) (Equation 17). Therefore, when sampling \( G_{\phi y_{1:t+1}} \), we need to make sure this dependency is satisfied, i.e. we need to sample a set of (independent) Gumbel variables conditionally upon their maximum.

### 3.5. Sampling a Set of Gumbel Variables Conditionally on Their Maximum

To sample a set of Gumbel variables \( G_{\phi_i} \sim \text{Gumbel}(\phi_i) \) with a given maximum \( T \), we can:

- Sample \( G_{\phi_i} \) for all \( i \) independently
- Let \( Z = \max_i G_{\phi_i} \) be the observed maximum
- Let \( \hat{G}_{\phi_i} = F^{-1}_{\phi_i}(F_{\phi_i,Z}(G_{\phi_i})) \), where \( F_{\phi_i,Z} \) and \( F^{-1}_{\phi_i,T} \) are the CDF and inverse CDF of truncated Gumbel Distributions.

The CDF of a Gumbel(\( \phi \)) distribution truncated at \( T \) is given by

\[
F_{\phi,T}(g) = P(G \leq g | G \leq T) = \frac{P(G \leq g \cap G \leq T)}{P(G \leq T)} = \frac{P(G \leq \min(g,T))}{P(G \leq T)} = \frac{F_\phi(\min(g,T))}{F_\phi(T)},
\]
where \( F_\phi(g) = \exp(-\exp(\phi - g)) \) is the CDF of the (untruncated) Gumbel(\( \phi \)) distribution (Equation 4). The inverse CDF of the truncated Gumbel distribution is given by

\[
F_{\phi,T}^{-1}(u) = \phi - \log(\exp(\phi - T) - \log u),
\]

such that the transformation \( \tilde{G}_{\phi_i} = F_{\phi_i,T}^{-1}(F_{\phi_i,Z}(G_{\phi_i})) \) can be written explicitly as

\[
\tilde{G}_{\phi_i} = F_{\phi_i,T}^{-1}(F_{\phi_i,Z}(G_{\phi_i})) = \phi_i - \log(\exp(\phi_i - T) - \exp(\phi_i - Z) + \exp(\phi_i - G_{\phi_i})) = -\log(\exp(-T) - \exp(-Z) + \exp(-G_{\phi_i})).
\]  

(18)

This shows that the samples with maximum \( Z \) are effectively ‘shifted’ towards their desired maximum \( T \), in (negative) exponential space through the (inverse) truncated Gumbel CDF. See Appendix A for a numerically stable implementation.

The validity of the sampling procedure is shown by conditioning on \( \arg \max_j G_{\phi_j} \).

For \( i = \arg \max_j G_{\phi_j} \) it holds that

\[
P(\tilde{G}_{\phi_i} \leq g | i = \arg \max_j G_{\phi_j}) = \mathbb{E}_Z \left[ P(\tilde{G}_{\phi_i} \leq g | i = \arg \max_j G_{\phi_j}, \max_j G_{\phi_j} = Z) \right] = \mathbb{E}_Z \left[ P(\tilde{G}_{\phi_i} \leq g | G_{\phi_i} = Z) \right] = \mathbb{E}_Z \left[ P(\tilde{G}_{\phi_i} \leq g | G_{\phi_i} = Z) \right] = \mathbb{E}_Z \left[ P(\tilde{G}_{\phi_i} \leq g | G_{\phi_i}, T) \right] = \mathbb{P}(T \leq g) = \mathbb{P}(G_{\phi_i} \leq g | \max_j G_{\phi_j} = T, i = \arg \max_j G_{\phi_j}).
\]

Combining the two cases, it holds that

\[
P(\tilde{G}_{\phi_i} \leq g) = P(\tilde{G}_{\phi_i} \leq g | i = \arg \max_j G_{\phi_j}) P(i = \arg \max_j G_{\phi_j}) + P(\tilde{G}_{\phi_i} \leq g | i \neq \arg \max_j G_{\phi_j}) P(i \neq \arg \max_j G_{\phi_j})
\]

\[
= P(G_{\phi_i} \leq g | \max_j G_{\phi_j} = T, i = \arg \max_j G_{\phi_j}) P(i = \arg \max_j G_{\phi_j}) + P(G_{\phi_i} \leq g | \max_j G_{\phi_j} = T, i \neq \arg \max_j G_{\phi_j}) P(i \neq \arg \max_j G_{\phi_j})
\]

\[
= P(G_{\phi_i} \leq g | \max_j G_{\phi_j} = T).
\]

This procedure allows us to recursively sample \( G_{\phi_1,T} \) for the complete tree diagram top-down, which is equivalent to sampling the perturbed log-probabilities for the complete tree diagram bottom-up. This means that for the leaves, or complete configurations \( y \), it holds that \( G_{\phi_y} \sim \text{Gumbel}(\phi_y) \) independently. The benefit of the top-down sampling procedure is that we can choose to sample only the parts of the tree diagram needed to obtain the top \( k \) leaves.
3.6. Ancestral Gumbel-Top-k Sampling

At the heart of the ancestral Gumbel-Top-k sampling is the idea that we can find the configurations corresponding to the top k perturbed log-probabilities, without instantiating the complete domain, by using top-down sampling. To do so, we maintain a queue $Q$ of partial configurations $y_{1:t}$, where no partial configuration $y_{1:t} \in Q$ is a prefix of another one. By doing so, each element $y_{1:t} \in Q$ can be seen as the root of a (disjoint) part of the domain $D|_{y_{1:t}} \subseteq D$. We do this in such a way that the queue represents a partitioning of the complete domain, so $D = \bigcup_{y_{1:t} \in Q} D|_{y_{1:t}}$.

For each of the partial configurations in the queue, we also keep track of the perturbed log-probability $G_{\phi y_{1:t}}$, which is obtained by top-down sampling. The queue represents the set of leaf nodes of a partially constructed probability tree diagram (see Figure 5). At any point in time, we can remove an element $y_{1:t}$ from the queue and add to the queue the set of extensions $y_{1:t+1}$ for $y_{t+1} \in D_{t+1}$, for which we can sample the perturbed log-probabilities $G_{\phi y_{1:t+1}}$ conditionally upon the value of their parent $G_{\phi y_{1:t}}$, as in standard ancestral sampling. Repeating this process will ultimately result in a queue $Q = D$, sampling all complete configurations $y \in D$, including their perturbed log-probabilities, such that the top $k$ determines the sample without replacement. To improve efficiency, we can bound the size of the queue, limiting the total number of expansions, without changing the result.

![Diagram of ancestral Gumbel-Top-k sampling](image)

Figure 5: Ancestral Gumbel-Top-k sampling, with $k = 3$ and $m = 1$. Configurations in the queue are shaded and in each iteration, the single ($m = 1$) incomplete configuration with the highest perturbed log-probability is expanded, sampling the perturbed log-probabilities for the extensions conditionally. This takes 5 iterations and has a total cost of 5 model evaluations to obtain $k = 3$ samples without replacement.
3.7. Bounding of the Queue Size at $k$

Assuming we are interested in sampling $k$ configurations without replacement, let $\kappa = \kappa(D)$ be the $k$-th largest of the perturbed log-probabilities $\{G_{\phi y} : y \in D\}$ of complete configurations $y \in D$. We call $\kappa$ the threshold since, using Gumbel-Top-$k$ sampling, $y$ will be part of the $k$ samples without replacement if $G_{\phi y} \geq \kappa$.

**Lemma 1** (Lower bound) Let $\kappa(Q)$ be the $k$-th largest perturbed log-probability of (partial) configurations in the queue $Q$. Then $\kappa(Q)$ is a lower bound for $\kappa$, i.e. $\kappa \geq \kappa(Q)$.

**Proof** For each of the $k$ largest perturbed log-probabilities $G_{\phi y_{1:t}}$ in the queue $Q$, it holds that $G_{\phi y_{1:t}} \geq \kappa(Q)$ (by definition). By the definition of $G_{\phi y_{1:t}}$ as the maximum of perturbed log-probabilities $G_{\phi y}$ for $y \in D_{\mid y_{1:t}}$ (Equation 15), there must necessarily be a completion $y \in D_{\mid y_{1:t}}$ for which $G_{\phi y} = G_{\phi y_{1:t}} \geq \kappa(Q)$. Since for $y_{1:t} \in Q$ the corresponding domains $D_{\mid y_{1:t}}$ are disjoint, this means that there are (at least) $k$ distinct configurations $y \in D$ such that $G_{\phi y} \geq \kappa(Q)$. Since there are at least $k$ configurations $y \in D$ for which $G_{\phi y} \geq \kappa(Q)$, and $\kappa$ is the $k$-th largest perturbed log-probability in $D$, it must hold that $\kappa \geq \kappa(Q)$.

**Lemma 2** (Upper bound) For $y_{1:t} \in Q$, the perturbed log-probability $G_{\phi y_{1:t}}$ is an upper bound for the perturbed log-probabilities of its possible completions, i.e. $G_{\phi y_{1:t}} \geq G_{\phi y}$ for $y \in D_{\mid y_{1:t}}$.

**Proof** It follows directly from the the definition of $G_{\phi y_{1:t}}$ (Equation 15) that

$$G_{\phi y_{1:t}} = \max_{y' \in D_{\mid y_{1:t}}} G_{\phi y'} \geq G_{\phi y} \quad \forall y \in D_{\mid y_{1:t}}.$$

**Theorem 3** (Limit queue size) All except the $k$ configurations with the largest perturbed log-probabilities in the queue $Q$ can be discarded while still resulting in a sample without replacement of size $k$.

**Proof** For $y_{1:t} \in Q$ that does not belong to one of the top $k$ perturbed log-probabilities, it holds that $G_{\phi y_{1:t}} < \kappa(Q)$ (assuming uniqueness of $G_{\phi y_{1:t}}$ since the domain is continuous) such that for its possible completions $y \in D_{\mid y_{1:t}}$ it holds that (using Lemma 1 and 2)

$$G_{\phi y} \leq G_{\phi y_{1:t}} < \kappa(Q) \leq \kappa \quad \forall y \in D_{\mid y_{1:t}}.$$

If $y_{1:t}$ is not one of the top $k$ in $Q$, then $G_{\phi y} < \kappa$ for all possible completions of $y_{1:t}$, and any such completion will not be in the sample without replacement of size $k$. This means that $y_{1:t}$ can be discarded from the queue without affecting the result.
As a consequence of Theorem 3, if we desire to sample $k$ configurations, then at any point in time, we need to keep at most the $k$ partial configurations with largest perturbed log-probabilities in the queue. This can be implemented by maintaining $Q$ as a priority queue of limited size $k$, where the priorities are given by the perturbed log-probabilities. Repeatedly removing the first incomplete configuration (with the highest perturbed log-probability) from the queue, and replacing it by its extensions, ultimately results in a queue with $k$ complete configurations, corresponding to the $k$ largest perturbed log-probabilities, as is illustrated in Figure 5. By its equivalence to explicit Gumbel-Top-$k$ sampling, this is a sample without replacement of size $k$ (in order!). By repeatedly expanding the partial configuration with highest perturbed log-probability, the algorithm is a best first search algorithm.

In a practical implementation, we can directly ‘yield’ the first complete configuration once it is found, remove it from the queue and decrease the queue size limit by 1. This will yield the first sample in exactly $T$ iterations (where $T$ is the number of variables), as if we used standard ancestral sampling. Then, it will ‘jump’ back to the partial configuration $y_{1:t}$ with the highest perturbed log-probability and repeat the best first search from there, lazily generating each next sample as the algorithm proceeds. By reusing partially sampled configurations, the algorithm evaluates the model for each possible partial configuration at most once. It is ‘anytime optimal’ in the sense that to yield $k' \leq k$ samples, it only computes model evaluations that are strictly necessary. The algorithm is a special case of Algorithm 1 with $m = 1$ (see Section 3.8) and a fixed variable selection strategy (Section 3.9).

Algorithm 1 AncestralGumbelTopkSampling($p_\theta, k, m$)

1: **Input:** Bayesian network $p_\theta(y_v|pa(v)) \forall v \in V$, max. sample size $k$, parallel expansion parameter $m$, variable selection strategy $\Upsilon$
2: Initialize result and queue empty
3: add $(S = \emptyset, y_S = \emptyset, \phi_{y_S} = 0, G_{\phi_{y_S}} = 0)$ to QUEUE
4: while QUEUE not empty do
5: expand $\leftarrow$ take and remove top $m$ from QUEUE according to $\tilde{G}$
6: for $(S, y_S, \phi_{y_S}, G_{\phi_{y_S}}) \in$ EXPAND (in parallel) do
7: if $S = V$ then
8: add $(S, y_S, \phi_{y_S}, G_{\phi_{y_S}})$ to QUEUE
9: else
10: select $v$ from $\{v \in V \setminus S : pa(v) \subseteq S\}$ according to variable selection strategy $\Upsilon$
11: let $S' = S \cup \{v\}$
12: compute $\phi_{y_{S'}} \leftarrow \phi_{y_S} + \log p_\theta(y_v|y_{pa(v)}) \quad \forall y_v \in D_v$ (here $y_{S'} = y_S \cup \{y_v\}$)
13: sample $G_{\phi_{y_{S'}}} \sim \text{Gumbel}(\phi_{y_{S'}}) \quad \forall y_v \in D_v$ conditionally given maximum $G_{\phi_{y_S}}$
14: add $(S', y_{S'}, \phi_{y_{S'}}, G_{\phi_{y_{S'}}})$ to QUEUE for all $y_v \in D_v$
15: end if
16: end for
17: QUEUE $\leftarrow$ take top $k$ from QUEUE according to $\tilde{G}$
18: while first element from QUEUE is complete, i.e. has $S = V$ do
19: remove first element from QUEUE and add it to RESULT
20: $k \leftarrow k - 1$
21: end while
22: end while
23: Return RESULT
3.8. Parallelizing the Algorithm

Instead of only expanding the first partial configuration in the queue, we can assume availability of $m \leq k$ parallel processors and expand $m$ partial configurations in parallel, replacing all of them by their (disjoint) expansions. This results in fewer required iterations (known as the span or depth) of the algorithm, but an increase of total cost of the algorithm, which is the number of iterations times $m$. This is because of the cost incurred for ‘idle’ parallel processors if initially the queue size is smaller than $m$, and because of redundant model evaluations if nodes are expanded which, in a fully sequential setting, would have been ‘pushed’ of the queue by other expansions (see Figure 6 for an example).

For $m = k$, the complete queue gets replaced in every iteration by the top $k$ of all expansions, and the resulting algorithm is a limited width breadth first search which was originally presented in Kool et al. (2019b) as Stochastic Beam Search. While this algorithm requires the fewest iterations, as it guarantees $k$ samples are generated in exactly $T$ iterations (where $T$ is the number of variables), it is the least efficient in terms of total cost.

Ultimately, the optimal choice of $m$ depends on properties of the model and available parallel hardware, e.g. the number of GPUs\(^3\) (or parallel processors on a single GPU) in modern Deep Learning applications. For example, if a single model evaluation can already fully utilize the GPUs (e.g. if it is a large and highly parallelizable neural network model), then setting $m = 1$ is optimal as this computes only strictly necessary model evaluations. On the other hand, if GPUs cannot be fully utilized for individual model evaluations, or additional GPUs are freely available, one should use $m = k$ to minimize the number of iterations which then directly translates to minimizing running time.

![Figure 6: Ancestral Gumbel-Top-k sampling, with $k = m$ (= 3), also known as Stochastic Beam Search (Kool et al., 2019b). This version expands in each iteration all partial configurations in the queue in parallel, traversing the tree diagram per level, i.e. using breadth first search. Compared to ancestral sampling with $m = 1$, this reduces the number of iterations from 5 to 3, but has a total cost of 3 iterations × 3 processors = 9. The total cost can be divided into a cost of 6 (rather than 5) for model evaluations, in this case since ‘CC’ is expanded unnecessarily, and a cost of 3 for idle processors if the queue (marked by shaded nodes) size is smaller than $m$: this cost is 2 in the first iteration and 1 in the second. The result is equivalent to using $m = 1$ (Figure 5) or any other $m < k$.](image-url)

\(^3\) GPU is an abbreviation of Graphical Processing Unit.
3.9. Variable Selection Strategy

We can consider a tree diagram as a tree representing the possible paths (from root to leaf) that standard ancestral sampling can take to generate a sample. In the tree diagram in Figure 4, each level $t$ corresponds to the fixed variable $y_t$, such that the sample is always generated in the order $y_1, ..., y_T$. However, if the graphical model has multiple topological orderings, any such order yields a valid tree diagram and thus valid order for ancestral sampling. Moreover, the order does not have to be globally fixed, but can vary depending on the partial configuration sampled so far.

In general, when ancestral sampling, we are free to choose any variable $y_t$ to sample next, as long as its parents $y_{pa(t)}$ have been sampled. For example, for a fully factorized model $p(y) = p(y_1)p(y_2)p(y_3)$, we are free to start sampling with $y_3$, and, depending on the value of $y_3$, select either $y_1$ or $y_2$ to sample next. This is an example of a variable selection strategy, which is represented by the tree diagram in Figure 7. While the resulting distribution is the same, the variable selection strategy affects which parts of the tree can be discarded (dashed nodes in Figure 7) and thus the number of iterations that are required by the algorithm. Sampling low-entropy variables first is a good idea, as this allows larger partial configurations to be ‘reused’ for subsequent samples (see Section 5.1.4). Algorithm 1 is presented in the form that can take an arbitrary variable selection strategy.

4. Related Algorithms

In this section we discuss a number of algorithms related to ancestral Gumbel-Top-$k$ sampling.

4.1. The ‘Trajectory Generator’

For $m = 1$, our ancestral Gumbel-Top-$k$ sampling algorithm is similar to the ‘trajectory generator’ with $\epsilon = 0$ described by Lorberbom et al. (2019), where it is used in the context of reinforcement learning (Sutton and Barto, 2018) to generate direct policy gradients. Our algorithm replaces a partial configuration $y_{1:t}$ in the queue by all expansions $y_{1:t+1}$, motivated by the idea that this is efficient if we consider computing $p(y_{1:t+1}|y_{1:t})$ for all $y_{t+1}$ to be a single model evaluation, e.g. the forward pass of a neural network with a softmax output layer (Equation 2). In the probability tree diagram, this means that if node is expanded, all childnodes are added to the tree. In the context of reinforcement learning, the algorithm by Lorberbom et al. (2019) is assumed to operate on a ‘state-reward tree’, where generating a new node corresponds to taking an action which requires an interaction with the environment. Therefore, they only add a single child node to the tree in each iteration.
In the implementation by Lorberbom et al. (2019), elements in the queue do not (only) represent leaf nodes of the tree, but the idea that they form a partition of the domain is still maintained. In particular, in our implementation, upon expansion we partition the domain \( D_{|y_1:t} \) represented by \( y_1:t \) as 
\[
D_{|y_1:t} = \bigcup_{y_{t+1} \in D_t} D_{|y_1:t+1}, \quad \text{whereas Lorberbom et al. (2019) expand only a single child node and partition the domain as}
\]
\[
D_{|y_1:t} = D_{|y_1:t+1} \cup (D_{|y_1:t} \setminus D_{|y_1:t+1}).
\]
When a single child node \( y_{t+1} \) is created, this corresponds to the maximum and thus it inherits the perturbed log-probability from the parent. For the parent, a new perturbed log-probability is sampled, truncated at the previous maximum, which then becomes the maximum of perturbed log-probabilities for the remaining domain \( D_{|y_1:t} \setminus D_{|y_1:t+1} \) corresponding to the child nodes that have not yet been instantiated.

4.2. **(Stochastic) Beam Search**

Beam search is a widely used method for approximate inference in various domains such as machine translation (Sutskever et al., 2014; Bahdanau et al., 2015) and image captioning (Vinyals et al., 2015). In machine learning, beam search is typically a test-time procedure, but there are works that include beam search in the training loop (Daumé et al., 2009; Wiseman and Rush, 2016; Edunov et al. 2018b; Negrinho et al., 2018; Edunov et al. 2018a). Beam search suffers from limited diversity of the results, and variants have been developed that encourage diversity (Li et al., 2016; Shao et al. 2017; Vijayakumar et al., 2018).

We argue that adding stochasticity is also a principled way to increase diversity in a beam search: this motivated the development of Stochastic Beam Search (Kool et al., 2019b). While Stochastic Beam Search is equivalent to ancestral Gumbel-Top-\( k \) sampling (only) for \( m = k \), the result (a sample without replacement) is equivalent for any \( m \leq k \). As such, we also consider ancestral Gumbel-Top-\( k \) sampling as a principled alternative to a heuristically randomized/diversified beam search, even though it is not (in general) a beam search.

We analyze the result of ancestral Gumbel-Top-\( k \) sampling by comparing Stochastic Beam Search to a naïve alternative implementation of a randomized beam search. In particular, imagine that we use an ordinary beam search, but replace the deterministic top-\( k \) operation by sampling without replacement in each step, e.g. using Gumbel-Top-\( k \) sampling, but without top-down sampling conditionally. In this naïve approach, a low-probability partial configuration will only be extended to completion, if it gets to be re-sampled, independently, again with low probability, at each step during the beam search. The result is a much lower probability to sample this configuration than assigned by the model.

Intuitively, we should somehow commit to a sampling ‘decision’ made at step \( t \). However, a hard commitment to generate exactly one completion for each of the \( k \) partial configurations at step \( t \) would prevent generating any two completions from the same partial configuration. By using top-down sampling, Stochastic Beam Search propagates the Gumbel perturbation of a partial configuration to its extensions, which can be seen as a soft commitment to that partial configuration. This means that it gets extended as long as its total perturbed log-probability is among the top \( k \), but ‘falls off’ the beam if, despite the consistent perturbation, another partial configuration becomes more promising. By this procedure, the result is a sample without replacement, which is true to the model, suggesting that Stochastic Beam Search, or equivalently, ancestral Gumbel-Top-\( k \) sampling, is a principled alternative to a heuristically randomized beam search.
4.3. Threshold Sampling

While Gumbel-Top-k sampling uses a fixed sample size of k, a related algorithm is threshold sampling (Duffield et al., 2005), which instead sets a fixed threshold and returns a variable sized sample. In analogy with Gumbel-Top-k sampling, threshold sampling would sample all configurations in the domain for which the perturbed log-probability exceeds a fixed threshold, instead of the k largest. Threshold sampling is a special case of Poisson sampling as each configuration is included in the sample with probability independent of the other configurations. This allows to use the sample for statistical estimation using the Horvitz-Thompson estimator (Horvitz and Thompson, 1952). To overcome the limitation of a variable sample size, priority sampling (Duffield et al., 2007) uses the (k+1)-th largest value as empirical threshold to obtain a fixed sample size k. This is equivalent to Gumbel-Top-k sampling and we experiment with using the resulting estimator in Section 5.3. An alternative method to control the sample size is to use an adaptive threshold (Ting, 2017).

4.4. Rejection Sampling

As an alternative to our algorithm, we can draw samples without replacement by rejecting duplicates from samples drawn with replacement. However, if the domain is large and the entropy low (e.g. with a translation model where there are only a few valid translations), then rejection sampling requires many samples and consequently many (expensive) model evaluations. Also, we have to estimate how many samples to draw (in parallel) or draw samples sequentially. Our algorithm allows to set \( m = k \), which guarantees generating k unique samples with exactly \( kT \) model evaluations in a single pass of \( T \) steps, which is equal to the computational requirement for taking \( k \) samples with replacement. When allowing our algorithm to generate samples sequentially, setting \( m < k \), it can generate \( k \) samples with fewer model evaluations than standard sampling with replacement.

4.5. Naïve Ancestral Sampling Without Replacement

As an alternative to ancestral Gumbel-Top-k sampling, we can also derive an algorithm based on ‘standard’ sampling without replacement. This means sampling the first configuration \( y_1^* \), then renormalizing the probability distribution to the domain \( D \setminus \{ y_1^* \} \) to sample \( y_2^* \), et cetera. The naïve method we propose is similar to Wong and Easton (1980), but instead of an arbitrary binary tree, we use the tree structure (Figure 4) induced by ancestral sampling from the graphical model. This method was concurrently proposed by Shi et al. (2020).

We represent configurations \( y \in D \) as ‘buckets’ on a horizontal axis, with size equal to their probability \( p(y) \). Sampling from \( p(y) \) is equivalent to uniformly selecting a point on this axis and returning the configuration of the corresponding bucket. See an example in Figure 8. Sampling without replacement means that to obtain the second sample, we should remove the bucket corresponding to the first sample, as indicated by the dark shading in Figure 8, and sample uniformly a position on the horizontal axis that is not shaded. When using ancestral sampling, we can determine the bucket by first determining the larger bucket, corresponding to the assignment of the first variable \( y_1 \), then sequentially narrowing it down to the final bucket by sampling \( y_2, \ldots, y_T \). Note that, although this consists of multiple sampling steps, this can be done using a single random number, as illustrated by the dotted
line in Figure 8, such that the result will be equivalent to the result without ancestral sampling with the same random number.

When sampling without replacement, we desire unique complete configurations, but we may have duplicate partial configurations. This means that if we desire to use naïve ancestral sampling to obtain a second sample without replacement, we cannot remove any partial configurations from the domain, but we need to remove the probability mass already sampled from the partial configurations, as illustrated in Figure 8. We can keep track of the probability mass already sampled for each partial configuration by building a prefix tree (or trie) representing the samples, removing the probability mass from each partial configuration by backtracking the path that was used to generate each sample. Each next sample can then be generated by sampling a path down this tree using the adjusted probabilities to sample without replacement, until we arrive at a partial configuration that has not been sampled before, which means that the remaining variables in the configuration can be sampled straightforwardly and the tree is grown to include this path.

Similar to ancestral Gumbel-Top-k sampling, naïve ancestral sampling without replacement has the benefits that it can reuse model evaluations for partial configurations and generate samples one at the time. As downside, this naïve method requires more iterations as it starts from the root for each sample, whereas ancestral Gumbel-Top-k sampling directly jumps to a partial configuration from the priority queue. Also, this naïve method requires careful administration of a complete tree structure whereas ancestral Gumbel-Top-k sampling only requires to keep track of a queue of partial configurations (without any tree structure). Lastly, naïve ancestral sampling without replacement is inherently sequential.

Figure 8: Sampling from a discrete domain visualized as selecting a point uniformly on a horizontal axis. The bucket (with size equal to sampling probability) which contains the random point is the sample. Ancestral sampling can be seen as sequentially narrowing down the ‘bucket’ that contains the sampled point, indicated by light shading. Sampling without replacement requires removing probability mass corresponding to the buckets that have already been sampled (indicated by dark shaded areas) and selecting the random point uniformly from the remaining area. When using ancestral sampling ‘naïvely’, we can use a tree to keep track of the probability mass that should be removed for partial configurations that have been sampled before.
4.6. Weighted Reservoir Sampling with Exponential Jumps

Weighted reservoir sampling (Efraimidis and Spirakis, 2006) is an algorithm for sampling without replacement from a stream of items given weights (unnormalized probabilities) to sample each item. Mathematically, it is equivalent to Gumbel-Top-k sampling (see Section 2.4), but it is executed in a streaming manner, i.e. perturbing (unnormalized) log-probabilities as they arrive while keeping track of the k largest perturbations so far.

In the streaming implementation, each next item replaces an item in a priority queue if its perturbed log-probability exceeds the k-th largest so far. Efraimidis and Spirakis (2006) note that the total weight that is ‘skipped’ before a next item enters the queue follows an exponential distribution, which can also be sampled directly, such that one can directly make an ‘exponential jump’ to the next item to insert the queue, without sampling the perturbations for each item individually. This can be seen as a form of top-down sampling.

Using weighted reservoir sampling, or equivalently, ‘streaming’ Gumbel-Top-k sampling, we can derive yet another algorithm for sampling without replacement from a Bayesian Network, by iterating over the complete domain (leaf nodes in the tree diagram in Figure 4) using a tree traversal algorithm, and keeping track of the k-largest perturbed log-probabilities so far. While this needs to enumerate the complete domain, we can use the ‘exponential jumps’ to directly jump to the next sequence in the domain. In doing so, complete subtrees can be skipped directly as the total weight (probability mass) in the subtree is given by the probability of the partial configuration corresponding to the root of the subtree. This means it can be executed without instantiating the complete tree. Additionally, traversing the tree by expanding the highest probability children first will limit the number of additions to the priority queue as more likely elements are inserted earlier.

While this algorithm is closely connected to our method, it is not equivalent. In particular, using streaming Gumbel-Top-k sampling with exponential jumps, the queue will be initially filled with k complete configurations, which get replaced later by other complete configurations with higher perturbed log-probabilities. The result can only be returned after the complete domain has been traversed. By contrast, ancestral Gumbel-Top-k sampling maintains partial configurations in the queue, which can be used to bound parts of the domain, and additionally returns samples without replacement as they are found, in order.

5. Experiments

This section presents the experiments and results.

5.1. Different Methods for Sampling Without Replacement

In our first experiment, we analyze different methods for sampling without replacement:

- **Ancestral Gumbel-Top-k sampling** (Section 3), where we experiment with different values of m to control the parallelizability of the algorithm.

- **Rejection sampling** (Section 4.4) which generates samples with replacement (using standard ancestral sampling) sequentially and rejects duplicates. We also implement a ‘parallel’ version of this, that generates m samples with replacement in parallel, then rejects the duplicates and repeats this procedure until k unique samples are found.
• *Naïve ancestral sampling* without replacement (Section 4.5). This algorithm is inherently sequential, but we also implement a ‘naïve’ parallelizable version similar to rejection sampling. This version generates \( m \) samples (with replacement) in parallel, removes duplicates and then constructs the tree in Figure 8 to remove the corresponding probability mass from the distribution. Then it uses this tree to generate (again in parallel) \( m \) new samples, which are distinct from the first \( m \) samples but may contain duplicates, which are again removed, after which the tree is updated. This is repeated until in total \( k \) unique samples have been generated.

5.1.1. Data Generation

We generate a random Bayesian Network of 10 variables, where each variable \( y_i \) has a Bernoulli distribution. We generate the variables in topological order, and variable \( i \) has a dependency on variable \( j < i \) (so \( j \in \text{pa}(i) \)) with probability \( c \in [0,1] \), where \( c \) is the *connectivity* factor of the graph. For \( c = 0 \), the resulting graph will be fully independent (Figure 1a), while for \( 0 < c < 1 \) the result will (most likely) be a sparse Bayesian network (Figure 1c) and for \( c = 1 \) the result is a ‘fully connected’ sequence model (Figure 1d). The Bernoulli distribution for \( y_i \in \{0,1\} \) is a mixture between an independent prior and a distribution depending on the parents, which is given by

\[
P(y_i = 1 | y_{\text{pa}(i)}) = \alpha_i p_i^{y_{\text{pa}(i)}} + (1 - \alpha_i)p_i.
\]

Here \( \alpha_i \sim \text{Uniform}(0,1) \) is a parameter that determines how much the variable \( y_i \) is influenced by its parents \( \text{pa}(i) \), \( p_i \sim \text{Uniform}(0,1) \) is the ‘independent’ or ‘prior’ probability that \( y_i = 1 \), and \( p_{y_{\text{pa}(i)}} \sim \text{Uniform}(0,1) \forall y_{\text{pa}(i)} \in D_{\text{pa}(i)} \) determines the influence of the parents on the probability that \( y_i = 1 \), given as a table for each of the \( 2^{|\text{pa}(v)|} \) possible values in \( D_{\text{pa}(i)} \).

5.1.2. Number of Iterations

First we set \( m = 1 \), making all algorithms fully sequential. For our algorithm, we measure the number of iterations (removing \( m = 1 \) element from the queue and adding its expansions), which is equal to the number of model evaluations. For rejection sampling and the naïve ancestral sampling algorithm, we count as iteration the sampling of one variable, such that the number of iterations is also equal to the number of model evaluations and generating a single sample \( y \) with \( T \) variables takes \( T \) iterations/model evaluations. In Figure 9a we compare the number of iterations (model evaluations) needed for generating different numbers of samples using the different methods. We clearly observe how our method requires the fewest iterations/model evaluations to generate the same number of samples.

5.1.3. Parallelizing the Algorithms

By choosing the number of parallel processors \( m \) we can trade off the total time (iterations) it takes to run the algorithm, also known as the *depth* or *span*, against the total cost of running the algorithm, which is \( m \) times the depth or span. Figure 9b shows for different numbers of parallel processors how many iterations it takes to generate \( k = 100 \) samples, i.e. it summarizes the result for the previous experiment run for different values of \( m \). It is clear how increasing the parallelization quickly decreases the number of iterations required although the effect diminishes as \( m \) increases.
Ancestral Gumbel-Top-$k$ Sampling

Figure 9: Results of sampling from 100 randomly generated models (Bayesian networks) with 10 Bernoulli variables each and a connectivity of $c = 0.5$. For each model, we generated $k = 100$ samples without replacement, using different methods and different numbers of parallel processors $m$. Rejection sampling and naïve ancestral sampling show a peak at $m = 50$ since they typically generate $3 \times 50 = 150$ samples (cost 1500) to obtain 100 unique samples, whereas $m = 40$ and $m = 60$ can suffice with generating $3 \times 40$ or $2 \times 60 = 120$ samples. Results are presented as mean and standard deviation over the 100 different models. Results with $c = 0$ or $c = 1$ (not shown) are similar in terms of iterations and costs.
To visualize the overhead from the parallelization, Figure 9c visualizes the total cost, which is the number of iterations (Figure 9b) multiplied by $m$. For ancestral Gumbel-Top-$k$ sampling, the increased costs for larger $m$ has two different sources as explained in Section 3.8: redundant model evaluations of partial configurations which would have been pushed of the queue in the sequential setting, and ‘idle’ parallel resources if the queue size is smaller than $m$. For rejection sampling, the parallel version is suboptimal because too many samples may be generated since they are generated in batches of size $m$. For the naïve ancestral sampling algorithm, by the batching in parallel, there may still be duplicates compared to the sequential ($m = 1$) algorithm, and more than $k$ samples may be generated in total.

Ultimately, $m$ determines the trade-off between the number of iterations and the total cost of the algorithm, which is visualized in Figure 9d. Our algorithm with $m = k = 100$ (i.e. Stochastic Beam Search) uses 10 iterations (since there are 10 variables) and has a cost of $10 \times 100 = 1000$, which is the minimum for naïve ancestral sampling and rejection sampling. On the other extreme, the sequential algorithm with $m = 1$ uses around 300 iterations with a cost of 300. The difference is large because of the small setting, where, with $m = 100$, many processors are idle in 6 of the 10 iterations since $2^6 = 64 < 100$. We expect the difference to be smaller in real applications, but this experiment clearly shows that there is a trade-off, and suggests that limited parallelization (in this case $m = 10$) rapidly decreases the number of iterations without increasing total cost too much.

5.1.4. Variable Selection Strategy

As explained in Section 3.9, we can select any variable from $\{v \in V \setminus S : \text{pa}(v) \subseteq S\}$ to sample in each iteration. We experiment with the following variable selection strategies:

- **Fixed** uses the order in which the variables are generated. Since we generated the Bayesian network in topological order, this is valid.

- **Random** chooses the variable from $\{v \in V \setminus S : \text{pa}(v) \subseteq S\}$ to sample next uniformly at random.

- **Minimum entropy** computes the (conditional) entropy for all variables $\{v \in V \setminus S : \text{pa}(v) \subseteq S\}$ as $-\hat{p}_v \log(\hat{p}_v) - (1 - \hat{p}_v) \log(1 - \hat{p}_v)$, where $\hat{p}_v = P(y_v = 1|y_{\text{pa}(v)})$ and selects the variable $v$ with minimum entropy.

- **Maximum entropy** selects the variable $v$ with maximum entropy.

We note that selecting $v$ based on entropy requires the model to be evaluated, which may be undesirable in some cases. However, theoretically these model evaluations can be cached/reused as most evaluations will be required eventually to sample the remaining variables. In the extreme case of $c = 0$, i.e. the fully independent model in Figure 1a, we can simply precompute all model evaluations (which is a good idea anyway), and using the entropy variable selection strategy reduces to sorting the variables by their entropy before starting the Gumbel-Top-$k$ sampling algorithm.

In Figure 10a we clearly see that selecting variables with minimum entropy first is the best strategy (it requires fewest iterations), whereas selecting based on maximum entropy performs worst, with a random strategy or using the fixed generation order (which can also be considered random) in between. As explained in Section 3.9, this is because selecting
variables with low entropy allows maximum reuse of partial configurations. In Figure 10b we show how the difference between different variable selection strategies decreases as there are more dependencies between variables, where for a ‘fully connected’ sequence model \((c = 1, \text{Figure 1d})\) there is no freedom in variable selection and all strategies perform the same.

5.2. Diverse Beam Search

This experiment was originally presented in Kool et al. (2019b). The results were obtained using Stochastic Beam Search, but are valid for any implementation of ancestral Gumbel-Top-\(k\) sampling with parallelism factor \(m < k\). In this experiment we compare Stochastic Beam Search as a principled (stochastic) alternative to Diverse Beam Search (Vijayakumar et al., 2018) in the context of neural machine translation to obtain a diverse set of translations for a single source sentence \(x\). Following the setup by Vijayakumar et al. (2018) we report both diversity as measured by the fraction of unique \(n\)-grams in the \(k\) translations as well as mean and maximum BLEU score (Papineni et al., 2002) as an indication of the quality of the sample. The maximum BLEU score corresponds to ‘oracle performance’ reported by Vijayakumar et al. (2018), but we report the mean as well since a single good translation and \(k - 1\) completely random sentences scores high on both maximum BLEU score and diversity, while being undesirable. A good method should increase diversity without sacrificing mean BLEU score.

We compare four different sentence generations methods: Beam Search (BS), Sampling (with replacement), Stochastic Beam Search (SBS) (sampling without replacement) and Diverse Beam Search with \(G\) groups (DBS(G)) (Vijayakumar et al., 2018). For Sampling and Stochastic Beam Search, we control the diversity of samples generated using the softmax temperature \(\tau\) (see Equation 2) used to compute the model probabilities. We use
\( \tau = 0.1, 0.2, \ldots, 0.8 \), where a higher \( \tau \) results in higher diversity. Heuristically, we also vary \( \tau \) for computing the scores with (deterministic) Beam Search. The diversity of Diverse Beam Search is controlled by the diversity strengths parameter, which we vary between 0.1, 0.2, \ldots, 0.8. We set the number of groups \( G \) equal to the sample size \( k \), which Vijayakumar et al. (2018) reported as the best choice.

We modify the Beam Search in fairseq (Ott et al., 2019) to implement Stochastic Beam Search\(^4\), and use the fairseq implementations for Beam Search, Sampling and Diverse Beam Search. For theoretical correctness of the Stochastic Beam Search, we disable length-normalization (Wu et al., 2016) and early stopping (and therefore also do not use these parameters for the other methods). We use the pretrained model from Gehring et al. (2017) and use the wmt14.v2.en-fr.newstest2014 test set\(^5\) consisting of 3003 sentences. We run the four methods with sample sizes \( k = 5, 10, 20 \) and plot the minimum, mean and maximum BLEU score among the \( k \) translations (averaged over the test set) against the average \( d = \frac{1}{4} \sum_{n=1}^{4} d_n \) of 1, 2, 3 and 4-gram diversity, where \( n \)-gram diversity is defined as

\[
d_n = \frac{\text{# of unique } n\text{-grams in } k \text{ translations}}{\text{total # of } n\text{-grams in } k \text{ translations}}.
\]

In Figure 11, we represent the results as curves, indicating the trade-off between diversity and BLEU score. The points indicate datapoints and the dashed lines indicate the (averaged) minimum and maximum BLEU score. For the same diversity, Stochastic Beam Search achieves higher mean/maximum BLEU score. Looking at a certain BLEU score, we observe that Stochastic Beam Search achieves the same BLEU score as Diverse Beam Search with a significantly larger diversity. For low temperatures (\(< 0.5\)), the maximum BLEU score of Stochastic Beam Search is comparable to the deterministic Beam Search, so the increased diversity does not sacrifice the best element in the sample. Note that Sampling achieves higher mean BLEU score at the cost of diversity, which may be because good translations are sampled repeatedly. However, the maximum BLEU score of both Sampling and Diverse Beam Search is lower than with Beam Search and Stochastic Beam Search.

---

\(\text{Figure 11: Minimum, mean and maximum BLEU score vs. diversity for different sample sizes } k.\) Points indicate different temperatures/diversity strengths, from 0.1 (low diversity, left in graph) to 0.8 (high diversity, right in graph).

---

\(^4\) Our code is available at https://github.com/wouterkool/stochastic-beam-search.

5.3. BLEU Score Estimation

In our second experiment, also presented originally in Kool et al. (2019b), we use sampling without replacement to evaluate the expected sentence level BLEU score for a translation \( y \) given a source sentence \( x \). Although we are often interested in corpus level BLEU score, estimation of sentence level BLEU score is useful, for example when training using minibatches to directly optimize BLEU score (Ranzato et al., 2016). We leave dependence of the BLEU score on the source sentence \( x \) implicit, and write \( f(y) = \text{BLEU}(y, x) \). We want to estimate the following expectation:

\[
\mathbb{E}_{y \sim p(y|x)}[f(y)] = \sum_{y \in D} p(y|x) f(y).
\]  

Under a Monte Carlo (MC) sampling with replacement scheme with size \( k \), we write \( S \) as the set of sampled sequences and estimate (19) using

\[
\mathbb{E}_{y \sim p(y|x)}[f(y)] \approx \frac{1}{k} \sum_{y \in S} f(y).
\]  

If the distribution \( p(y|x) \) has low entropy (for example if there are only few valid translations), then MC estimation may be inefficient since repeated samples are uninformative. We can use sampling without replacement as an improvement, but we need to use importance weights to correct for the changed sampling probabilities. Using Gumbel-Top-\( k \) sampling, we can implement an estimator equivalent to the estimator described by Vieira (2017), which can be seen as a Horvitz-Thompson estimator (Horvitz and Thompson, 1952) used with priority sampling (Duffield et al., 2007):

\[
\mathbb{E}_{y \sim p_y(y|x)}[f(y)] \approx \sum_{y \in S} p(y|x) \frac{q_{\theta,\kappa}(y|x)}{q_{\theta,a}(y|x)} f(y).
\]  

Using this estimator, we ‘sacrifice’ the \( k \)-th sample to obtain the empirical threshold \( \kappa \) (which is the \( k \)-th largest perturbed log-probability, see Section 3.7), and we define \( S \) as the set of the \( k - 1 \) sequences corresponding to the \( k - 1 \) largest perturbed log-probabilities. It holds that \( y \in S \) if \( G_{\phi_y} > \kappa \), which highlights the relation to threshold sampling (see Section 4.3). We define

\[
q_{\theta,a}(y|x) = P(G_{\phi_y} > a) = 1 - \exp(-\exp(\phi_y - a)).
\]  

If we would assume a fixed threshold \( a \) and variably sized sample \( S = \{ y \in D : G_{\phi_y} > a \} \), then \( q_{\theta,a}(y|x) = P(y \in S) \) and \( \frac{p(y|x)}{q_{\theta,a}(y|x)} \) is a standard importance weight. Surprisingly, using a fixed sample size \( k \) (and empirical threshold \( \kappa ) \) also yields in an unbiased estimator, and we include a proof adapted from Duffield et al. (2007) and Vieira (2017) in Appendix B.

Empirically, the estimator (21) has high variance, and in practice we find it is preferred to normalize the importance weights by \( W(S) = \sum_{y \in S} \frac{p(y|x)}{q_{\theta,\kappa}(y|x)} \) (Hesterberg, 1988):

\[
\mathbb{E}_{y \sim p_y(y|x)}[f(y)] \approx \frac{1}{W(S)} \sum_{y \in S} \frac{p(y|x)}{q_{\theta,\kappa}(y|x)} f(y).
\]  

---

6. Formally, when sampling with replacement, \( S \) is a multiset.
The estimator (23) is biased but consistent: in the limit $k = |D|$ we sample the entire domain, so we have empirical threshold $\kappa = -\infty$ and $q_{\theta, \kappa}(y|x) = 1$ and $W(S) = 1$, such that (23) is equal to (19).

We have to take care computing the importance weights as depending on the entropy the terms in the quotient $\frac{p_{\theta}(y|x)}{q_{\theta, \infty}(y|x)}$ can become very small, and in our case the computation of $P(G_{\phi_y} > a) = 1 - \exp(-\exp(\phi_y - a))$ can suffer from catastrophic cancellation. For details, see Appendix C.

Because the model is not trained to use its own predictions as input, at test time errors can accumulate. As a result, when sampling with the default temperature $\tau = 1$, the expected BLEU score is very low (below 10). To improve quality of generated sentences we use lower temperatures and experiment with $\tau = 0.05, 0.1, 0.2, 0.5$. We then use different methods to estimate the BLEU score:

- **Monte Carlo** (MC), using Equation (20).
- **Stochastic Beam Search** (SBS), where we compute estimates using the estimator in Equation (21) and the normalized variant in Equation (23).
- **Beam Search** (BS), where we compute a deterministic beam $S$ (the temperature $\tau$ affects the scoring) and compute $\sum_{y \in S} p_{\theta}(y|x)f(y)$. This is not a statistical estimator, but a lower bound to the target (19) which serves as a validation of the implementation and gives insight on how many sequences we need at least to capture most of the mass in (19). We also compute the normalized version $\sum_{y \in S} \frac{p_{\theta}(y|x)f(y)}{\sum_{y \in S} p_{\theta}(y|x)}$, which can heuristically be considered as a ‘determinstic estimate’.

In Figure 12 we show the results of computing each estimate 100 times (BS only once as it is deterministic) for three different sentences\footnote{These are sentences 1500, 2000 and 2500 from the WMT14 data set.} for temperatures $\tau = 0.05, 0.1, 0.2, 0.5$ and sample sizes $k = 1$ to 250. We report the empirical mean and 2.5-th and 97.5-th percentile. The normalized SBS estimator indeed achieves significantly lower variance than the unnormalized version and for $\tau < 0.5$, it significantly reduces variance compared to MC, without adding observable bias. For $\tau = 0.5$ the results are similar, but we are less interested in this case as the overall BLEU score is lower than for $\tau = 0.2$.

### 5.4. Conditional Entropy Estimation

Additionally to estimating the BLEU score we use $f(y) = -\log p_{\theta}(y|x)$ such that Equation (19) becomes the model entropy (conditioned on the source sentence $x$)

$$E_{y \sim p_{\theta}(y|x)} \left[ -\log p_{\theta}(y|x) \right].$$

Entropy estimation is useful in optimization settings where we want to include an entropy loss to ensure diversity. It is a different problem than BLEU score estimation as high BLEU score (for a good model) correlates positively with model probability, while for entropy rare $y$ contribute the largest terms $-\log p_{\theta}(y|x)$. We use the same experimental setup as for the BLEU score and present the results in Figure 13. The results are similar to the BLEU score experiment: the normalized SBS estimate has significantly lower variance than MC for $\tau < 0.5$ while for $\tau = 0.5$, results are similar. This shows that Stochastic Beam Search can be used to construct practical statistical estimators.
Ancestral Gumbel-Top-k Sampling

Figure 12: BLEU score estimates for three sentences sampled/decoded by different estimators for different temperatures.

Figure 13: Entropy estimates for three sentences sampled/decoded by different estimators for different temperatures.
6. Possible Extensions of Ancestral Gumbel-Top-k Sampling

Our algorithm can be extended in various ways, to give efficient implementations of two existing algorithms or to perform efficient sampling with replacement.

6.1. Memory Augmented Policy Optimization

Our algorithm can be extended to give an efficient implementation of Memory Augmented Policy Optimization (MAPO) (Liang et al., 2018), an extension of REINFORCE (Williams, 1992) that optimizes a policy in reinforcement learning using a replay buffer of good experiences. MAPO computes an exact gradient for the $b$ experiences (configurations) in the buffer, which requires $b$ model evaluations, and uses a sample from the model outside of the buffer, obtained using (potentially inefficient) rejection sampling. We can adapt ancestral Gumbel-Top-k sampling to implement MAPO efficiently, by modifying the priority queue to use a hierarchical criterion, that will put in the front partial configurations which correspond to partial configurations in the buffer, and sort the remaining configurations based on the perturbed log-probability. When using $k = b + 1$, the result will be the $b$ configurations in the buffer (with their model evaluations), and a single sample outside of the buffer.

6.2. Rao-Blackwellized Stochastic Gradients

While MAPO is an estimator that computes an exact gradient for a number of ‘good’ configurations, Liu et al. (2019) proposed a similar estimator that computes an exact gradient for high probability configurations (instead of ‘good’ ones), which may have a greater contribution overall. While they consider 1-dimensional categorical distributions only, we note that the estimator can also be used in multi-dimensional settings, where high probability categories can be found by an approximate algorithm such as (deterministic) beam search, but in this case it is challenging to obtain a sample from the remaining domain.

Similar to MAPO, we can make a modification of our algorithm, to yield both the high probability configurations, and a sample from the remaining domain. In particular, we can modify Stochastic Beam Search (our algorithm with $m = k$) to maintain in the queue the $k - 1$ partial configurations with highest log-probability (without Gumbel perturbation) and one partial configuration which has the highest perturbed log-probability among the remaining configurations. This means that we have a ‘deterministic’ beam of size $k - 1$ as well as a single sample (partial configuration) outside of the beam.

In general, it may be preferred to have fewer exact and more sampled configurations; see the relevant discussions in Fearnhead and Clifford (2003); Liu et al. (2019). Using our algorithm, we could maintain $k - \ell$ configurations by their log-probability, and $\ell$ configurations by the perturbed log-probability. The resulting $\ell$ samples without replacement can be converted to (at least) $\ell$ samples with replacement using resampling (see below), or one can use an estimator based on sampling without replacement (Kool et al., 2019a, 2020).

---

8. Existence of a partial configuration/trajectory in the buffer can be efficiently checked by representing the replay buffer as a prefix tree (trie), similar to Figure 8.

9. The result may be slightly different than deterministic beam search with size $k - 1$, since extensions of the $k$-th configuration (the sample) may have high log-probability and push other configurations off the beam. If desired, this can be heuristically prevented.
6.3. Sampling with Replacement with Sampling Without Replacement

Sampling without replacement can be used to sample with replacement by using a resampling algorithm. The basic idea is that using ancestral Gumbel-Top-$k$ sampling (Algorithm 1) ‘lazily’, we can obtain the first sample $y_1^*$, for which sampling with/without replacement is the same. Then for the second sample, we can take $y_1^*$ with probability $p(y_1^*)$, or sample from $D \setminus \{y_1^*\}$ with probability $1 - p(y_1^*)$, which can be done by continuing incremental sampling without replacement. At any point, if $y_1^*, ..., y_n^*$ are the samples without replacement so far, we can get another sample with replacement by choosing $y_1^*$ with probability $p(y_1^*)$, $y_2^*$ with probability $p(y_2^*)$, et cetera or find the next sample $y_{n+1}^*$ (without replacement) with probability $1 - \sum_{i=1}^{n} p(y_i^*)$. Repeating this algorithm to obtain a desired number of samples has a lower computational cost (as measured by model evaluations) than standard sampling with replacement, as resampling is cheap (it does not require additional model evaluations) and the inner ancestral Gumbel-Top-$k$ sampling algorithm uses model evaluations efficiently.

7. Discussion

We introduced ancestral Gumbel-Top-$k$ sampling, an algorithm that allows to efficiently draw samples without replacement from a probability distribution represented as a Bayesian network. It allows to control the parallelizability of the algorithm, trading off the number of required iterations versus the total cost of running the algorithm, such that it can make efficient use of modern hardware. We discussed possible extensions of the algorithm, which allow implementations of the gradient estimators by Liang et al. (2018) and Liu et al. (2019) and a resampling algorithm to efficiently obtain samples with replacement.

We have discussed related algorithms and empirically shown the benefits of using ancestral Gumbel-Top-$k$ sampling: it enables to generate samples without replacement using a significantly lower computational cost than alternatives. We have analyzed the influence of the number of parallel processors experimentally, suggesting that limited parallelism quickly decreases the number of required iterations without increasing total cost too much. Additionally, we have shown how selecting the order of sampling based on entropy can reduce the cost of the algorithm, which is especially useful for sampling from fully independent models where the ‘optimal’ order of sampling can be determined upfront.

In the context of sequence models, Gumbel-Top-$k$ sampling relates to sampling (with replacement) and (deterministic) beam search and can be seen as a method that combines the advantages of both. Our experiments support the idea that Gumbel-Top-$k$ sampling can be used as a drop-in replacement in places where sampling (with replacement) or (deterministic) beam search is used. In fact, our experiments show that sampling without replacement can be used to yield lower-variance estimators and high-diversity samples from a neural machine translation model.

We hope that our method motivates future work to develop improved statistical learning methods, especially in the context of Deep Learning, based on sampling without replacement, a direction of research that has, in fact, already started (Kool et al., 2019a, 2020). We believe that ancestral Gumbel-Top-$k$ sampling has potential to increase both computational and statistical efficiency in Deep Learning applications that involve discrete computations, such as neural machine translation (Sutskever et al., 2014; Bahdanau et al., 2015) and image captioning (Vinyals et al., 2015).
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Appendix A. Sampling a Set of Gumbels with Maximum \( k \)

As explained in Section 3.4, to sample a set of Gumbels with maximum \( T \), let \( G_{\phi_i} \sim \text{Gumbel}(\phi_i) \), let \( Z = \max_i G_{\phi_i} \) and define

\[
\tilde{G}_{\phi_i} = F_{\phi_i,T}^{-1}(F_{\phi_i,Z}(G_{\phi_i})) = -\log(\exp(-T) - \exp(-Z) + \exp(-G_{\phi_i})).
\]  

(24)

Direct computation of (24) can be unstable as large terms need to be exponentiated. Instead, we compute

\[
v_i = T - G_{\phi_i} + \log1mexp(G_{\phi_i} - Z),
\]
\[
\tilde{G}_{\phi_i} = T - \max(0, v_i) - \log1pexp(-|v_i|)
\]

where we have defined

\[
\log1mexp(a) = \log(1 - \exp(a)), \quad a \leq 0
\]
\[
\log1pexp(a) = \log(1 + \exp(a)).
\]

This is equivalent as

\[
T - \max(0, v_i) - \log(1 + \exp(-|v_i|))
\]
\[
= T - \log(1 + \exp(v_i))
\]
\[
= T - \log(1 + \exp(T - G_{\phi_i} + \log(1 - \exp(G_{\phi_i} - Z))))
\]
\[
= T - \log(1 + \exp(T - G_{\phi_i})(1 - \exp(G_{\phi_i} - Z)))
\]
\[
= T - \log(1 + \exp(T - G_{\phi_i}) - \exp(T - Z))
\]
\[
= -\log(\exp(-T) + \exp(-G_{\phi_i}) - \exp(-Z))
\]
\[
= \tilde{G}_{\phi_i}
\]

The first step can be easily verified by considering the cases \( v_i < 0 \) and \( v_i \geq 0 \). \( \log1mexp \) and \( \log1pexp \) can be computed accurately using \( \log1p(a) = \log(1 + a) \) and \( \expm1(a) = \exp(a) - 1 \) (Mächler, 2012):

\[
\log1mexp(a) = \begin{cases} 
\log(-\expm1(a)) & a > -0.693 \\
\log1p(-\exp(a)) & \text{otherwise},
\end{cases}
\]
\[
\log1pexp(a) = \begin{cases} 
\log1p(\exp(a)) & a < 18 \\
x + \exp(a) & \text{otherwise}.
\end{cases}
\]
Appendix B. Unbiasedness of the Importance Weighted Estimator

We give a proof of unbiasedness of the importance weighted estimator, which is adapted from the proofs by Duffield et al. (2007) and Vieira (2017). For generality of the proof, we enumerate categories in the domain $D$ by $i = 1, \ldots, n$ and we consider general random keys $h_i$ for $i = 1, \ldots, n$ (not necessarily Gumbel perturbations). As was noted by Vieira (2017), the actual distribution of the keys does not influence the unbiasedness of the estimator, but does determine the effective sampling scheme. Using Gumbel perturbed log-probabilities as keys (e.g. $h_i = G_{\phi_i}$) is equivalent to the PPSWOR scheme described by Vieira (2017). For simplicity, we write $p_i = p_{\theta}(i)$ and $q_i(a) = q_{\theta,a}(i) = P(h_i > a)$ (see Equation 22), and we define $h_{1:n} = \{h_1, \ldots, h_n\}$ and $h_{-i} = \{h_1, \ldots, h_{i-1}, h_{i+1}, \ldots, h_n\} = h_{1:n} \setminus \{h_i\}$.

Given a fixed threshold $a$, it holds that $q_i(a) = P(i \in S)$ is the probability that category $i$ is included in the sample $S$, so it can be thought of as the inclusion probability of $i$. Given a fixed sample size $k$, let $\kappa$ be the $(k+1)$-th largest element of $h_{1:n}$, so $\kappa$ is the empirical threshold. Let $\kappa'_i$ be the $k$-th largest element of $h_{-i}$ (the $k$-th largest of all other elements). We first prove Lemma 4, which is then used to prove Theorem 5, which states that the importance weighted estimator is an unbiased estimator of $\mathbb{E}[f(i)]$, for a given function $f(i)$.

**Lemma 4** The expected weight of each term in the importance weighted estimator is 1:

$$\mathbb{E}_{h_{1:n}} \left[ \frac{1 \{i \in S\}}{q_i(\kappa)} \right] = 1.$$

**Proof** We make use of the observation (slightly rephrased) by Duffield et al. (2007) that conditioning on $h_{-i}$, we know $\kappa'_i$, and the event $i \in S$ implies that $\kappa = \kappa'_i$ since $i$ will only be in the sample if $h_i > \kappa'_i$, which means that $\kappa'_i$ is the $(k + 1)$-th largest value of $h_{-i} \cup \{h_i\} = h_{1:n}$. The reverse is also true: if $\kappa = \kappa'_i$ then $h_i$ must be larger than $\kappa'_i$ since otherwise the $(k + 1)$-th largest value of $h_{1:n}$ will be smaller than $\kappa'_i$. By separating the expectation over $h_{1:n}$ it follows that

\[
\mathbb{E}_{h_{1:n}} \left[ \frac{1 \{i \in S\}}{q_i(\kappa)} \right] = \mathbb{E}_{h_{-i}} \left[ \mathbb{E}_{h_i} \left[ \frac{1 \{i \in S\}}{q_i(\kappa)} \right] h_i \right] \\
= \mathbb{E}_{h_{-i}} \left[ \mathbb{E}_{h_i} \left[ \frac{1 \{i \in S\}}{q_i(\kappa)} \right] h_{-i}, i \in S \right] P(i \in S|h_{-i}) + \mathbb{E}_{h_i} \left[ \frac{1 \{i \in S\}}{q_i(\kappa)} \right] h_{-i}, i \notin S \right] P(i \notin S|h_{-i}) \\
= \mathbb{E}_{h_{-i}} \left[ \mathbb{E}_{h_i} \left[ \frac{1}{q_i(\kappa)} \right] h_{-i}, i \in S \right] q_i(\kappa'_i) + 0 \\
= \mathbb{E}_{h_{-i}} \left[ \mathbb{E}_{h_i} \left[ \frac{1}{q_i(\kappa)} \kappa = \kappa'_i \right] q_i(\kappa'_i) \right] \\
= \mathbb{E}_{h_{-i}} \left[ \frac{1}{q_i(\kappa'_i)} q_i(\kappa'_i) \right] = \mathbb{E}_{h_{-i}} [1] = 1.
\]
Theorem 5 The importance weighted estimator is an unbiased estimator of \( E[f(i)] \):

\[
\mathbb{E}_{h_{1:n}} \left[ \sum_{i \in S} \frac{p_i}{q_i(\kappa)} f(i) \right] = \mathbb{E}[f(i)].
\]

Proof

\[
\mathbb{E}_{h_{1:n}} \left[ \sum_{i \in S} \frac{p_i}{q_i(\kappa)} f(i) \right] = \mathbb{E}_{h_{1:n}} \left[ \sum_{i=1}^{n} \frac{p_i}{q_i(\kappa)} f(i) \mathbb{1}_{\{i \in S\}} \right]
= \sum_{i=1}^{n} p_i f(i) \cdot \mathbb{E}_{h_{1:n}} \left[ \mathbb{1}_{\{i \in S\}} \right]
= \sum_{i=1}^{n} p_i f(i) \cdot 1
= \mathbb{E}[f(i)].
\]

Appendix C. Numerical Stability of Importance Weights

We have to take care computing the importance weights as depending on the entropy the terms in the quotient \( \frac{p_i}{q_i(\kappa)} \) (using notation from Appendix B) can become very small, and in our case the computation of \( q_i(\kappa) = 1 - \exp(-\exp(\phi_i - \kappa)) \) (Equation (22)) can suffer from catastrophic cancellation. We can rewrite this expression using the more numerically stable implementation \( \expm1(x) = \exp(x) - 1 \) as \( q_i(\kappa) = -\expm1(-\exp(\phi_i - \kappa)) \) but in some cases this still suffers from instability as \( \exp(\phi_i - \kappa) \) can underflow if \( \phi_i - \kappa \) is small. Instead, for \( \phi_i - \kappa < -10 \) we use the identity

\[
\log(1 - \exp(-z)) = \log(z) - \frac{z}{2} + \frac{z^2}{24} - \frac{z^4}{2880} + O(z^6)
\]

to directly compute the log importance weight using \( z = \exp(\phi_i - \kappa) \) and \( \phi_i = \log p_i \) (we assume \( \phi_i \) is normalized). Since \( q_i(\kappa) = 1 - \exp(-z) \), we have

\[
\log \left( \frac{p_i}{q_i(\kappa)} \right) = \log p_i - \log q_i(\kappa)
= \log p_i - \log (1 - \exp(-z))
= \log p_i - \left( \log(z) - \frac{z}{2} + \frac{z^2}{24} - \frac{z^4}{2880} + O(z^6) \right)
= \log p_i - \left( \phi_i - \kappa - \frac{z}{2} + \frac{z^2}{24} - \frac{z^4}{2880} + O(z^6) \right)
= \kappa + \frac{z}{2} - \frac{z^2}{24} + \frac{z^4}{2880} + O(z^6).
\]

If \( \phi_i - \kappa < -10 \) then \( 0 < z < 10^{-6} \) so this computation will not lose any significant digits.
References


